

# 软件分析

南京大学

计算机科学与技术系

程序设计语言与  
静态分析研究组

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# Static Program Analysis

## Data Flow Analysis — Applications

Nanjing University

Yue Li

2020

# contents

1. Overview of Data Flow Analysis
2. Preliminaries of Data Flow Analysis
3. Reaching Definitions Analysis
4. Live Variables Analysis
5. Available Expressions Analysis

# Data Flow Analysis

# Data Flow Analysis

## How Data Flows on CFG?

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How Data  
Flows through the

CFG (a program)?

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Recall 1<sup>st</sup>  
lecture

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How Data Flows on CFG?

How application-specific Data ← Abstraction

Over-approximation → Flows through the

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Over-approximation → Flows through the

Nodes (BBs/statements) and  
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for most static analyses  
(may analysis)

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## Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction



Over-approximation → Flows through the

Nodes (BBs/statements) and  
Edges (control flows) of  
CFG (a program)?

may analysis:

outputs information that may be true (over-approximation)

must analysis:

outputs information that must be true (under-approximation)

Over- and under-approximations are both for safety of analysis

# Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction

Safe-approximation → Flows through the  
may analysis: safe=over      Nodes (BBs/statements) and  
must analysis: safe=under      Edges (control flows) of  
  CFG (a program)?

# Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction

Safe-approximation → Flows through the

Transfer function → Nodes (BBs/statements) and

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+,-,0, $\perp$ , $\top$

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$+ \ op \ - = - ; + \ op \ + = +$

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different data-flow analysis applications have  
different data abstraction and  
different flow safe-approximation strategies, i.e.,  
different transfer functions and control-flow handlings

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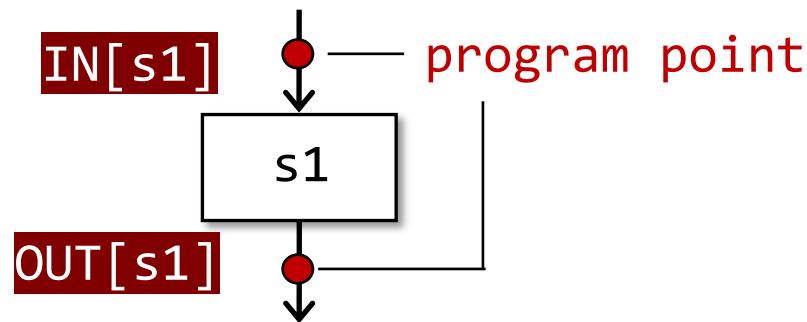
*union the signs at merges*

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# Preliminaries of Data Flow Analysis

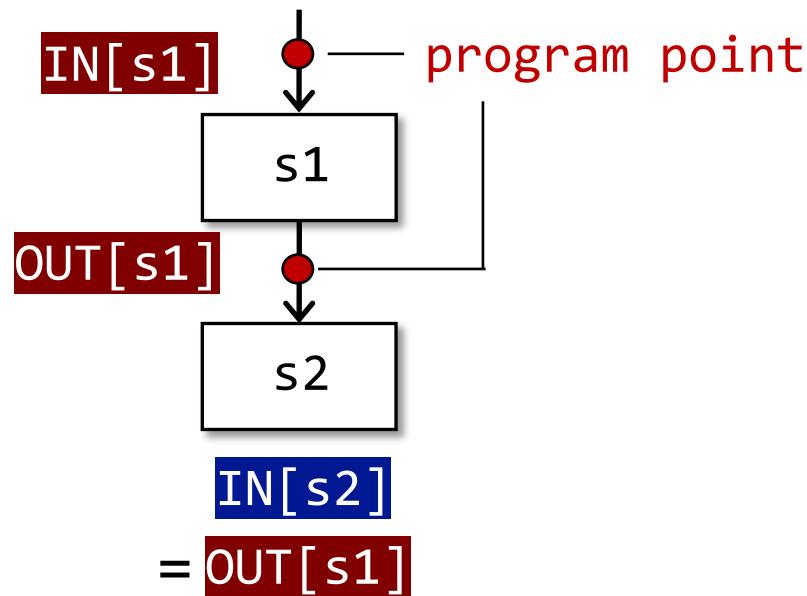
# Input and Output States

- Each execution of an IR statement transforms an **input state** to a new **output state**
- The input (output) state is associated with the **program point** before (after) the statement



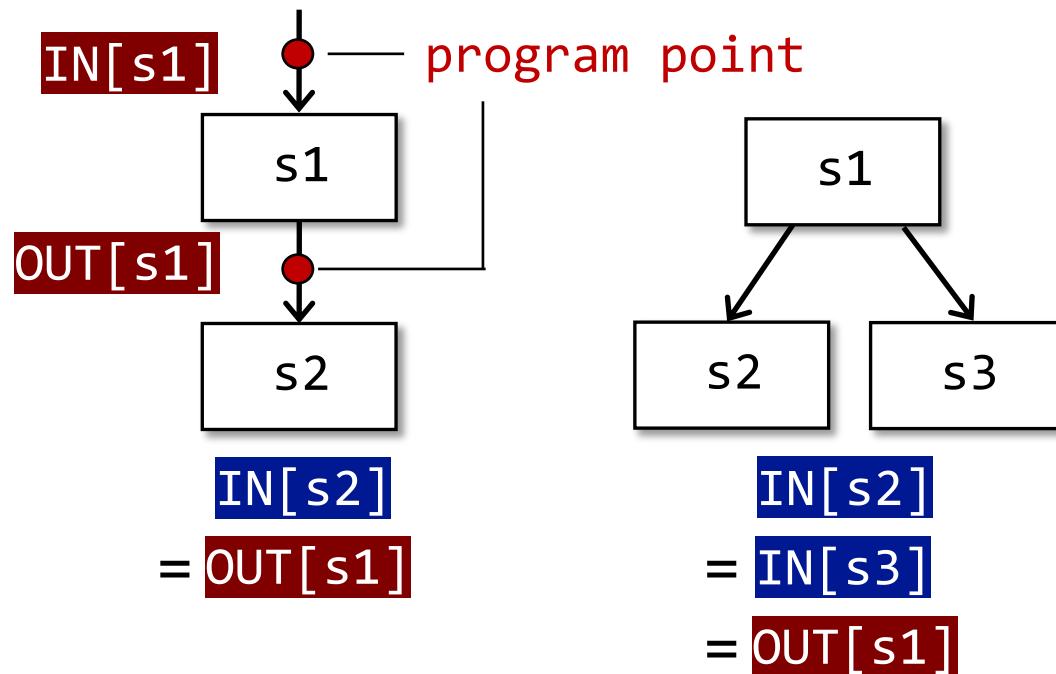
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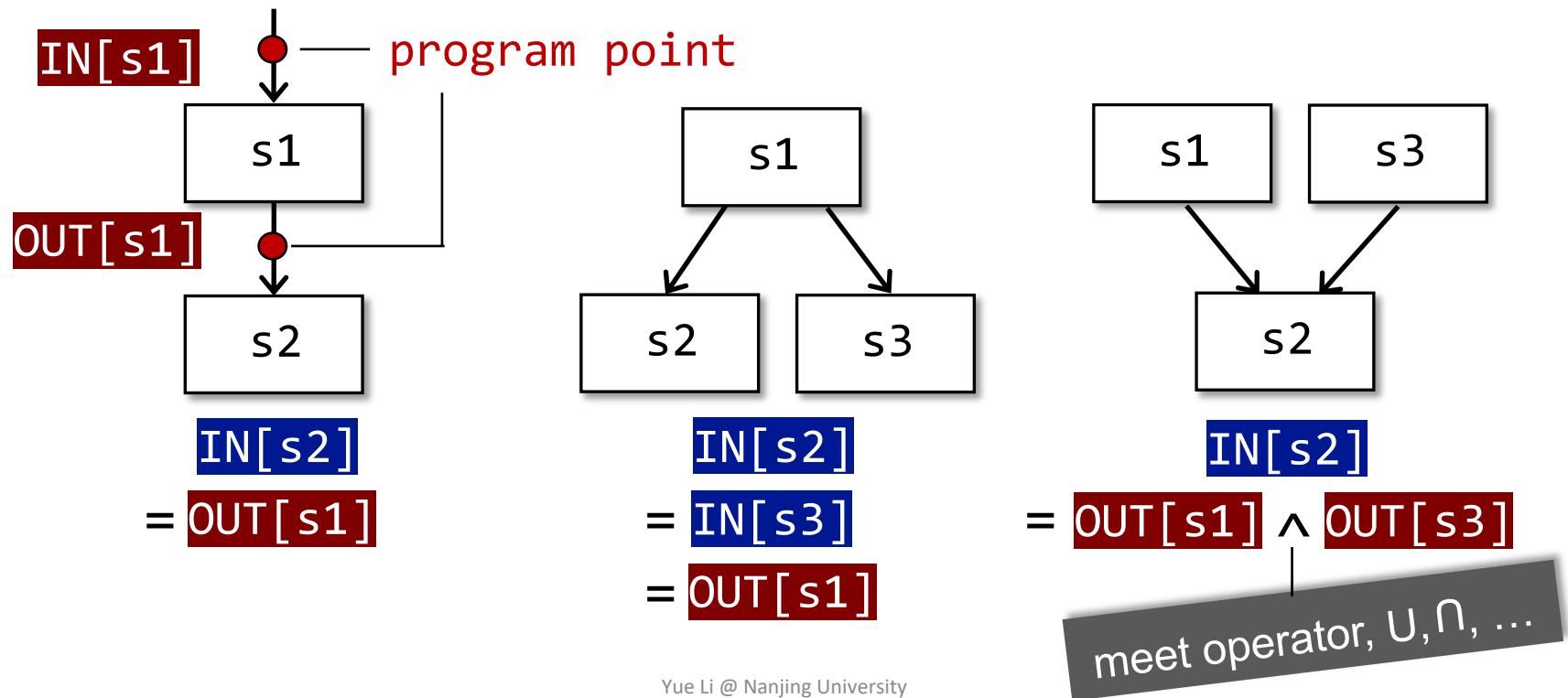
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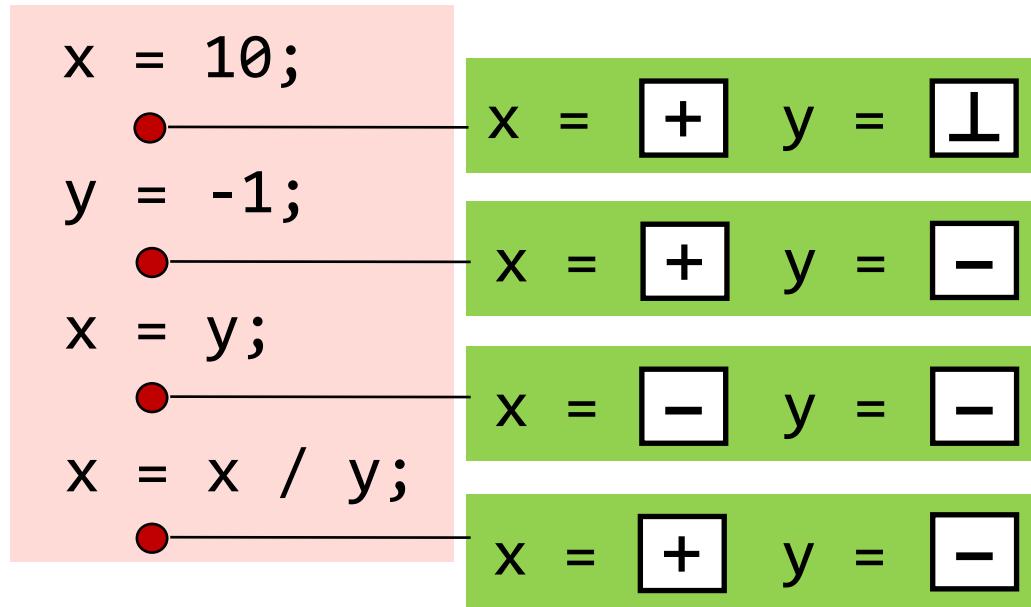
In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.

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```
x = 10;  
●  
y = -1;  
●  
x = y;  
●  
x = x / 6;  
●
```

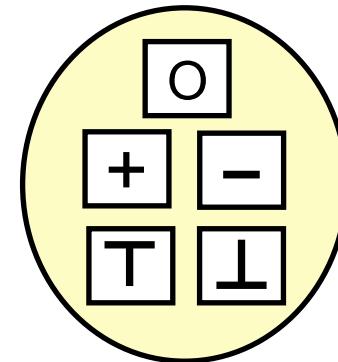
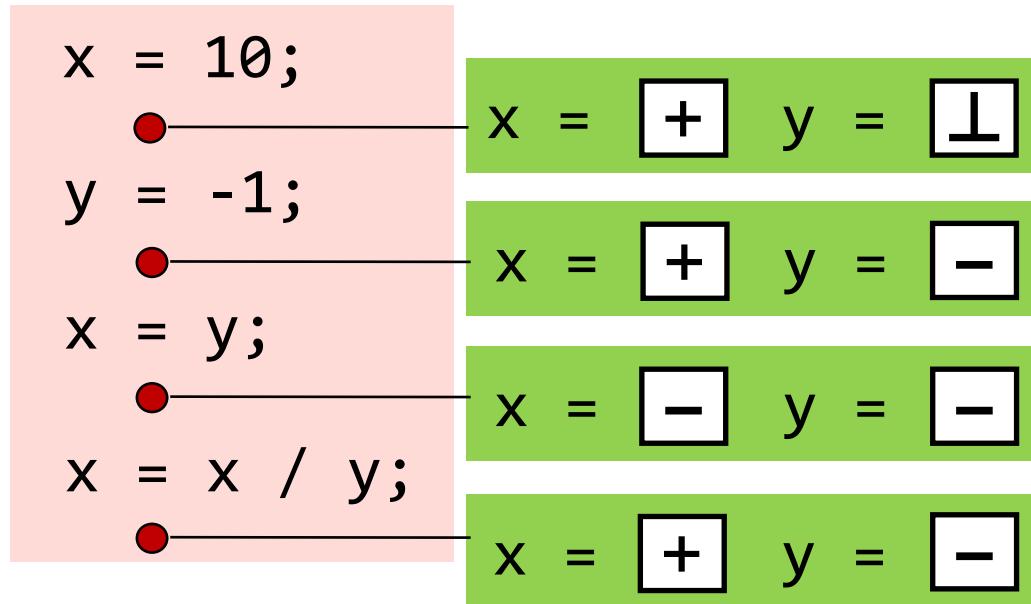
Recall 1<sup>st</sup>  
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Recall 1<sup>st</sup> lecture

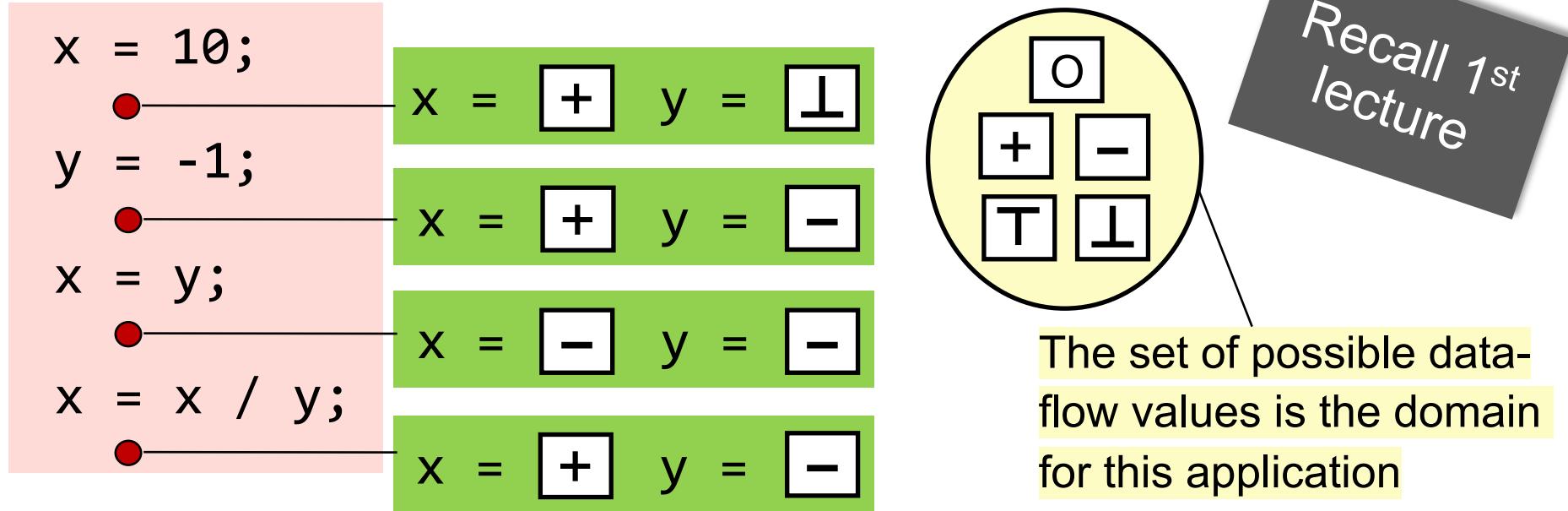
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Recall 1<sup>st</sup> lecture

The set of possible data-flow values is the domain for this application

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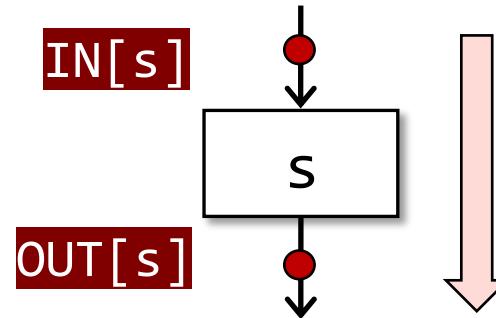
Data-flow analysis is to **find a solution** to a set of *safe-approximation-directed constraints* on the  $\text{IN}[s]$ 's and  $\text{OUT}[s]$ 's, for **all statements**  $s$ .

- *constraints* based on semantics of statements (*transfer functions*)
- *constraints* based on the *flows of control*

# Notations for Transfer Function's Constraints

- Forward Analysis

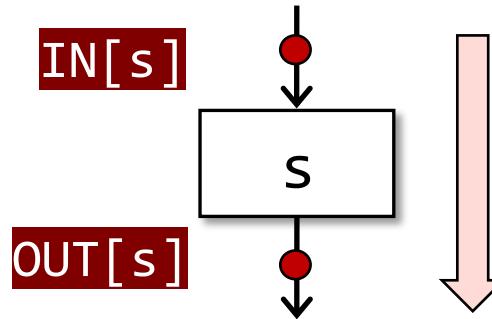
$$\text{OUT}[s] = f_s(\text{IN}[s])$$



# Notations for Transfer Function's Constraints

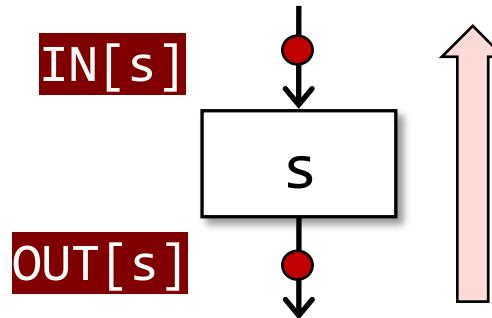
- Forward Analysis

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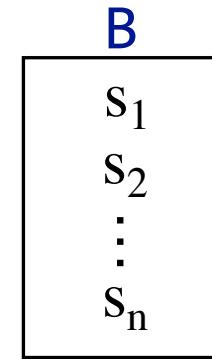
- Backward Analysis

$$\text{IN}[s] = f_s(\text{OUT}[s])$$



# Notations for Control Flow's Constraints

- Control flow within a BB
- Control flow among BBs

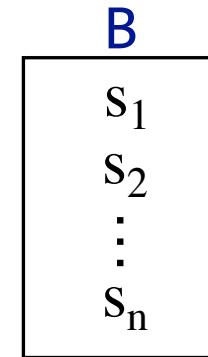


# Notations for Control Flow's Constraints

- Control flow within a BB

$$\text{IN}[s_{i+1}] = \text{OUT}[s_i], \text{ for all } i = 1, 2, \dots, n-1$$

- Control flow among BBs



# Notations for Control Flow's Constraints

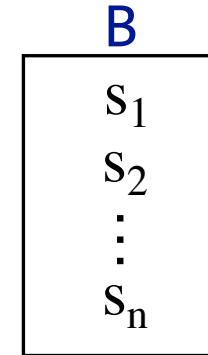
- Control flow within a BB

$$\text{IN}[s_{i+1}] = \text{OUT}[s_i], \text{ for all } i = 1, 2, \dots, n-1$$

- Control flow among BBs

$$\text{IN}[B] = \text{IN}[s_1]$$

$$\text{OUT}[B] = \text{OUT}[s_n]$$



# Notations for Control Flow's Constraints

- Control flow within a BB

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- Control flow among BBs

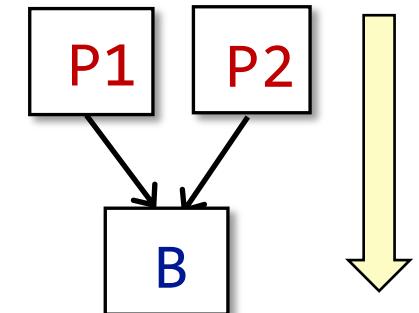
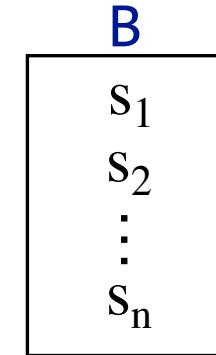
$$\text{IN}[B] = \text{IN}[s_1]$$

$$\text{OUT}[B] = \text{OUT}[s_n]$$

$$\text{OUT}[B] = f_B(\text{IN}[B]), \quad f_B = f_{s_n} \circ \dots \circ f_{s_2} \circ f_{s_1}$$

$$\text{IN}[B] = \bigwedge P \text{ a predecessor of } B \text{ OUT}[P]$$

The meet operator  $\bigwedge$  is used to summarize the contributions from different paths at the confluence of those paths



# Notations for Control Flow's Constraints

- Control flow within a BB

$$\text{IN}[s_{i+1}] = \text{OUT}[s_i], \text{ for all } i = 1, 2, \dots, n-1$$

- Control flow among BBs

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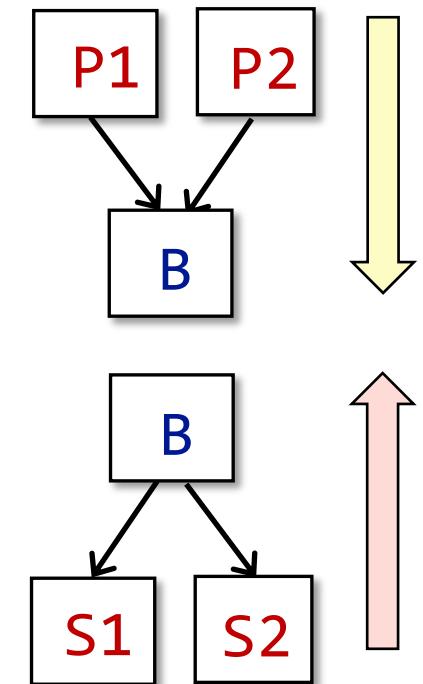
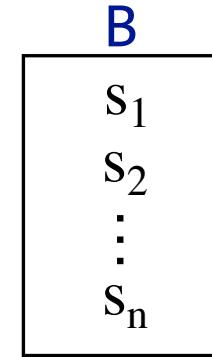
$$\text{OUT}[B] = \text{OUT}[s_n]$$

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$$\text{OUT}[B] = \bigwedge_S S \text{ a successor of } B \text{ } \text{IN}[S]$$



# Data Flow Analysis Applications

(I) Reaching Definitions Analysis

(II) Live Variables Analysis

(III) Available Expressions Analysis

# Issues Not Covered

- Method Calls
  - Intra-procedural CFG
  - Will be introduced in lecture: Inter-procedural Analysis
- Aliases
  - Variables have no aliases
  - Will be introduced in lecture: Pointer Analysis

# Reaching Definitions

A **definition d** at program point p *reaches* a point q if there is a path from p to q such that **d** is not “killed” along that path

# Reaching Definitions

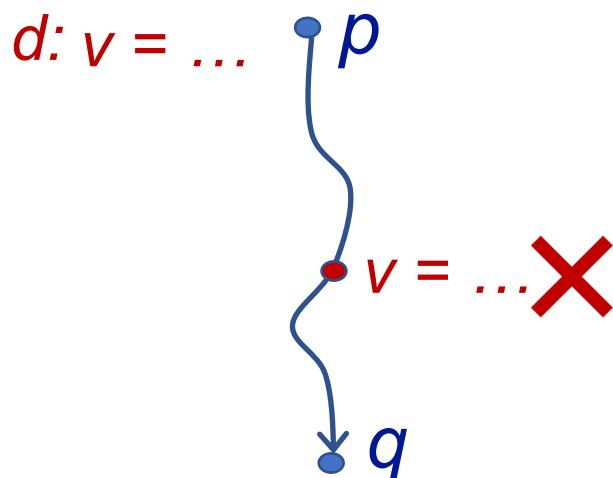
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- A **definition of a variable v** is a statement that assigns a value to v

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- Translated as: definition of variable v at program point p reaches point q if there is a path from p to q such that no new definition of v appears on that path



# Reaching Definitions

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- Translated as: definition of variable v at program point p reaches point q if there is a path from p to q such that no new definition of v appears on that path
- Reaching definitions can be used to **detect possible undefined variables**. e.g., introduce a dummy definition for each variable v at the entry of CFG, and if the dummy definition of v reaches a point p where v is used, then v may be used before definition (*as undefined reaches v*)

# Understanding Reaching Definitions

- Data Flow Values/Facts
  - The definitions of all the variables in a program

*Abstraction*

# Understanding Reaching Definitions

- Data Flow Values/Facts
    - The definitions of all the variables in a program
    - Can be represented by bit vectors
- e.g., D1, D2, D3, D4, ..., D100 (100 definitions)

A diagram showing a sequence of bits: 00000...0. A bracket underneath the first 100 bits is labeled "100 bits".

Bit i from the left represents definition Di

*Abstraction*

# Understanding Reaching Definitions

*Safe-approximation*

- Transfer Function
- Control Flow

# Understanding Reaching Definitions

D:  $v = x \ op \ y$

*Safe-approximation*

This statement “**generates**” a definition D of variable v and “**kills**” all the other definitions in the program that define variable v, while leaving the remaining incoming definitions unaffected.

- Transfer Function
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# Understanding Reaching Definitions

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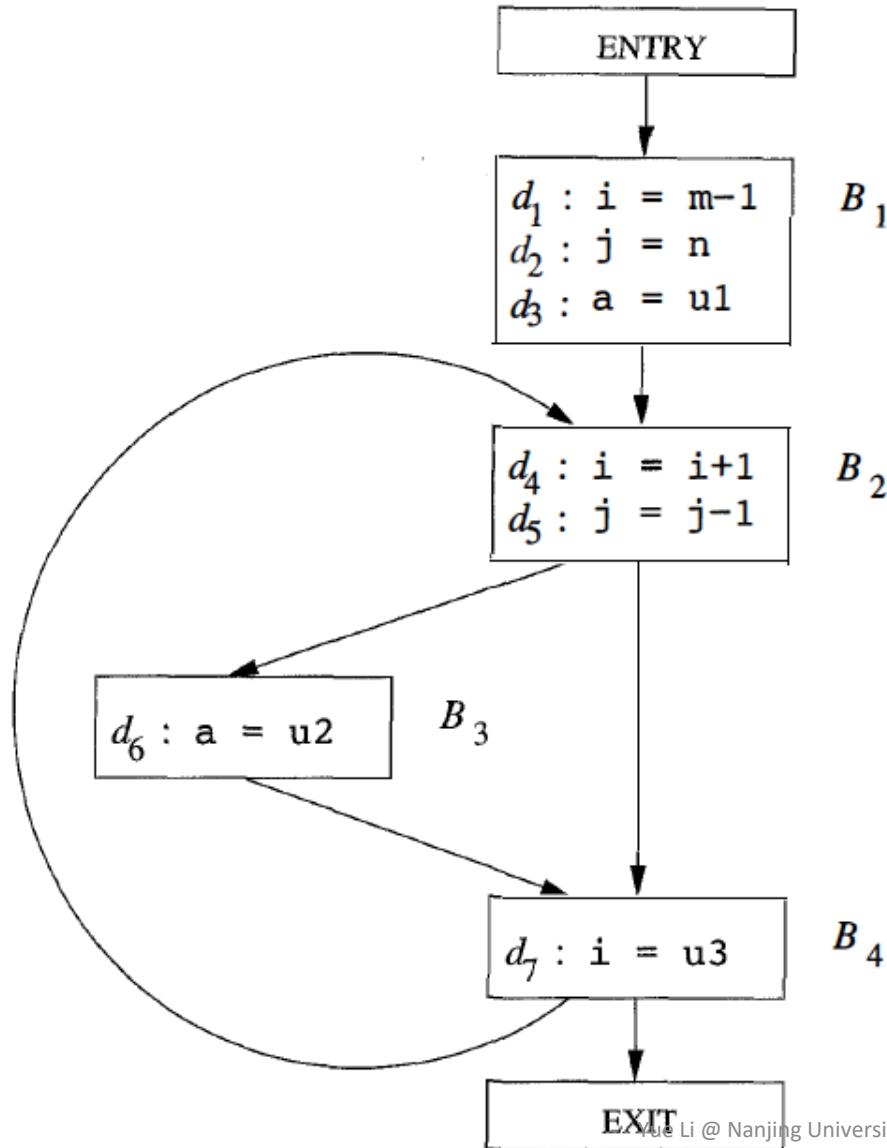
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- Transfer Function

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$

- Control Flow

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$



$$\text{gen}_{B_1} = \{ d_1, d_2, d_3 \}$$

$$\text{kill}_{B_1} = \{ d_4, d_5, d_6, d_7 \}$$

$$\text{gen}_{B_2} = \{ d_4, d_5 \}$$

$$\text{kill}_{B_2} = \{ d_1, d_2, d_7 \}$$

$$\text{gen}_{B_3} = \{ d_6 \}$$

$$\text{kill}_{B_3} = \{ d_3 \}$$

$$\text{gen}_{B_4} = \{ d_7 \}$$

$$\text{kill}_{B_4} = \{ d_1, d_4 \}$$

# Understanding Reaching Definitions

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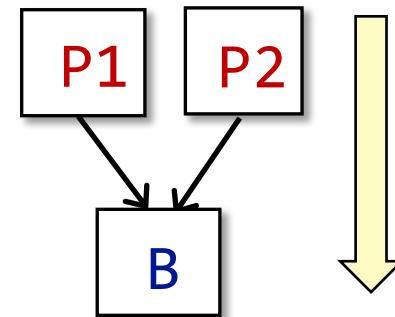
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$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

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$$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P]$$



# Understanding Reaching Definitions

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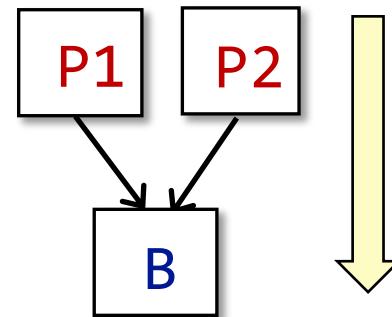
Safe-approximation

This statement “**generates**” a definition D of variable v and “**kills**” all the other definitions in the program that define variable v, while leaving the remaining incoming definitions unaffected.

A definition reaches a program point as long as there exists at least one path along which the definition reaches.

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# Understanding Reaching Definitions

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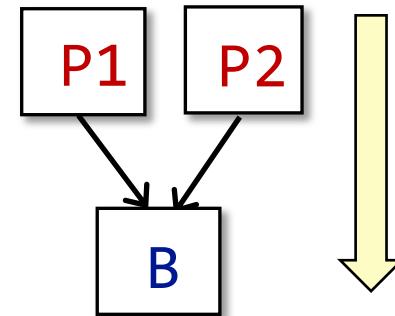
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# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[B] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[B] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P];$   
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```

*Why entry is excluded?*

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**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

$OUT[entry] = \emptyset;$

**for** (each basic block  $B \setminus entry$ )

$OUT[B] = \emptyset;$

**while** (changes to any  $OUT$  occur) ?

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$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P];$

$OUT[B] = gen_B \cup (IN[B] - kill_B);$

}

*Why this iterative algorithm  
can finally stop?*

# Algorithm of Reaching Definitions Analysis

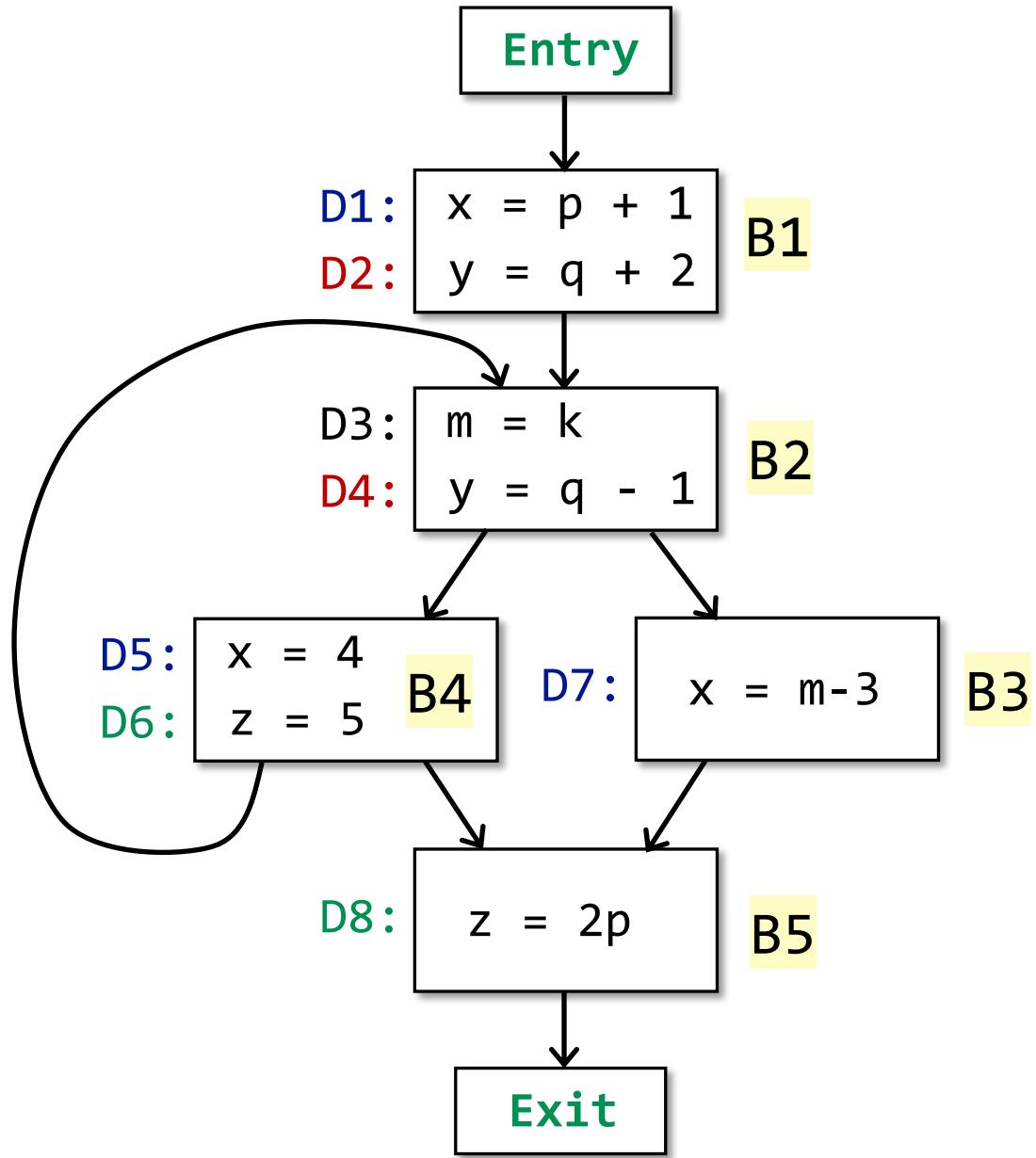
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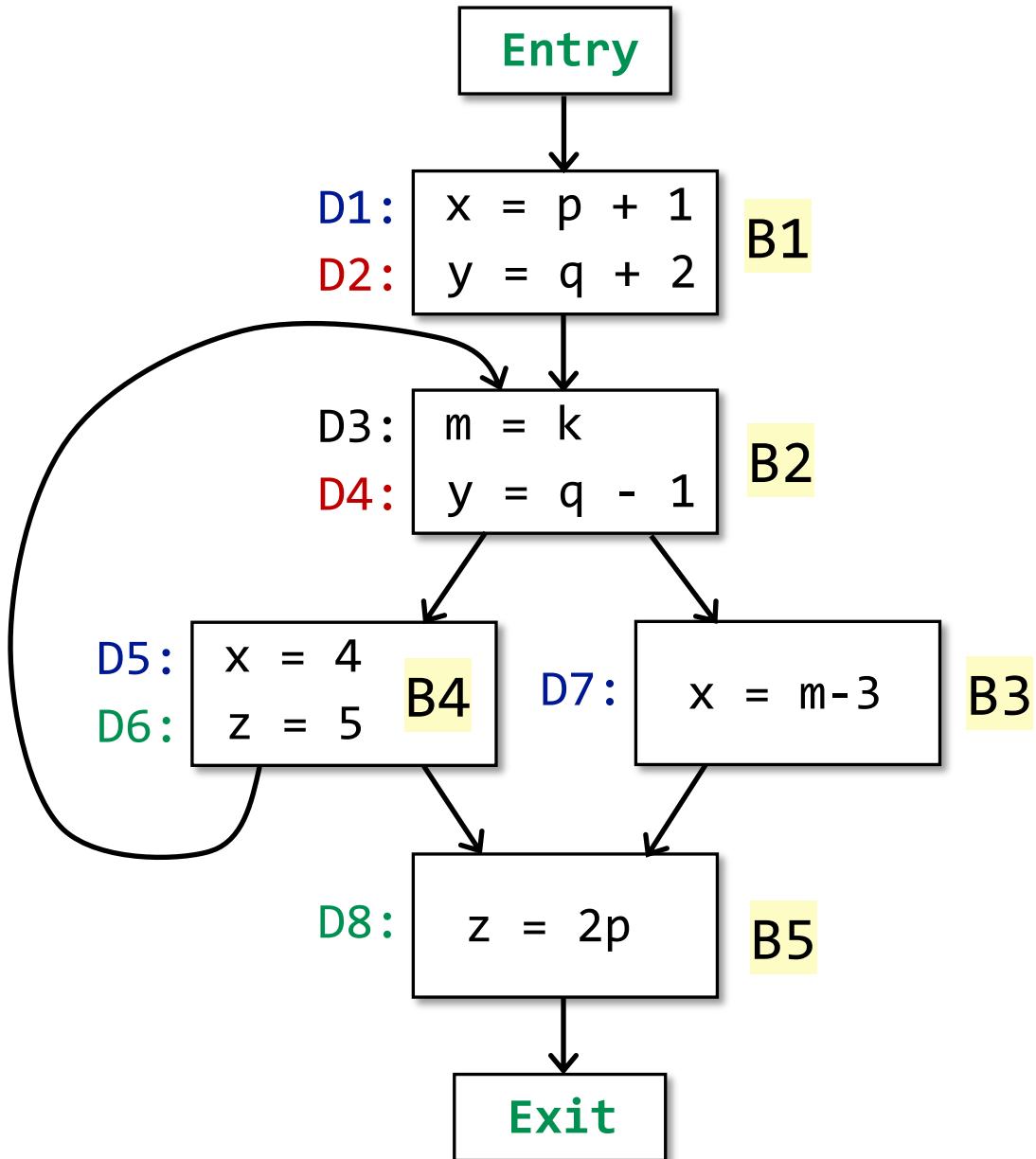


```

D1
D2
do {
    D3
    D4
    if(...) {
        D5
        D6
    } else {
        D7
        break
    }
} while(...)

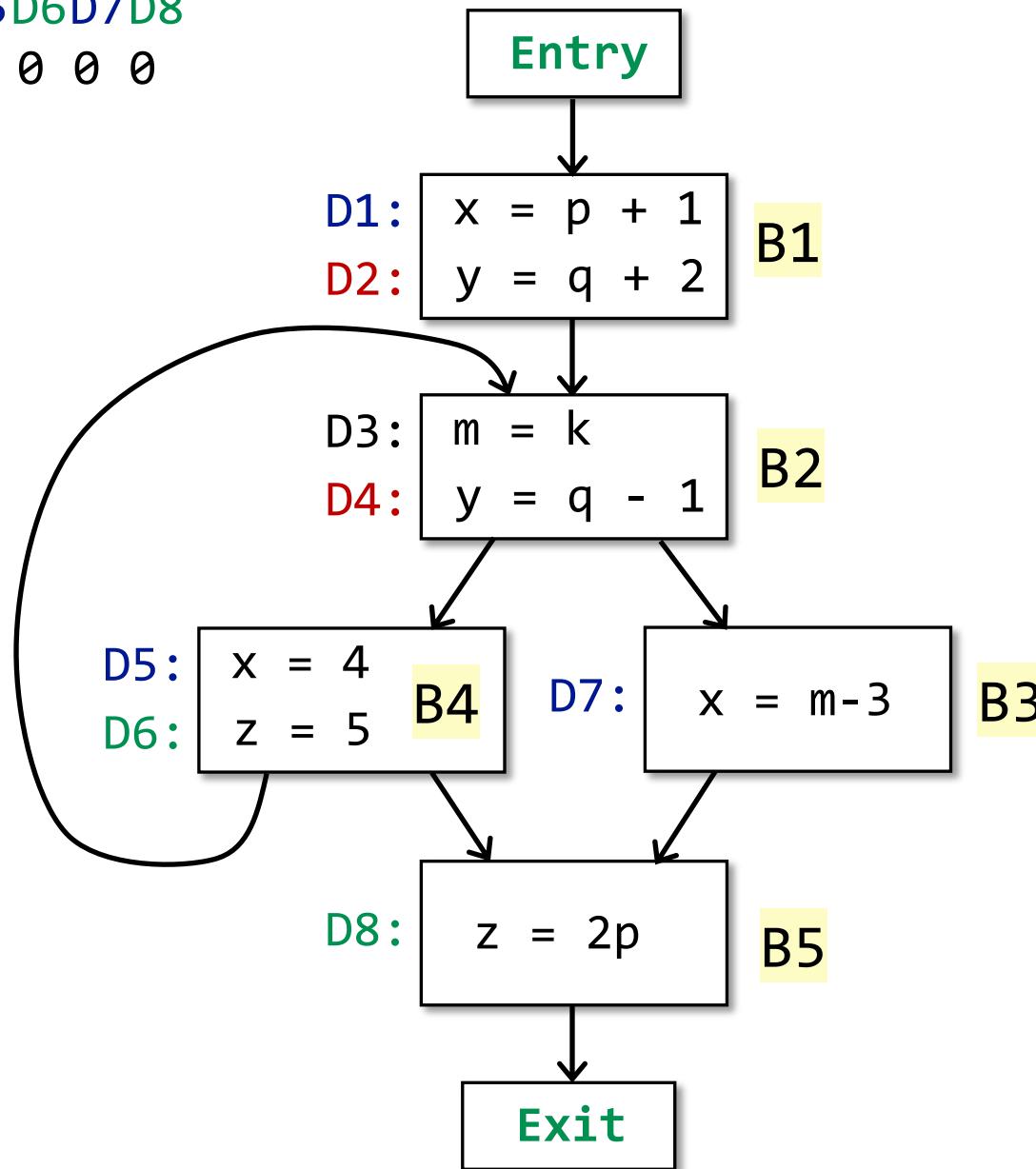
D8

```



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0



# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

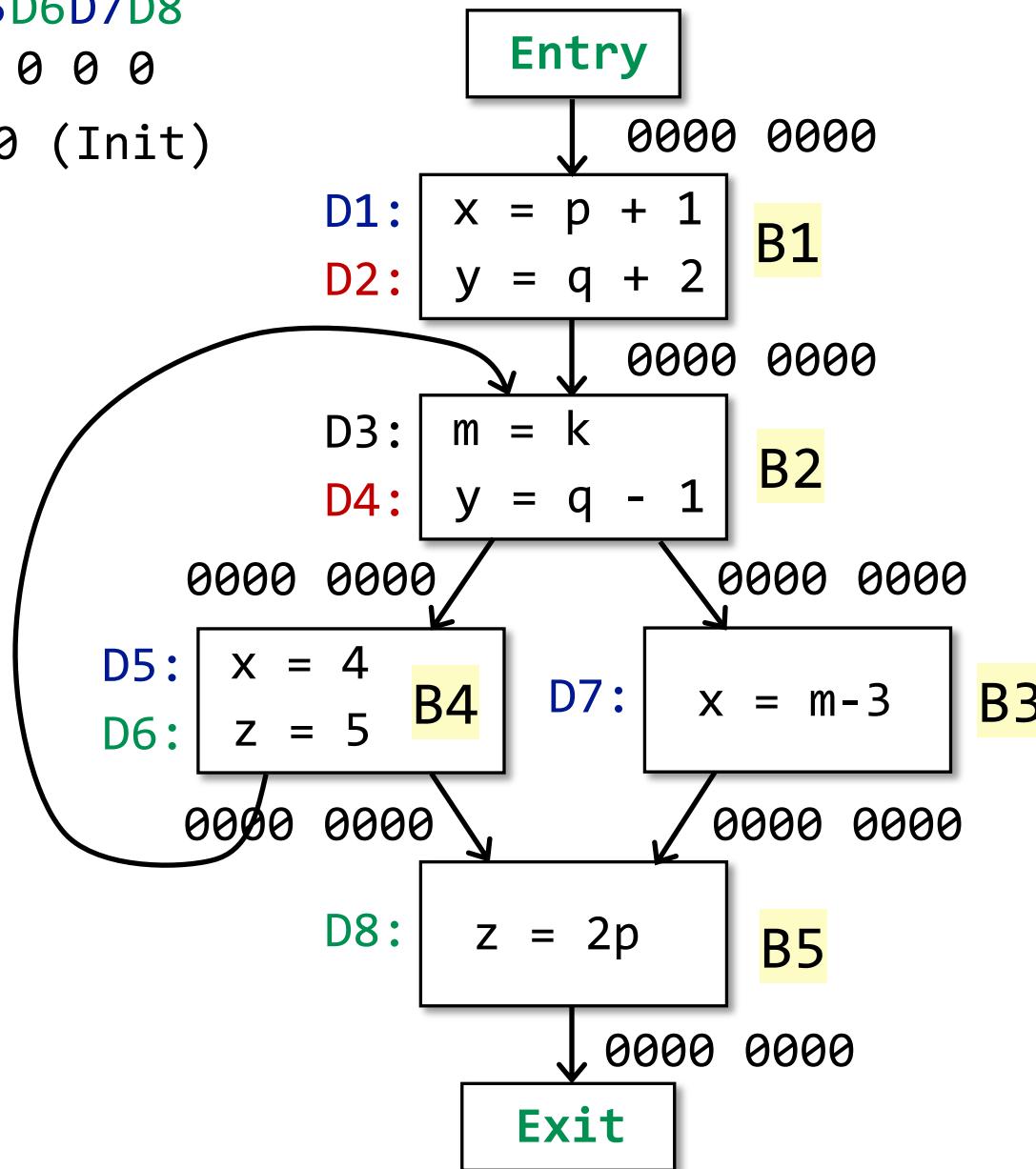
```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[B] = Ø;  
  
while (changes to any OUT occur)  
    for (each basic block  $B \setminus entry$ ) {  
        IN[B] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P];$   
        OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
    }
```



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)



# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[B] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[B] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P];$   
            OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
        }
```

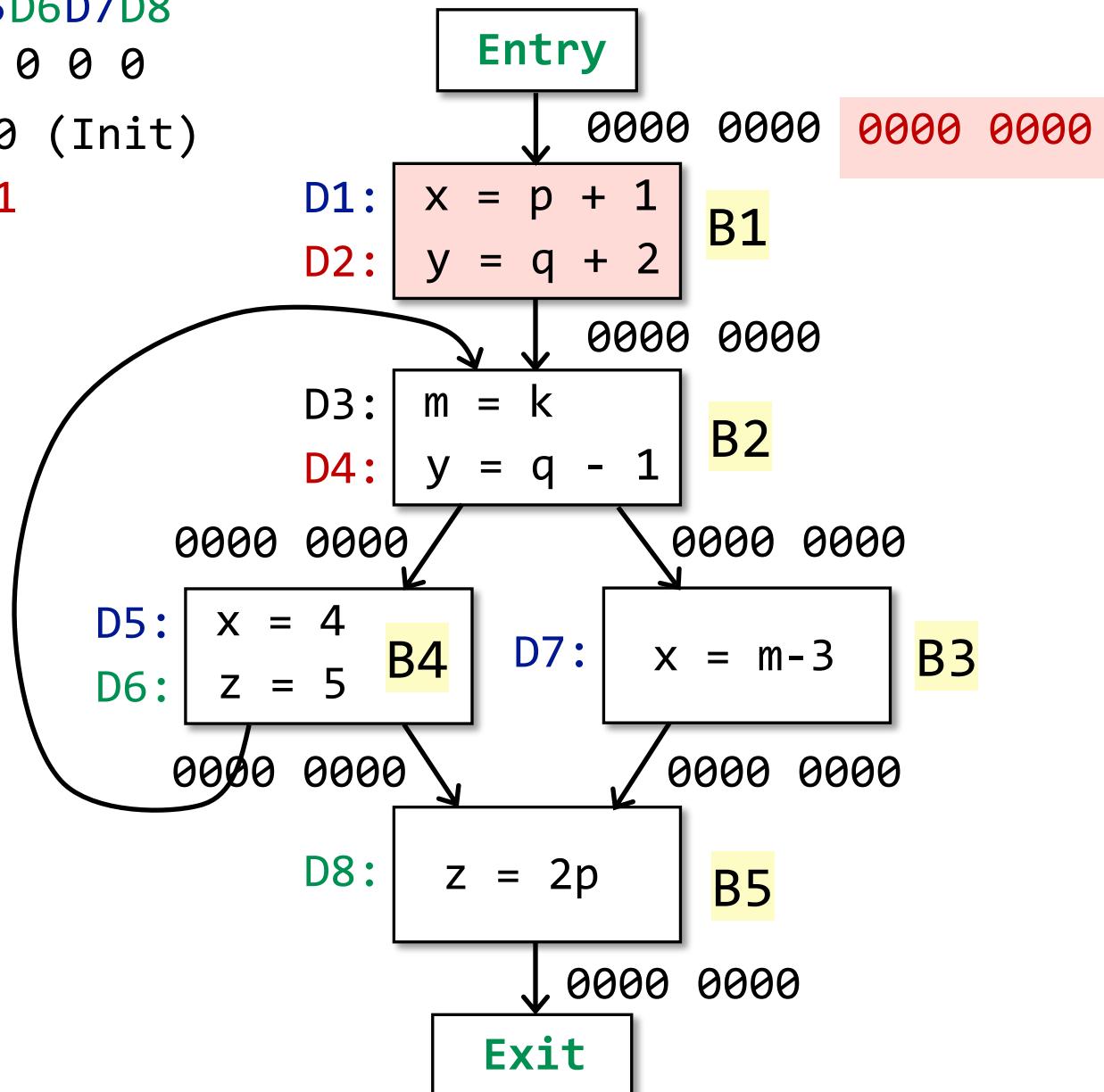


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1



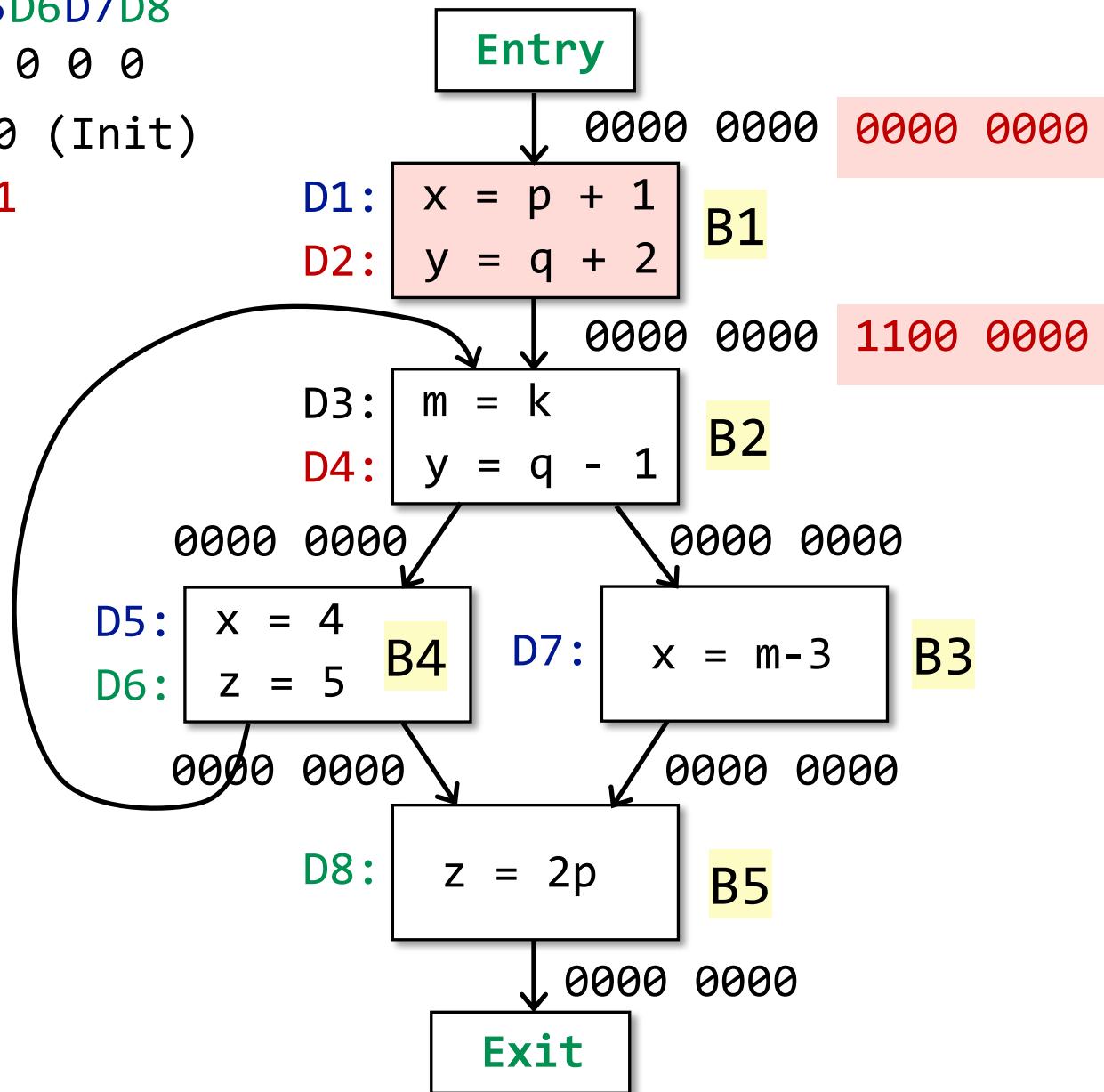
$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B);$  @ Nanjing University

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

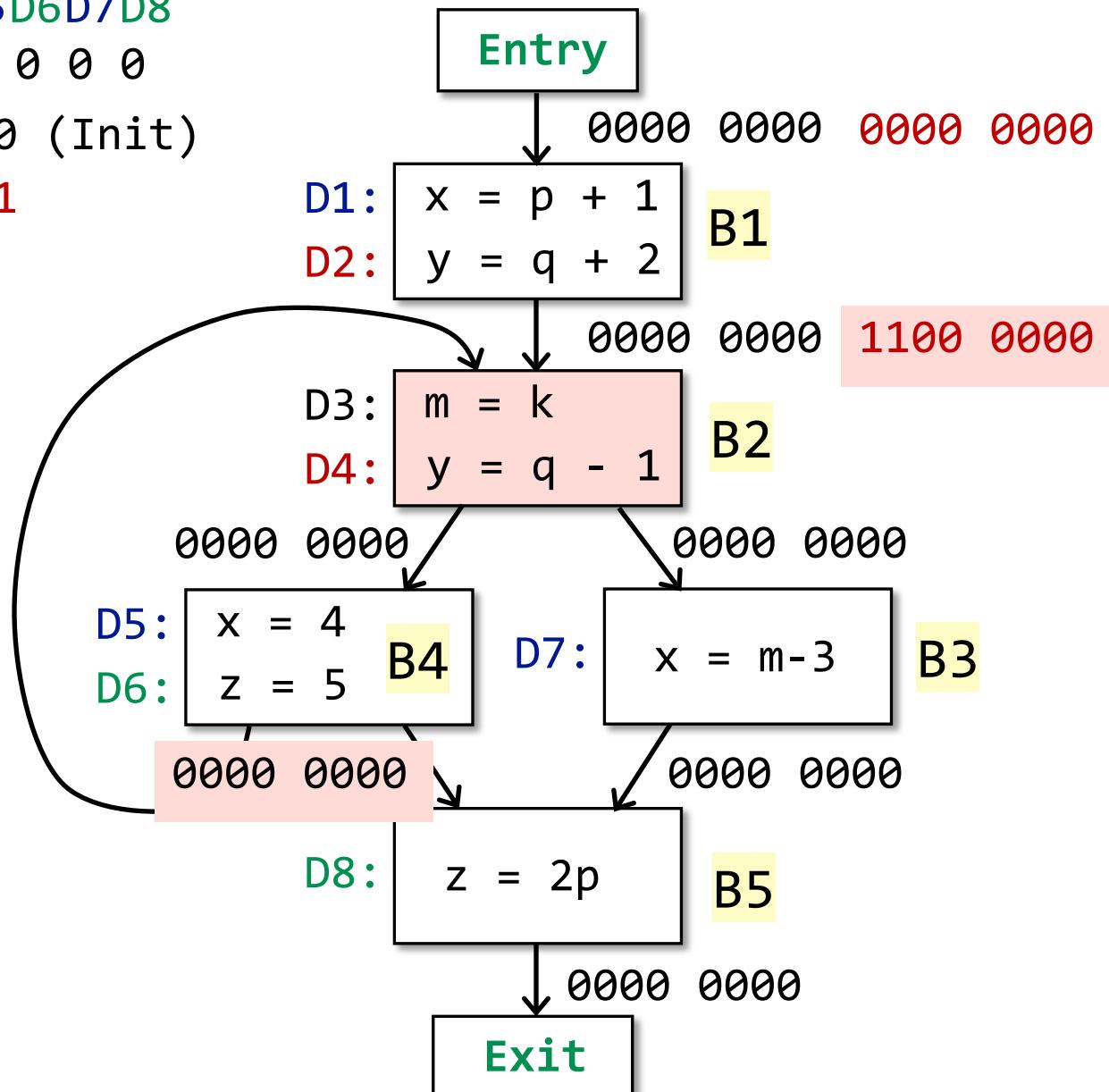


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1



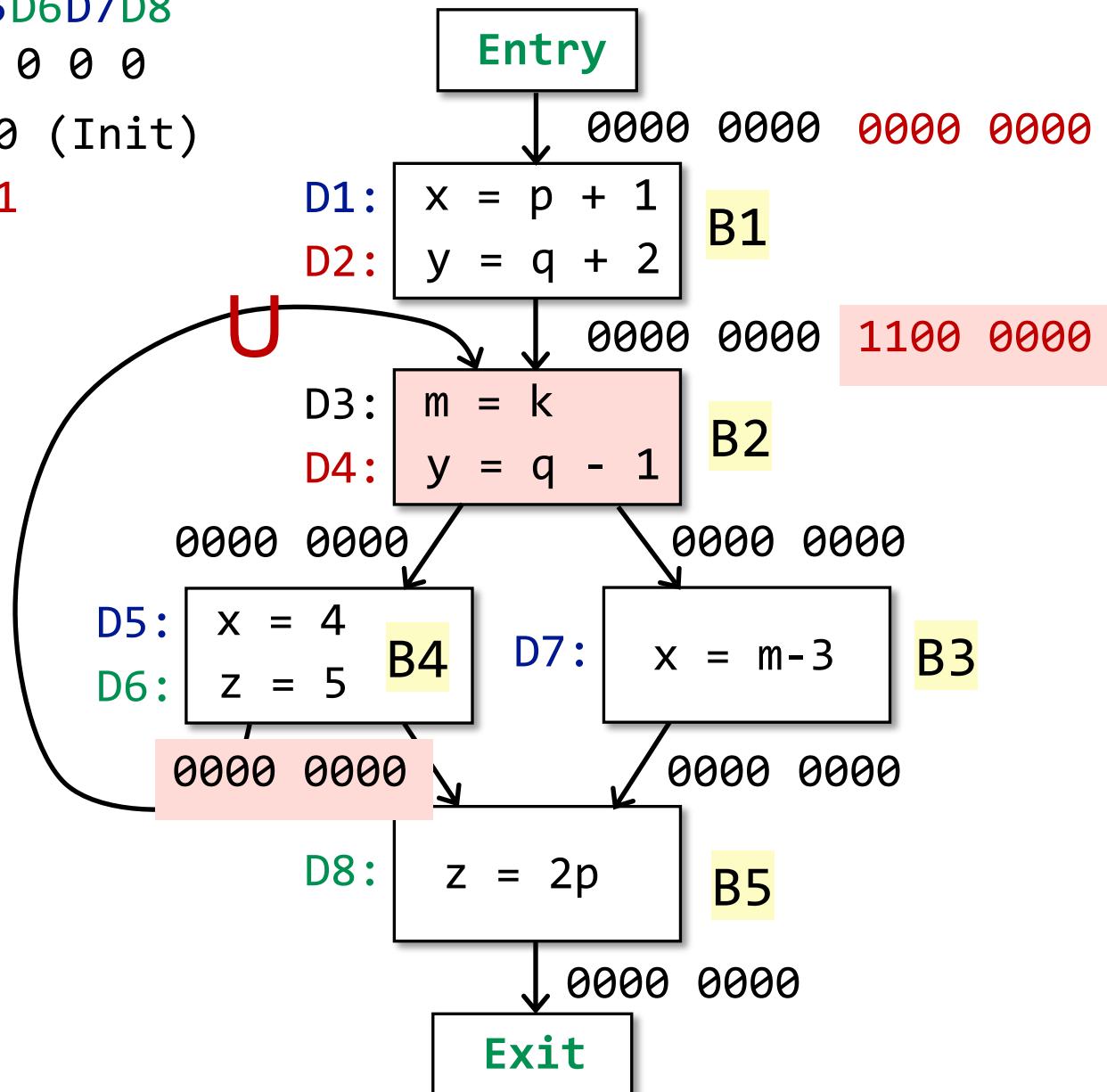
$$\text{IN}[B] = \bigcup_{P \text{ a predecessor of } B} \text{OUT}[P];$$

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

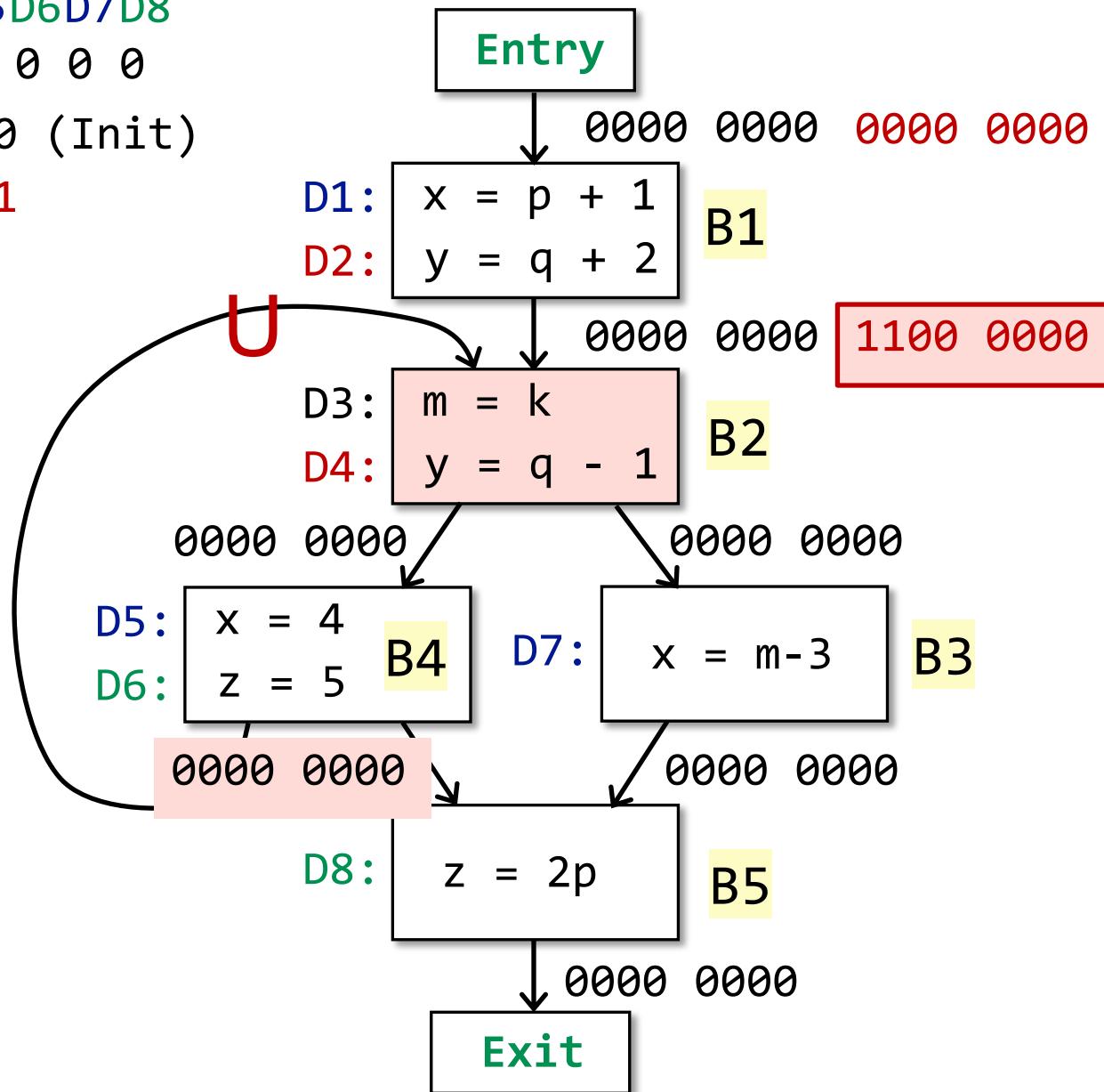


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

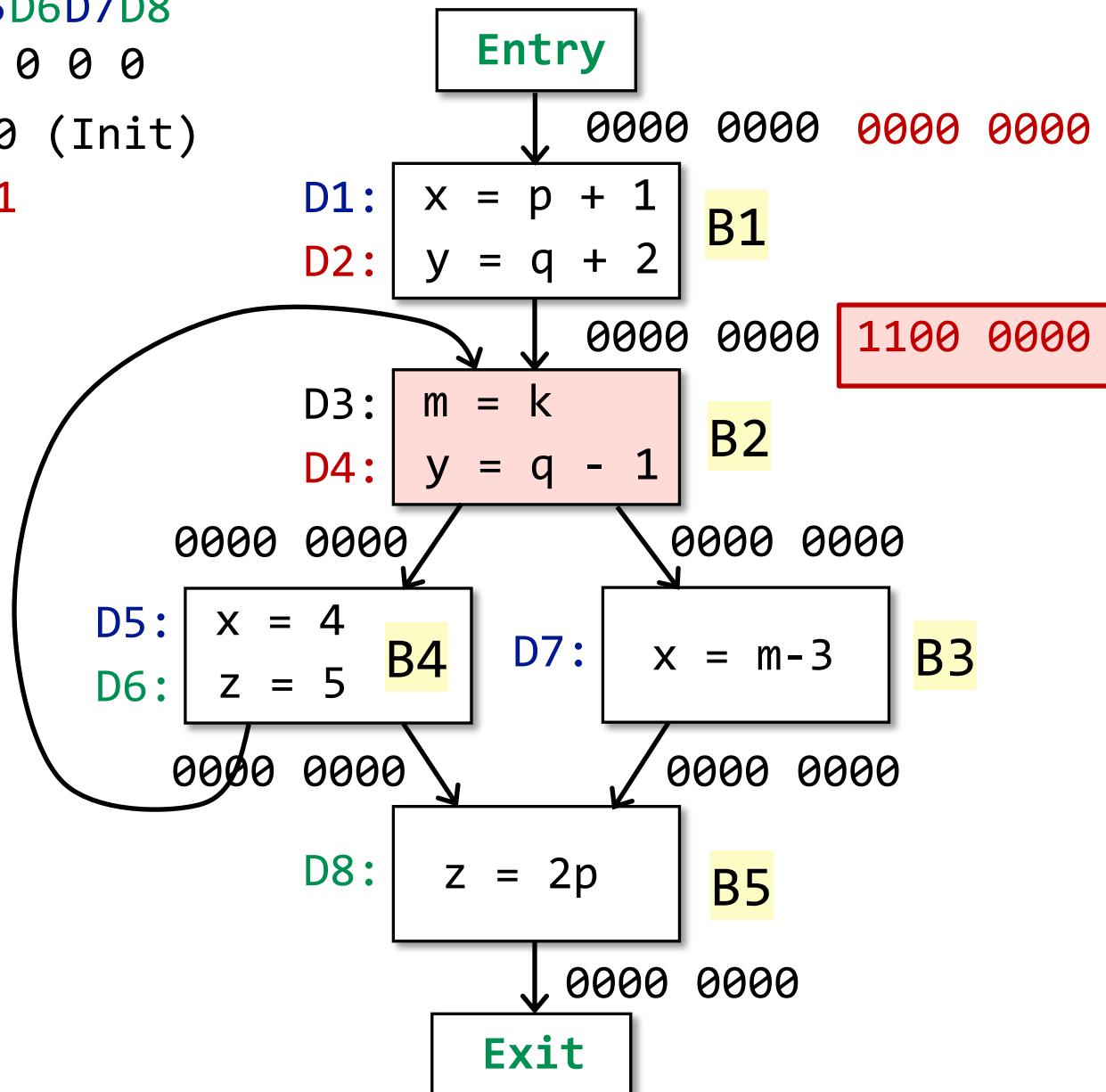


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

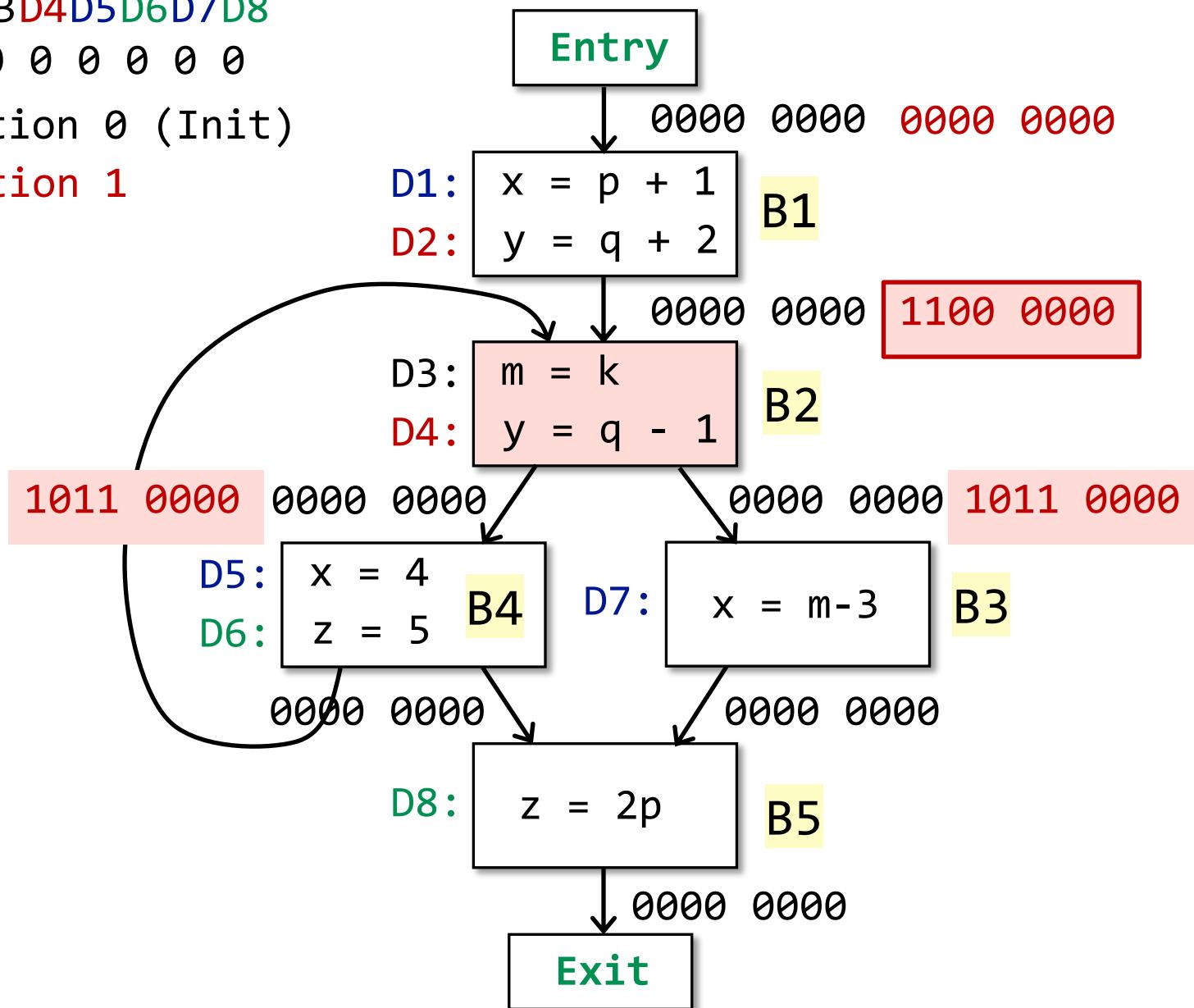


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

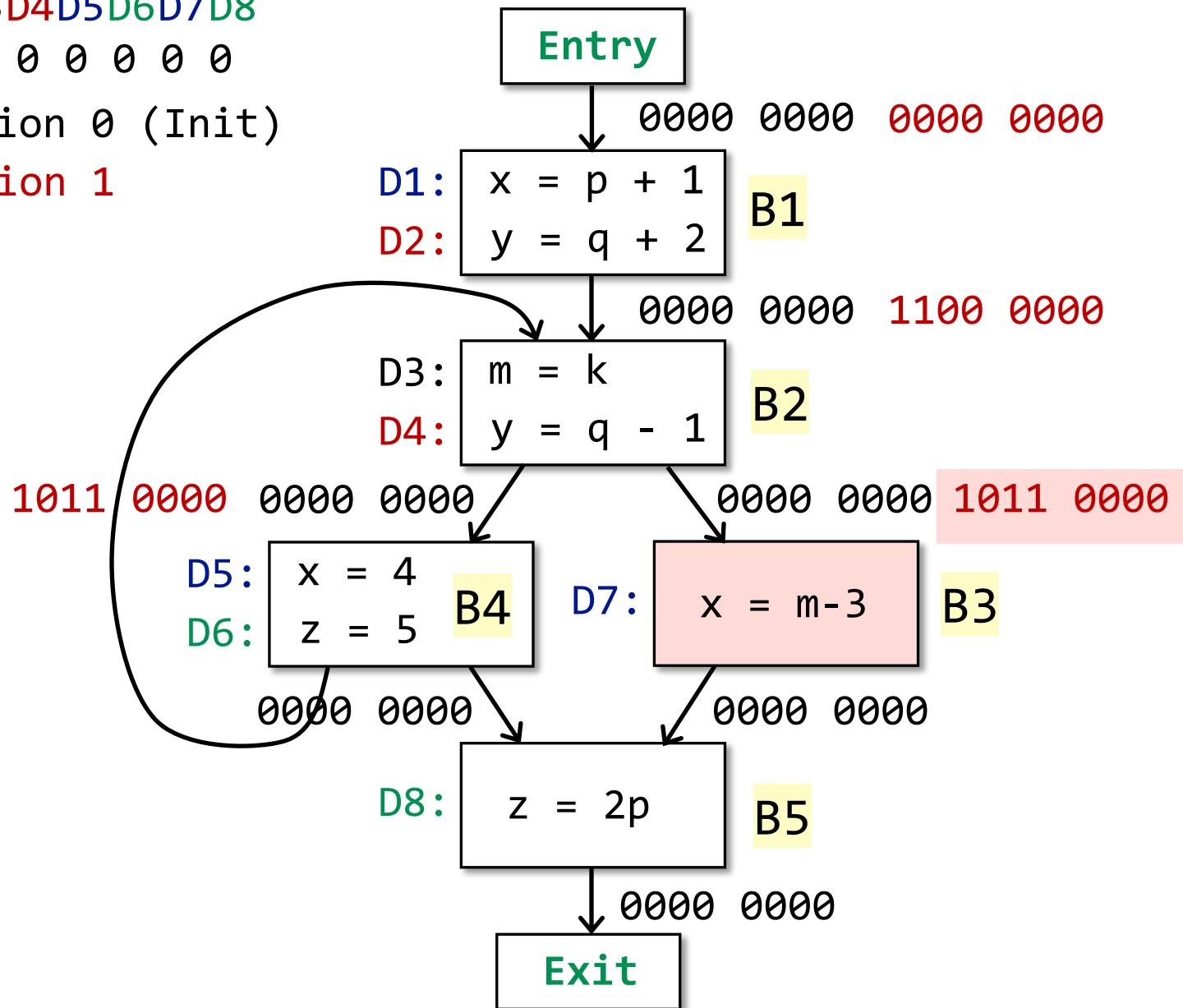


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

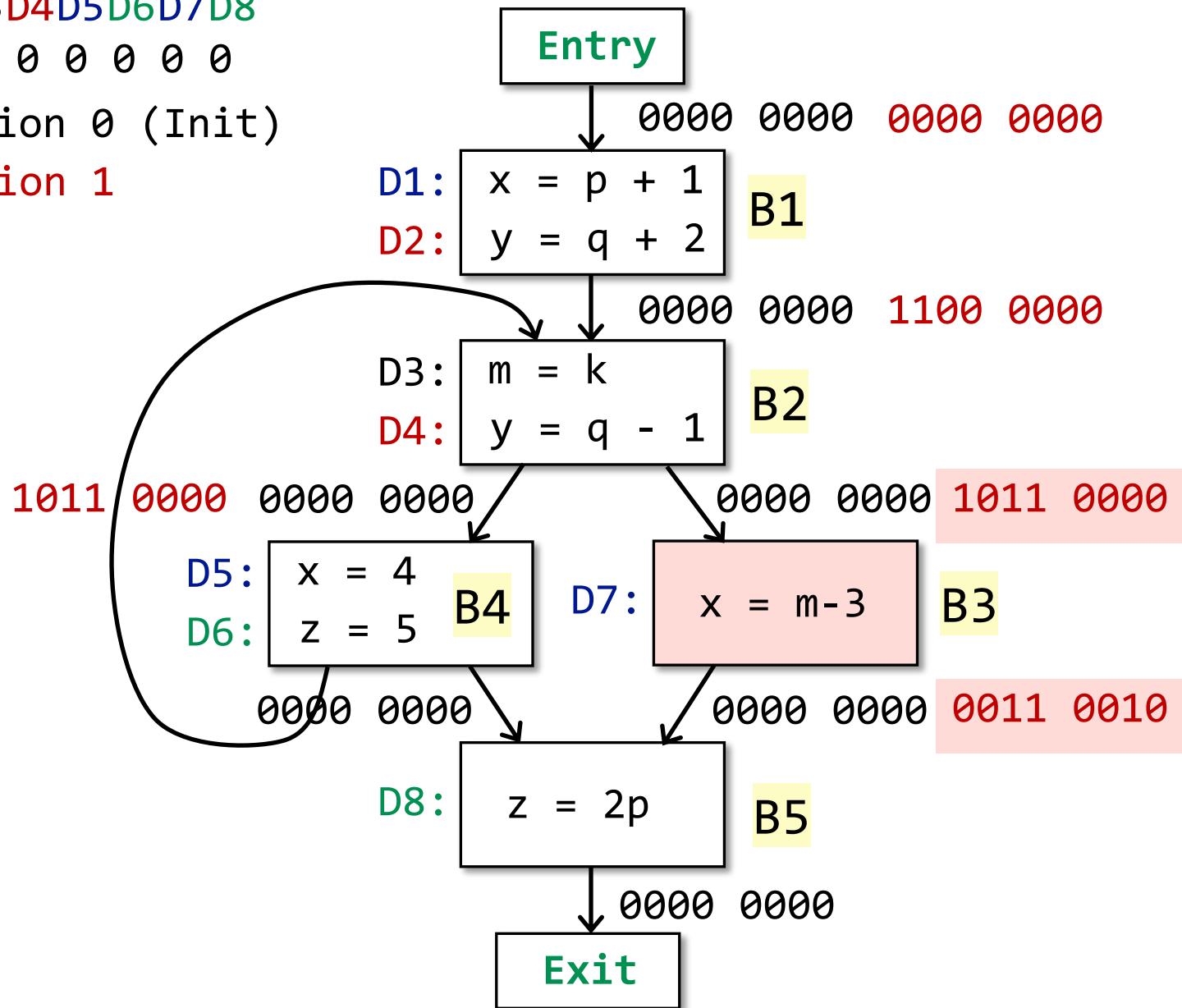


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

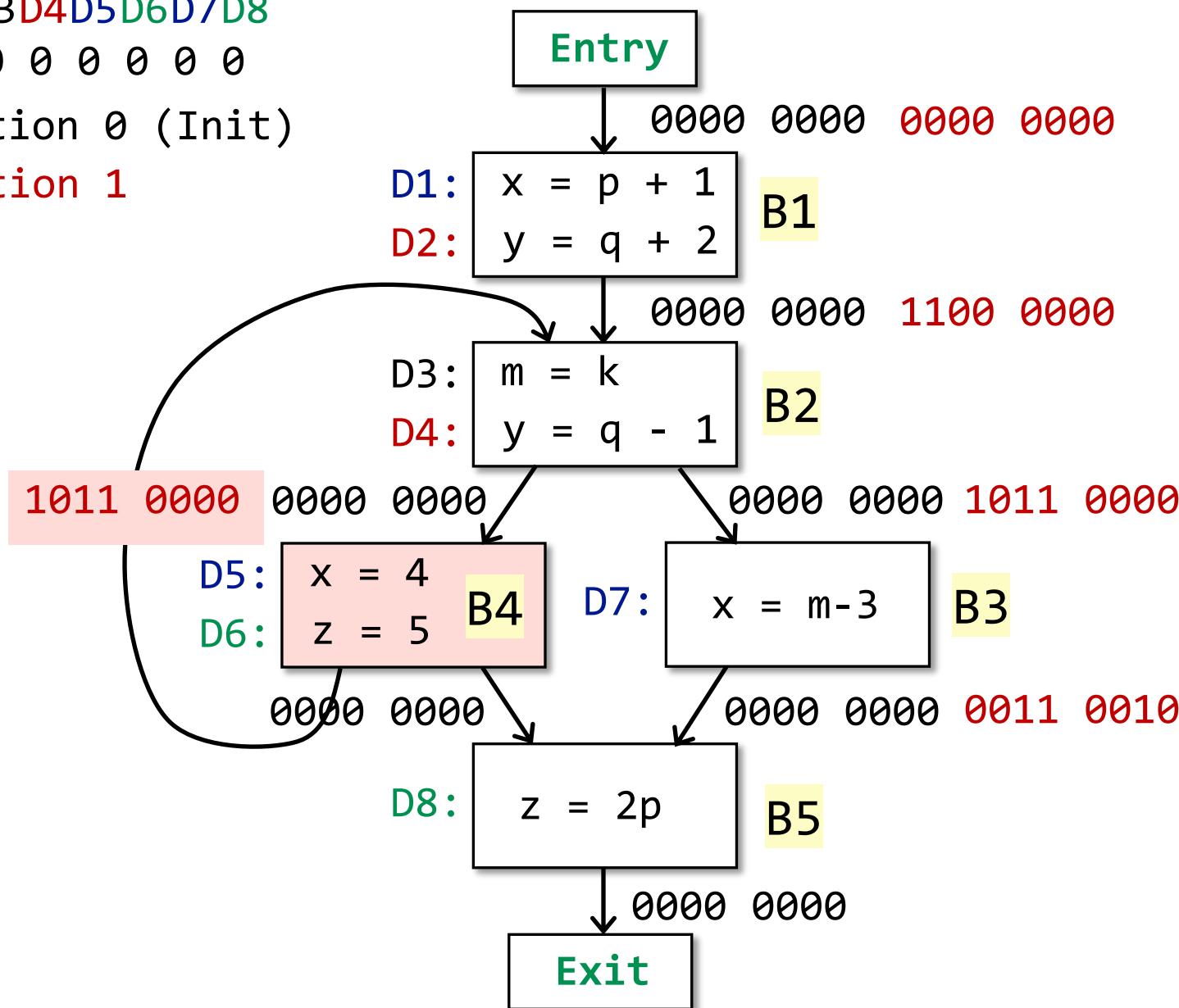


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

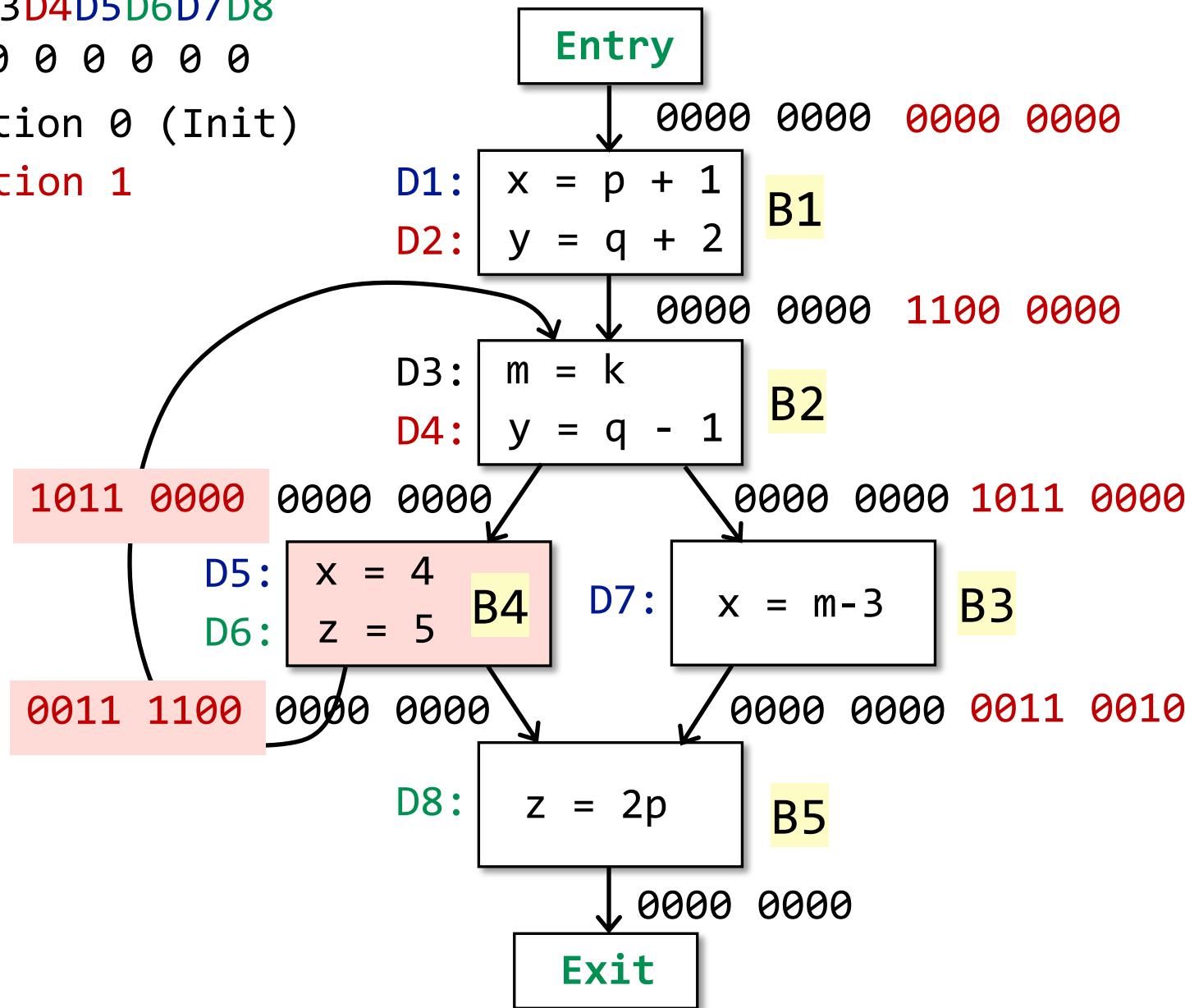


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

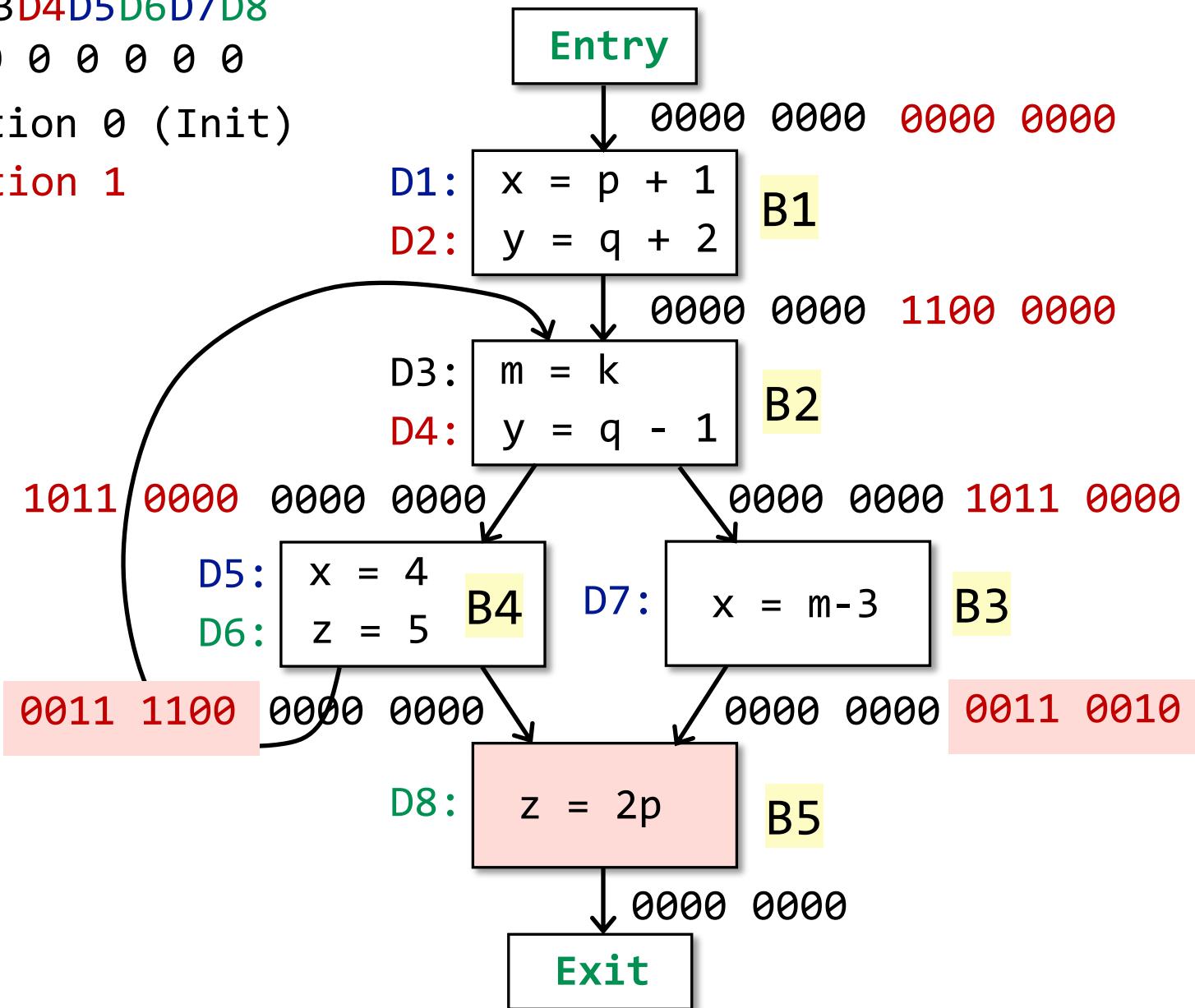


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

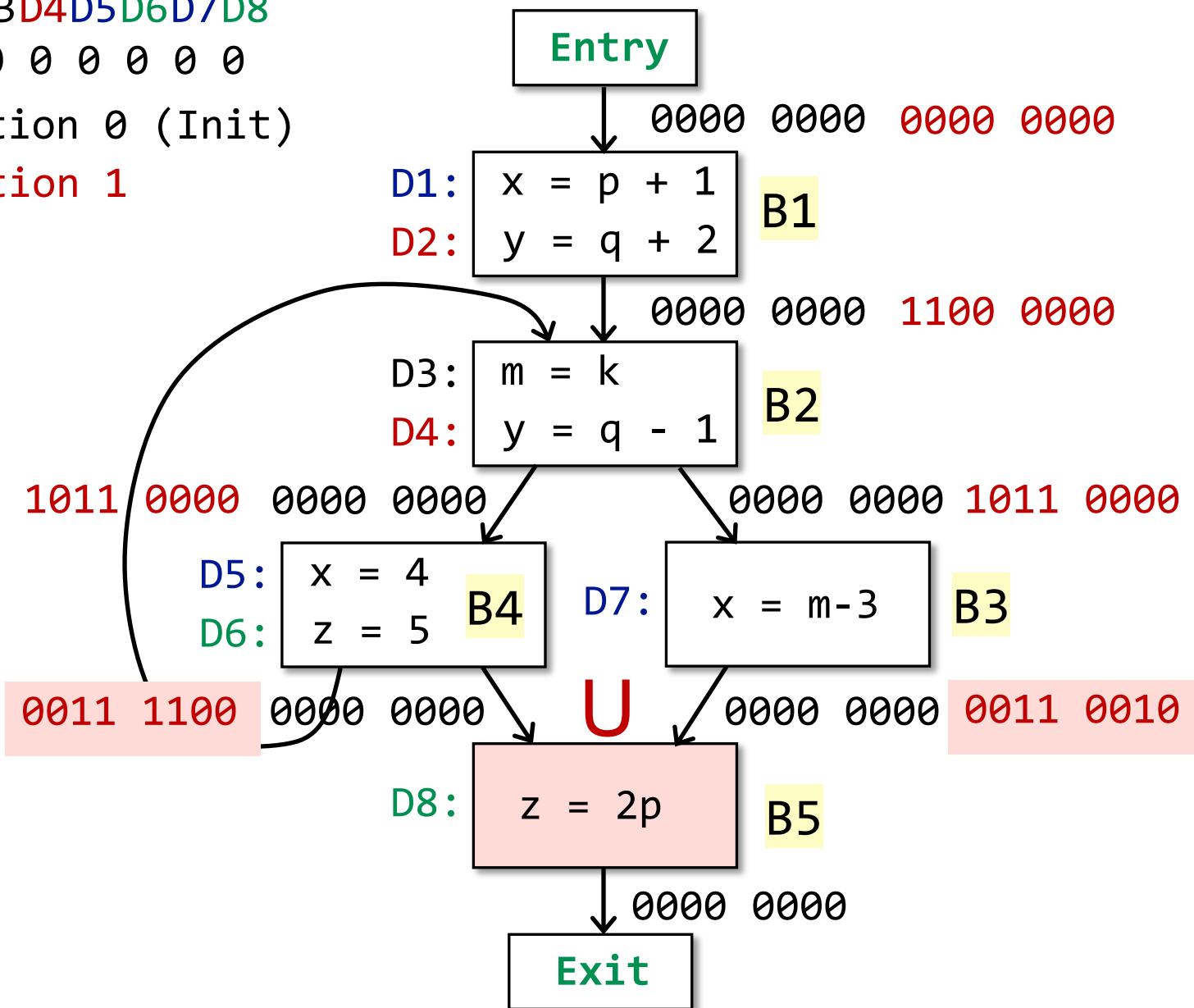


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

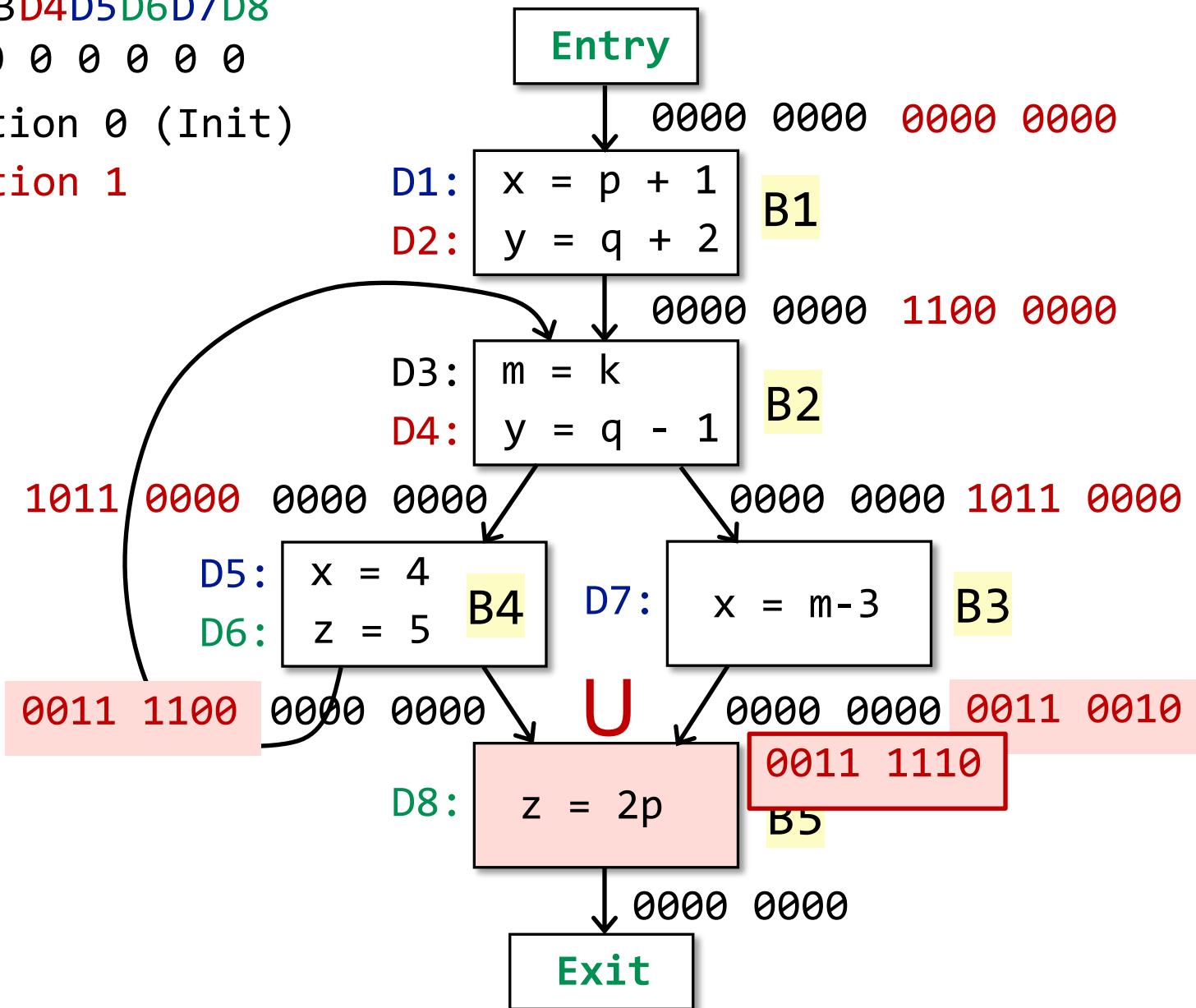


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

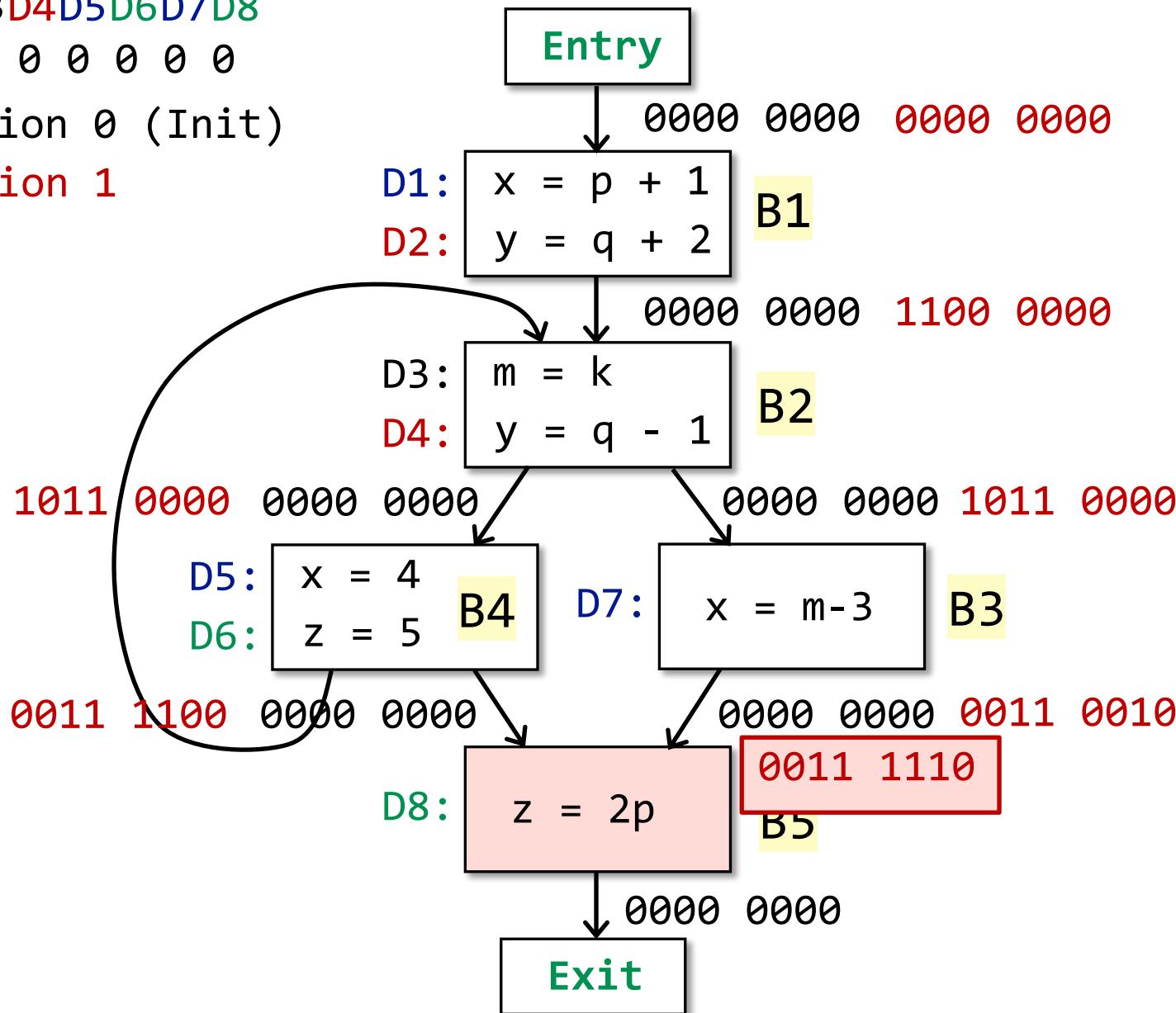


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

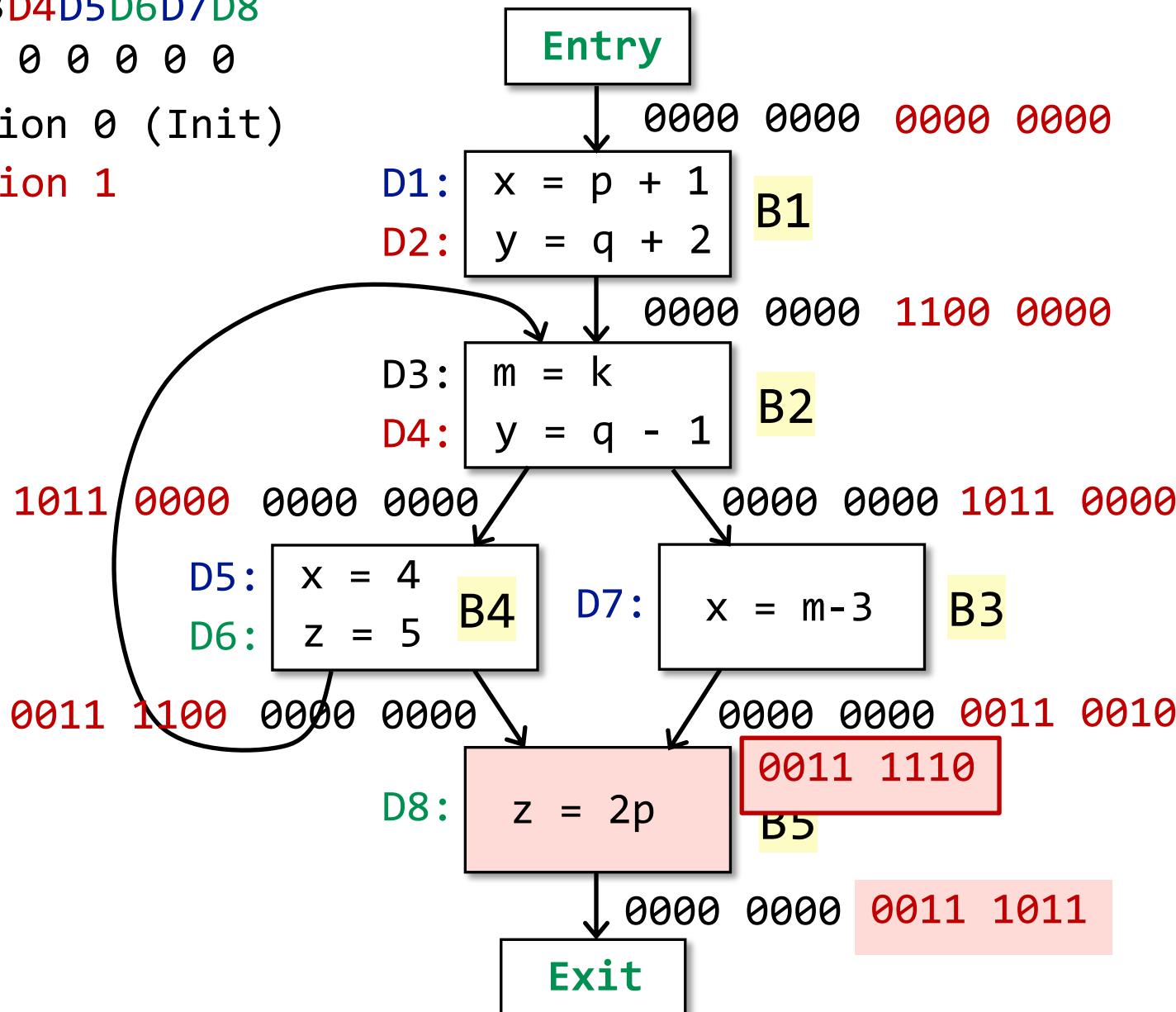


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

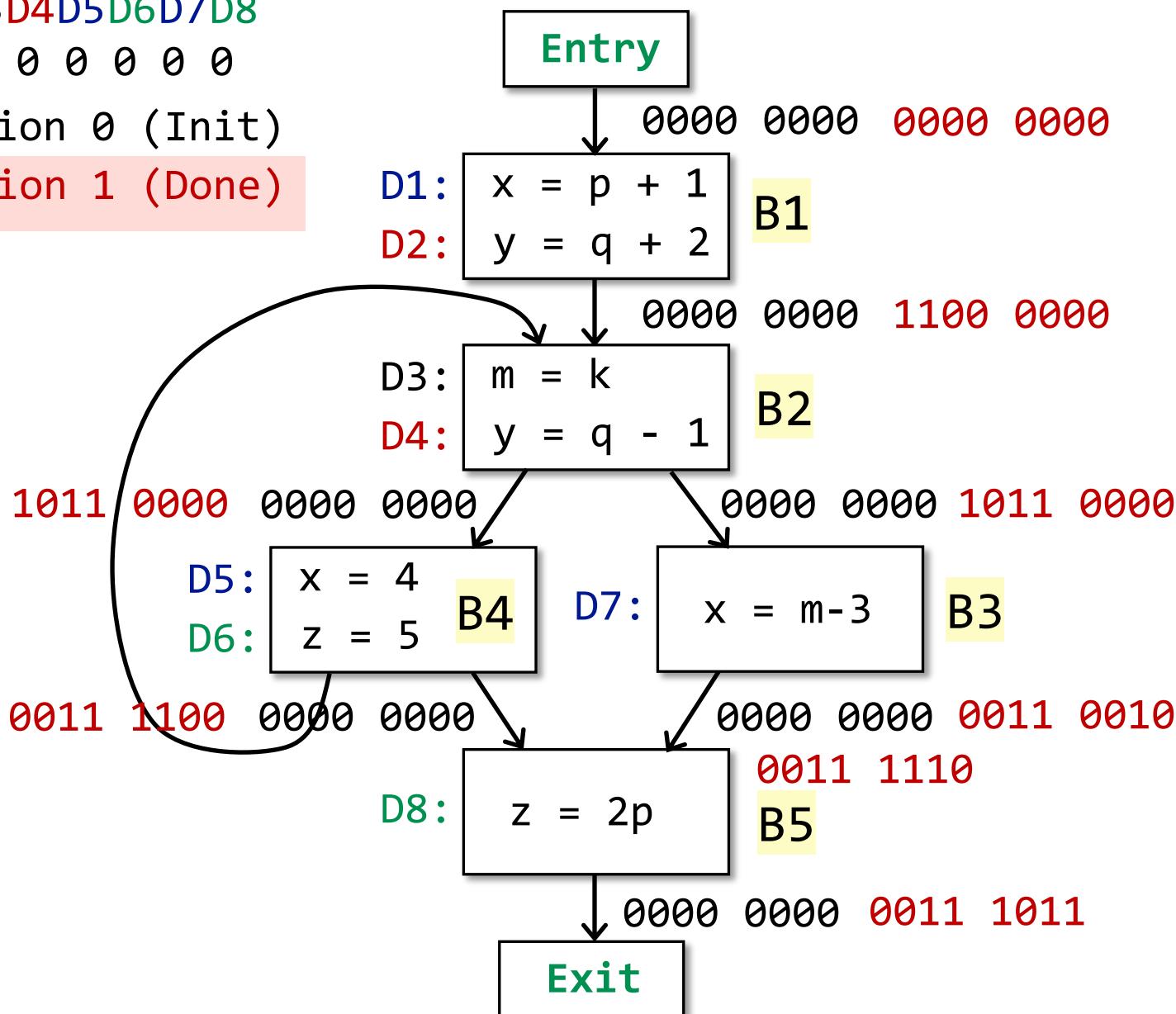


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)



# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

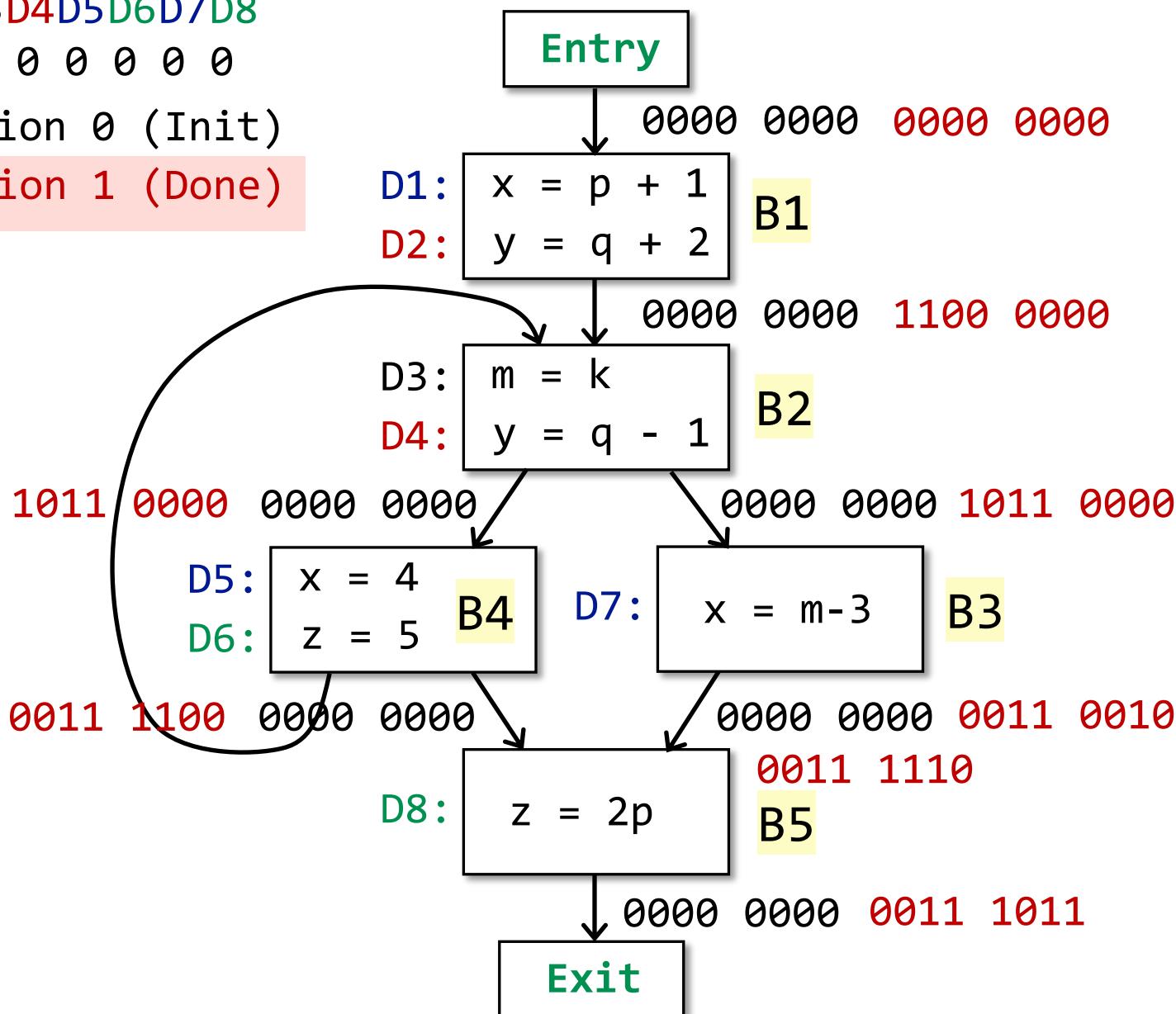
```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[ $B$ ] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[ $B$ ] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P]$ ;  
            OUT[ $B$ ] =  $gen_B \cup (IN[B] - kill_B)$ ;  
        }  
    }
```

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

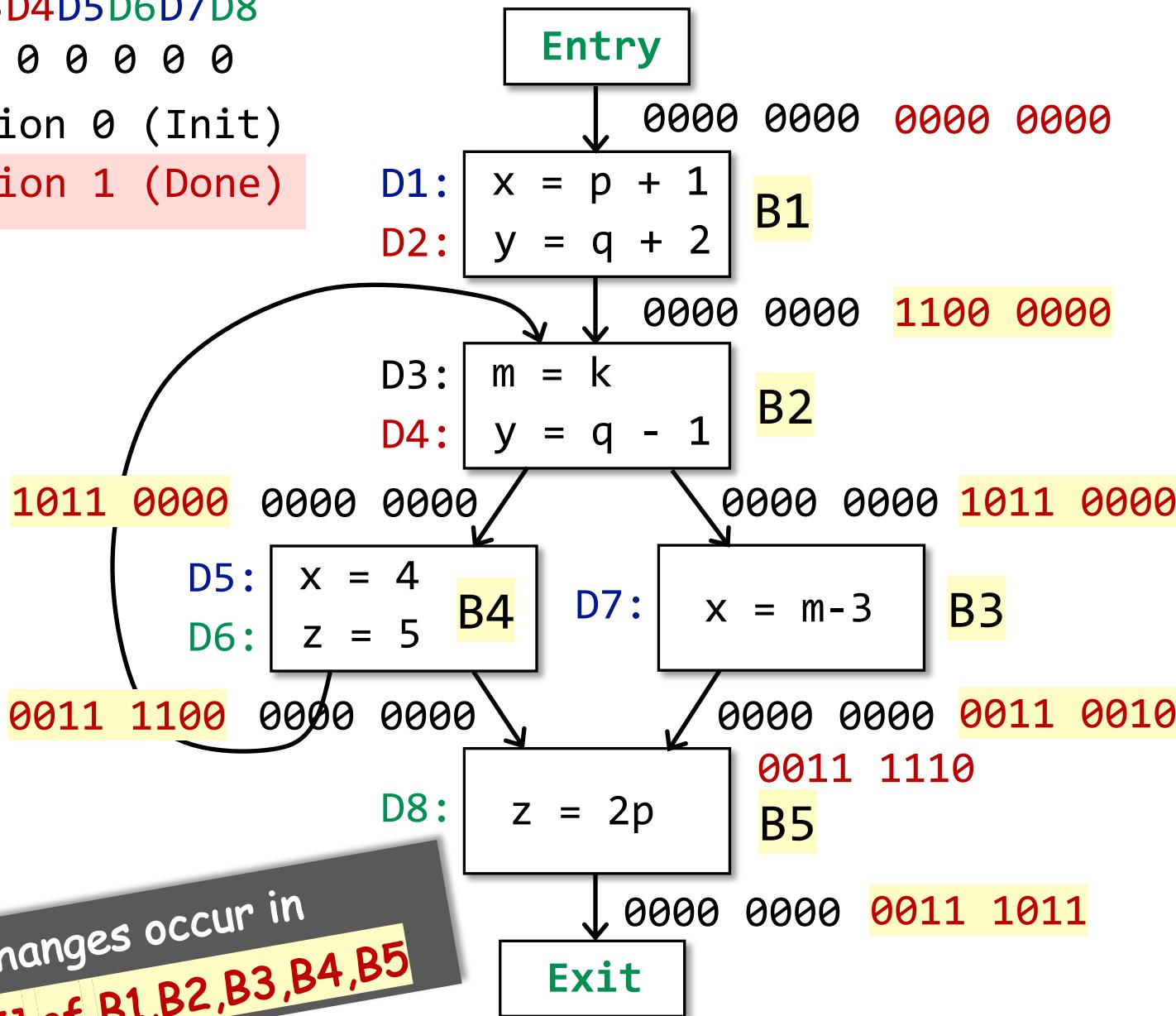


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)



Changes occur in  
**OUT[] of B1, B2, B3, B4, B5**

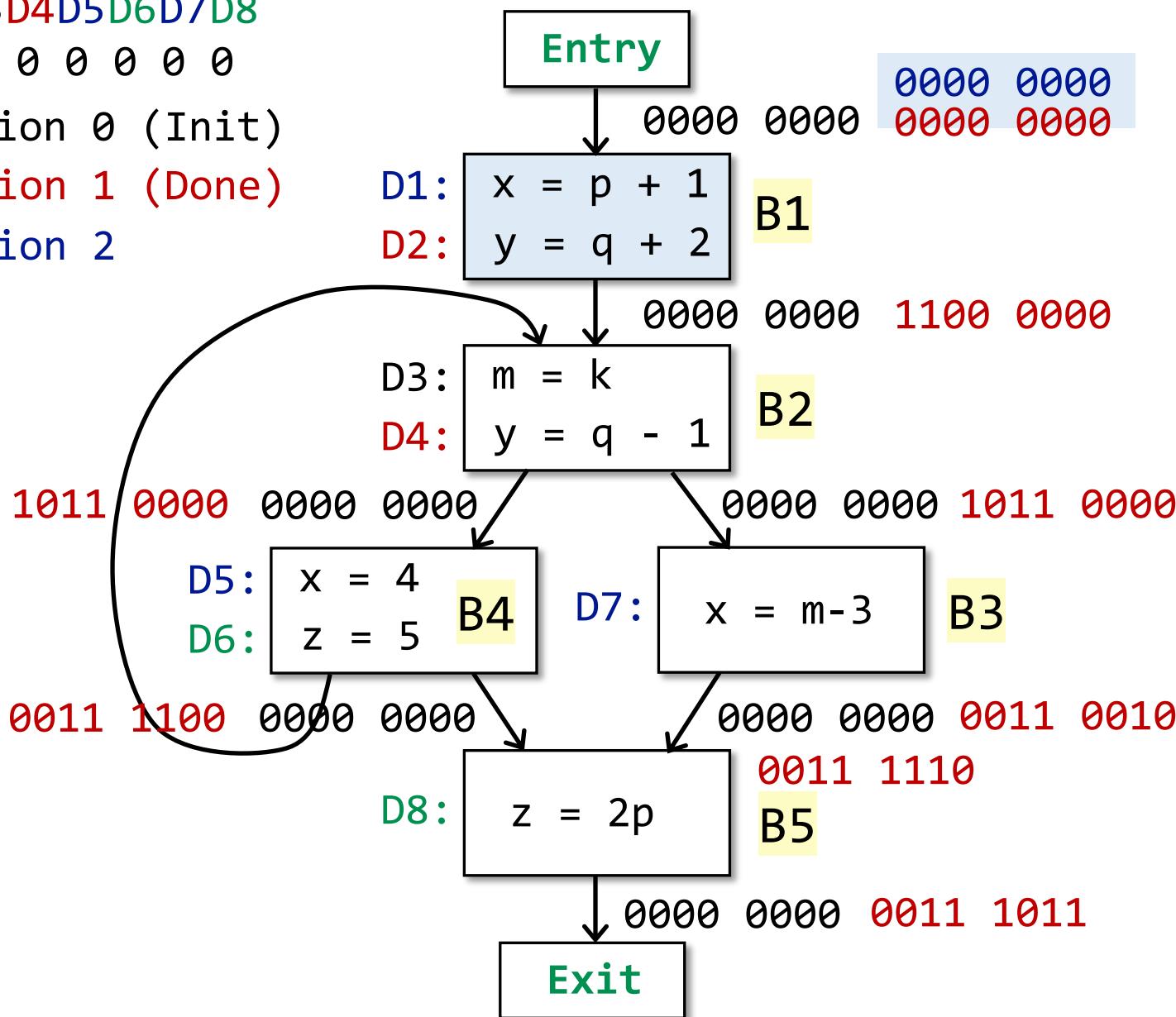
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



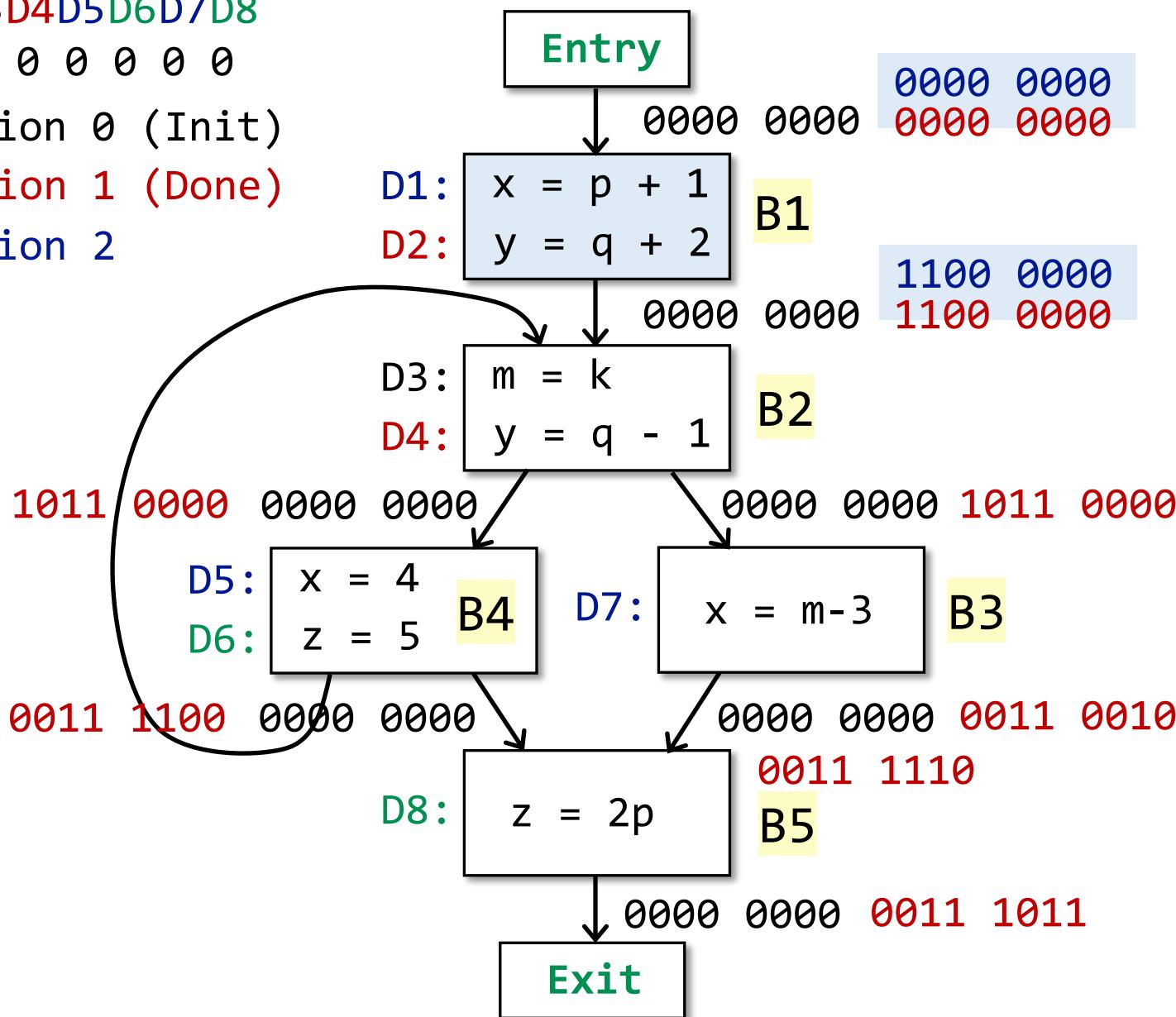
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



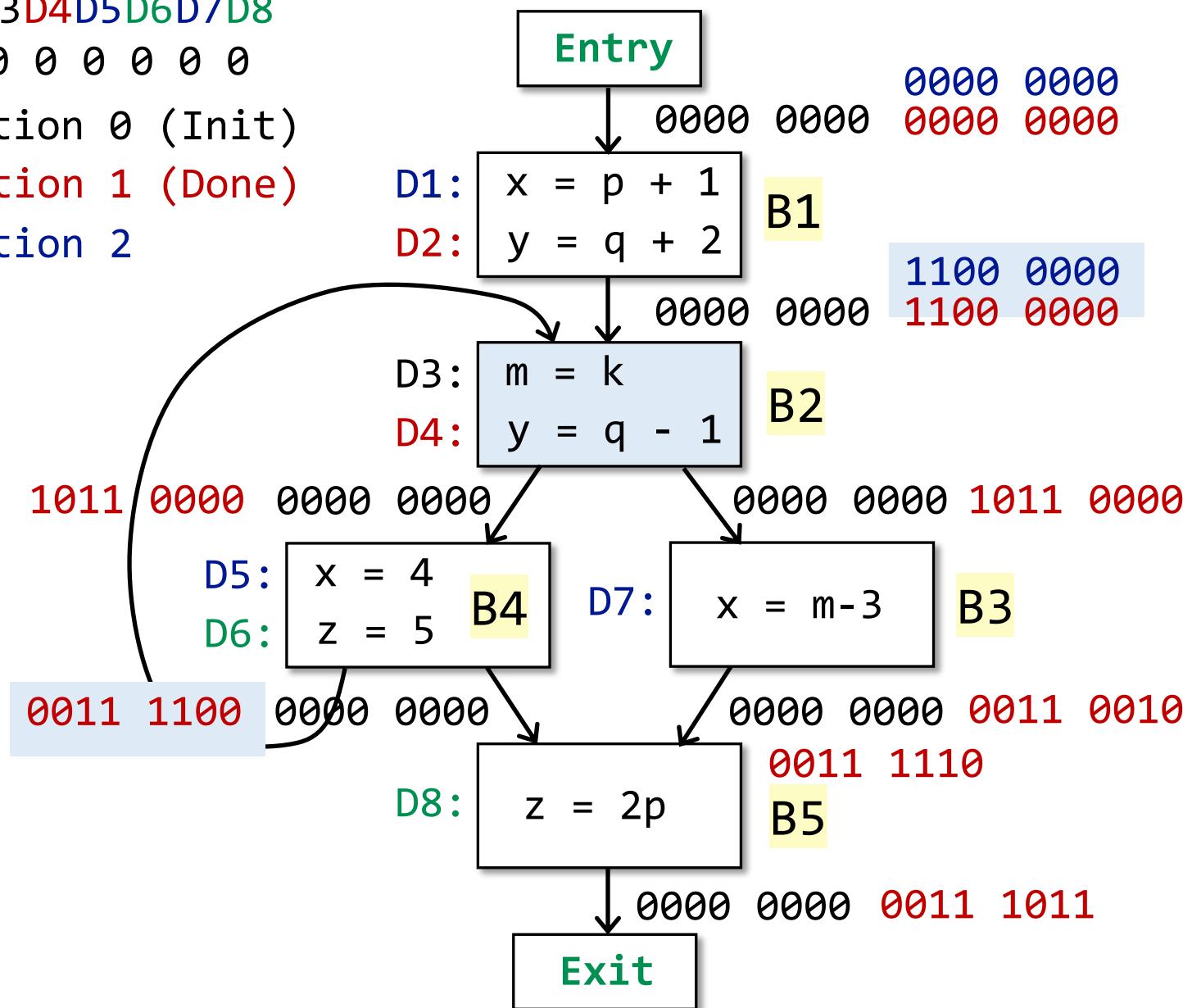
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



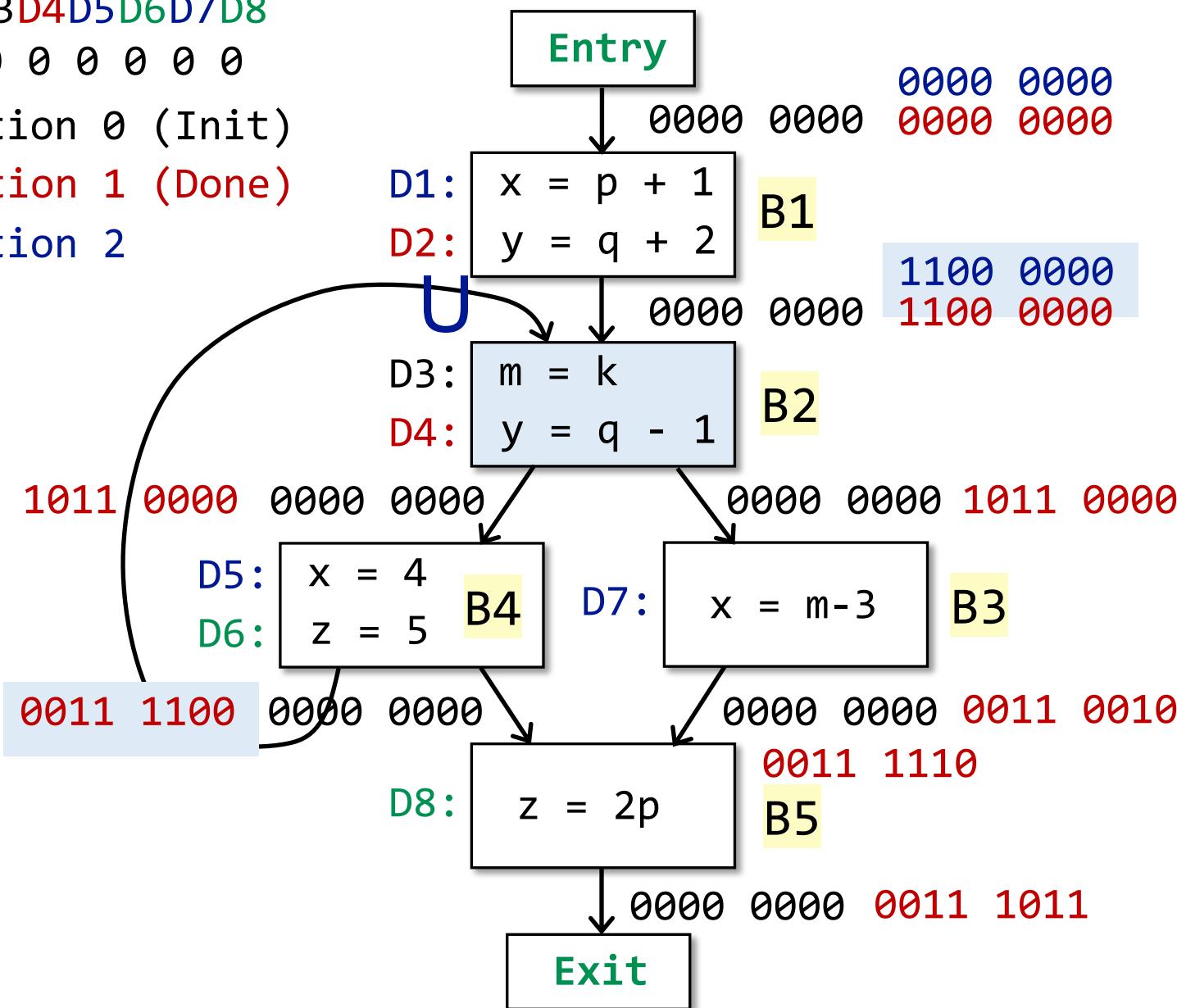
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



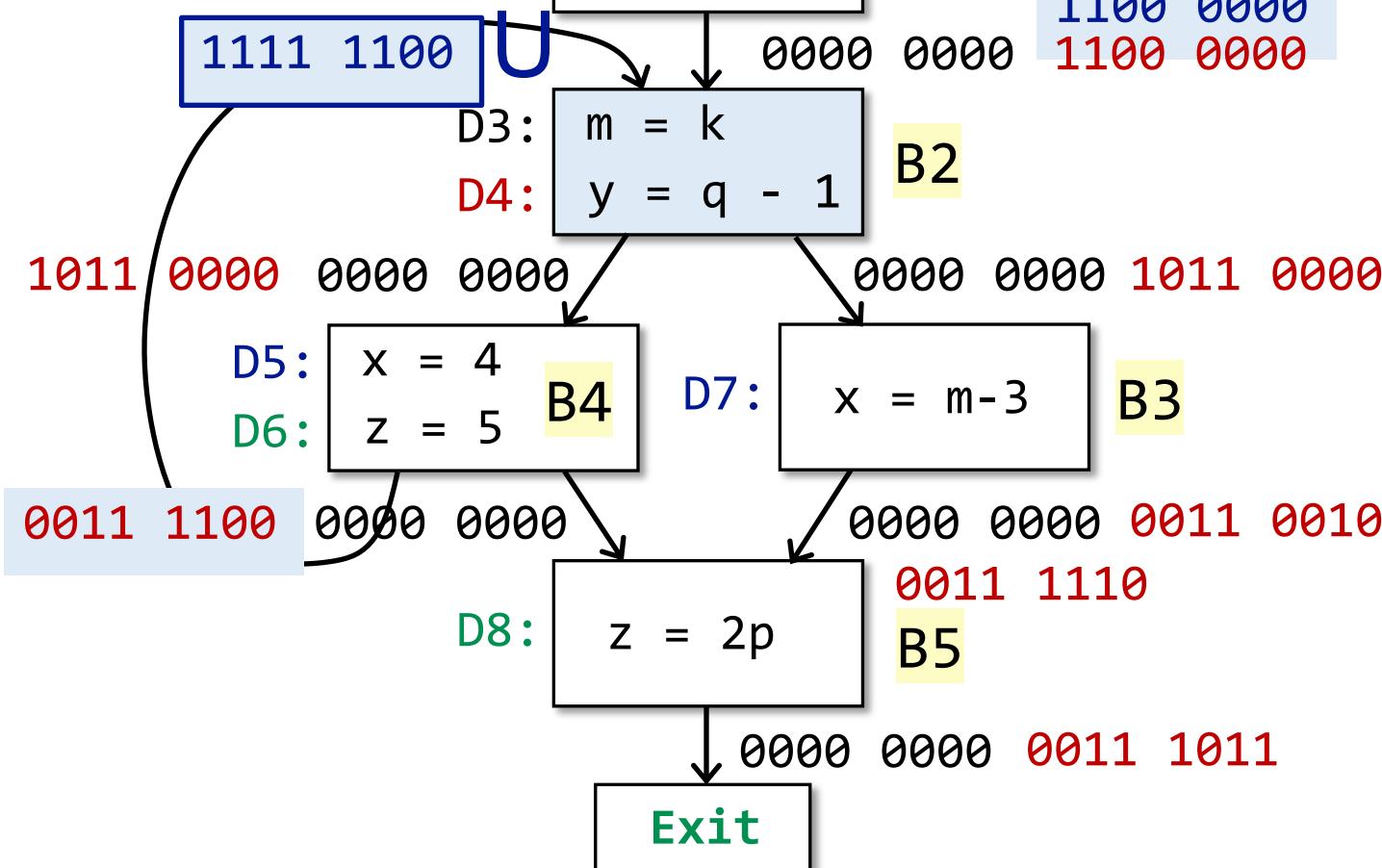
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



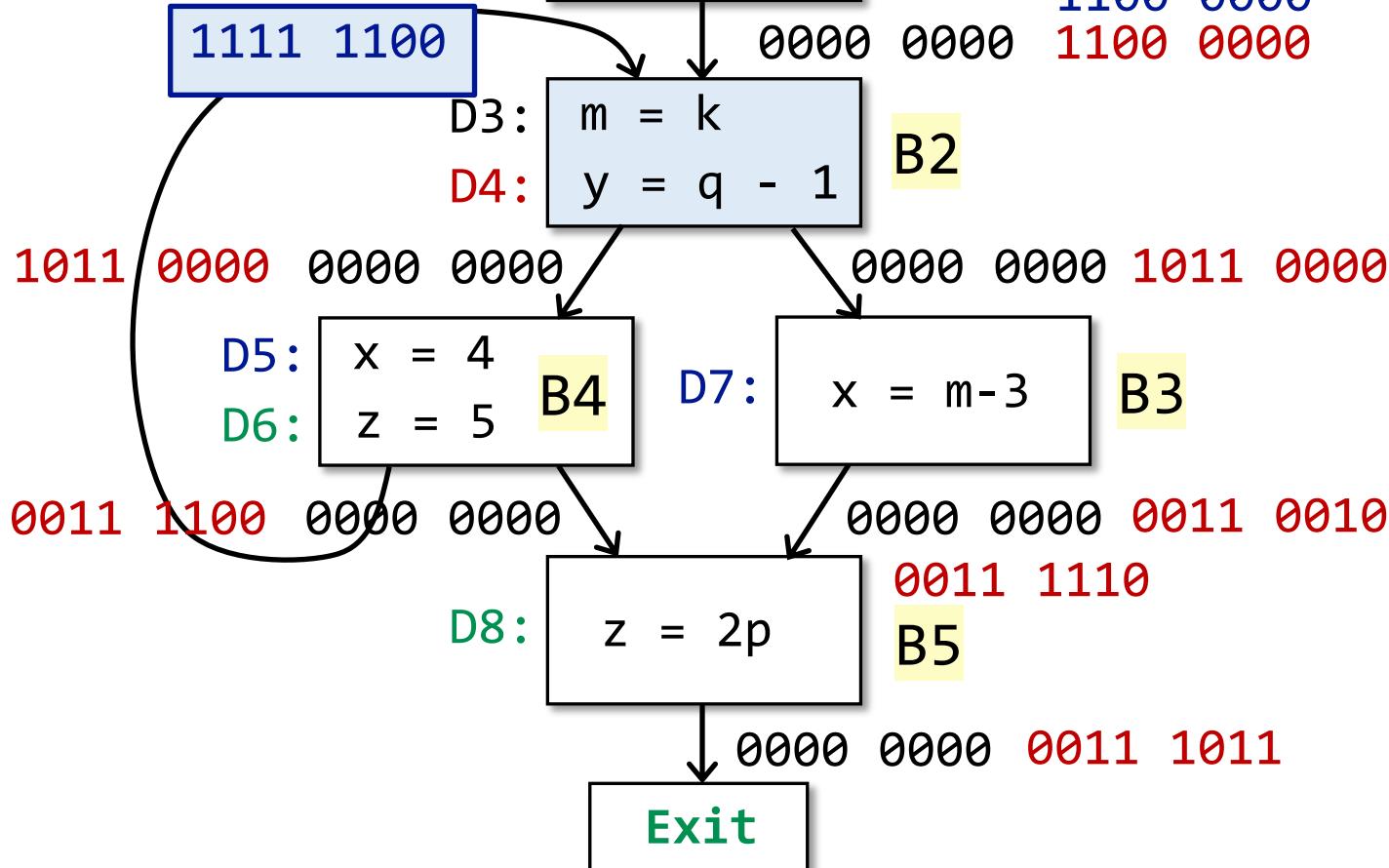
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



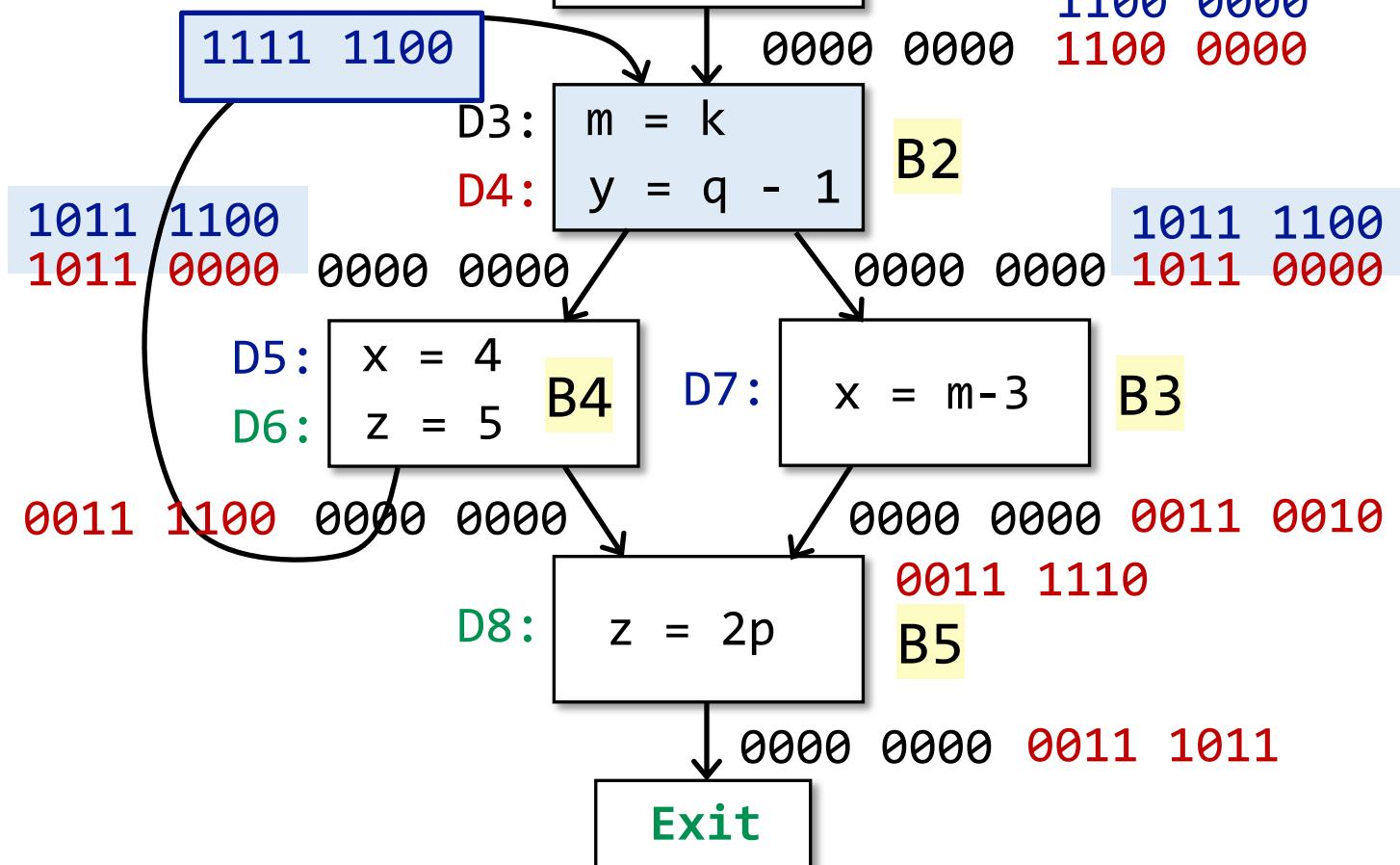
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100  
0000 0000

Entry

0000 0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:

$x = m - 3$

1011 1100  
1011 0000

B3

D8:  $z = 2p$

0011 1110  
B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

D5:  $x = 4$   
D6:  $z = 5$  B4

0011 1100 0000 0000

D8:  $z = 2p$  B5

0011 1110

Exit

Entry

0000 0000

0000 0000

0000 0000

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

D3:  $m = k$

D4:  $y = q - 1$

B2

D7:  $x = m - 3$

1011 1100  
1011 0000

B3

0011 0110  
0011 0010

0000 0000

0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100

Entry

0000 0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

D8:  $z = 2p$

0011 1110

B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

1011 1100  
1011 0000

B3

D8:  $z = 2p$

0011 1110  
B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

1011 1100  
1011 0000

0011 0110  
0011 0010

D8:  $z = 2p$

0011 1110  
B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

D8:  $z = 2p$

Exit

1100 0000  
1100 0000

1011 1100  
1011 0000

0011 0110  
0011 0010

0011 1110  
B5

0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

1011 1100  
1011 0000

B3

D8:  $z = 2p$

B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000

0000

0000 0000

0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

D8:  $z = 2p$

Exit

0000 0000

0011 1110

0011 0110

0011 0010

0011 1110

0000 0000

0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

D8:  $z = 2p$

B5

Exit

1100 0000  
1100 0000

1011 1100  
1011 0000

0011 0110  
0011 0010

0011 1110 0011 1110

0011 1011  
0011 1011

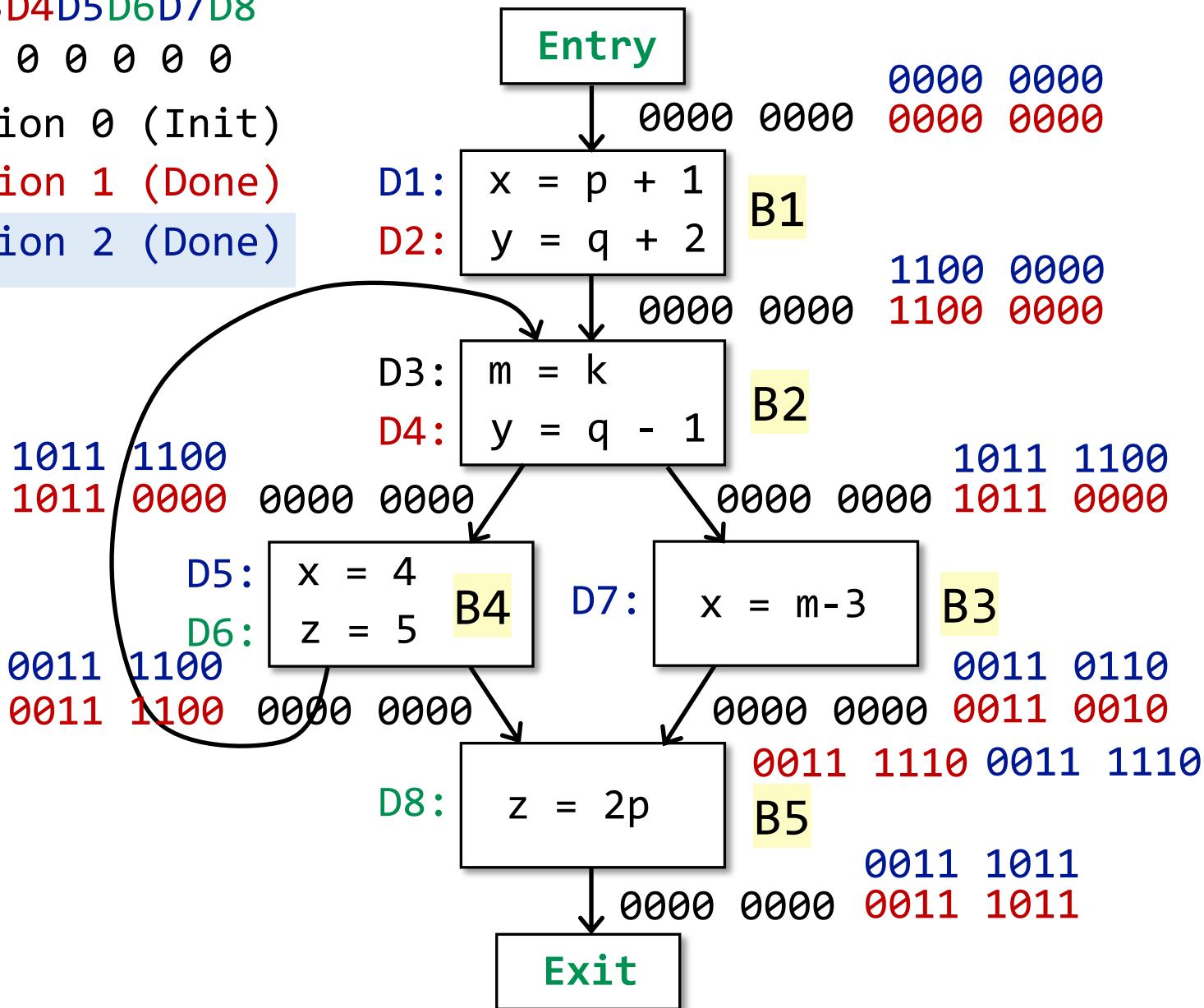
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

1011 1100  
1011 0000

0011 1100  
0011 1100

Changes occur in  
OUT[] of B2,B3

Entry

0000

0000

0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:

$x = m - 3$

1011 1100  
1011 0000

B3

D8:  $z = 2p$

B5

0011 1011  
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

|      |      |
|------|------|
| 0000 | 0000 |
| 0000 | 0000 |
| 0000 | 0000 |

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

D8:  $z = 2p$

0011 1110 0011 1110  
0011 0110  
0011 0010

0011 1011  
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

|      |      |
|------|------|
| 0000 | 0000 |
| 0000 | 0000 |
| 0000 | 0000 |

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

|      |      |
|------|------|
| 1100 | 0000 |
| 1100 | 0000 |
| 1100 | 0000 |

D3:  $m = k$   
D4:  $y = q - 1$

B2

|      |      |
|------|------|
| 1011 | 1100 |
| 1011 | 0000 |

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

|      |      |
|------|------|
| 0011 | 0110 |
| 0011 | 0010 |

D8:  $z = 2p$

B5

|      |      |      |      |
|------|------|------|------|
| 0011 | 1110 | 0011 | 1110 |
| 0011 | 1011 | 0011 | 1011 |

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

0000 0000  
0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

1100 0000  
1100 0000  
1100 0000

D3:  $m = k$   
D4:  $y = q - 1$

B2

1011 1100  
1011 0000

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

0011 0110  
0011 0010

D8:  $z = 2p$

B5

0011 1110 0011 1110  
0011 1011  
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 0000

0011 1100  
0011 1100

U

Entry

0000 0000

0000 0000  
0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

1100 0000  
1100 0000  
1100 0000

D3:  $m = k$   
D4:  $y = q - 1$

B2

1011 1100  
1011 0000

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

0011 0110  
0011 0010

D8:  $z = 2p$

0011 1110 0011 1110  
B5

0011 1011  
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 0000

0011 1100  
0011 1100

U

Entry

0000 0000

0000 0000  
0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

1100 0000  
1100 0000  
1100 0000

D3:  $m = k$   
D4:  $y = q - 1$

B2

1011 1100  
1011 0000

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:

$x = m - 3$

B3

0011 0110  
0011 0010

D8:  $z = 2p$

B5

0011 1011  
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 0000

0011 1100  
0011 1100

1111 1100

0000 0000

0000 0000  
0000 0000  
0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

1011 1100  
1011 0000

0011 0110  
0011 0010

D8:  $z = 2p$

0011 1110 0011 1110  
B5

0011 1011  
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

1111 1100

D3:  $m = k$

D4:  $y = q - 1$

B2

1011 1100  
1011 1100  
1011 0000

D5:  $x = 4$

D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

0011 0110  
0011 0010

D8:  $z = 2p$

B5

0011 1110 0011 1110  
0011 1011  
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100

1111 1100

0000 0000

|      |      |
|------|------|
| 0000 | 0000 |
| 0000 | 0000 |
| 0000 | 0000 |

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

|      |      |
|------|------|
| 1100 | 0000 |
| 1100 | 0000 |
| 1100 | 0000 |

D3:  $m = k$   
D4:  $y = q - 1$

B2

|      |      |
|------|------|
| 1011 | 1100 |
| 1011 | 1100 |
| 1011 | 0000 |

D5:  $x = 4$   
D6:  $z = 5$

B4

0000 0000

D7:  $x = m - 3$

B3

0000 0000

|      |      |
|------|------|
| 0011 | 0110 |
| 0011 | 0010 |

D8:  $z = 2p$

B5

0000 0000

|      |      |      |      |
|------|------|------|------|
| 0011 | 1110 | 0011 | 1110 |
| 0011 | 1011 | 0011 | 1011 |

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100

1111 1100

0000 0000

0000 0000

|      |      |
|------|------|
| 0000 | 0000 |
| 0000 | 0000 |
| 0000 | 0000 |

|      |      |
|------|------|
| 1100 | 0000 |
| 1100 | 0000 |
| 1100 | 0000 |

|      |      |
|------|------|
| 1011 | 1100 |
| 1011 | 1100 |
| 1011 | 0000 |

|      |      |
|------|------|
| 0011 | 0110 |
| 0011 | 0110 |
| 0011 | 0010 |

|      |      |
|------|------|
| 0011 | 1011 |
| 0011 | 1011 |

Entry

0000 0000

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

D3:  $m = k$   
D4:  $y = q - 1$

B2

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

D8:  $z = 2p$

B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

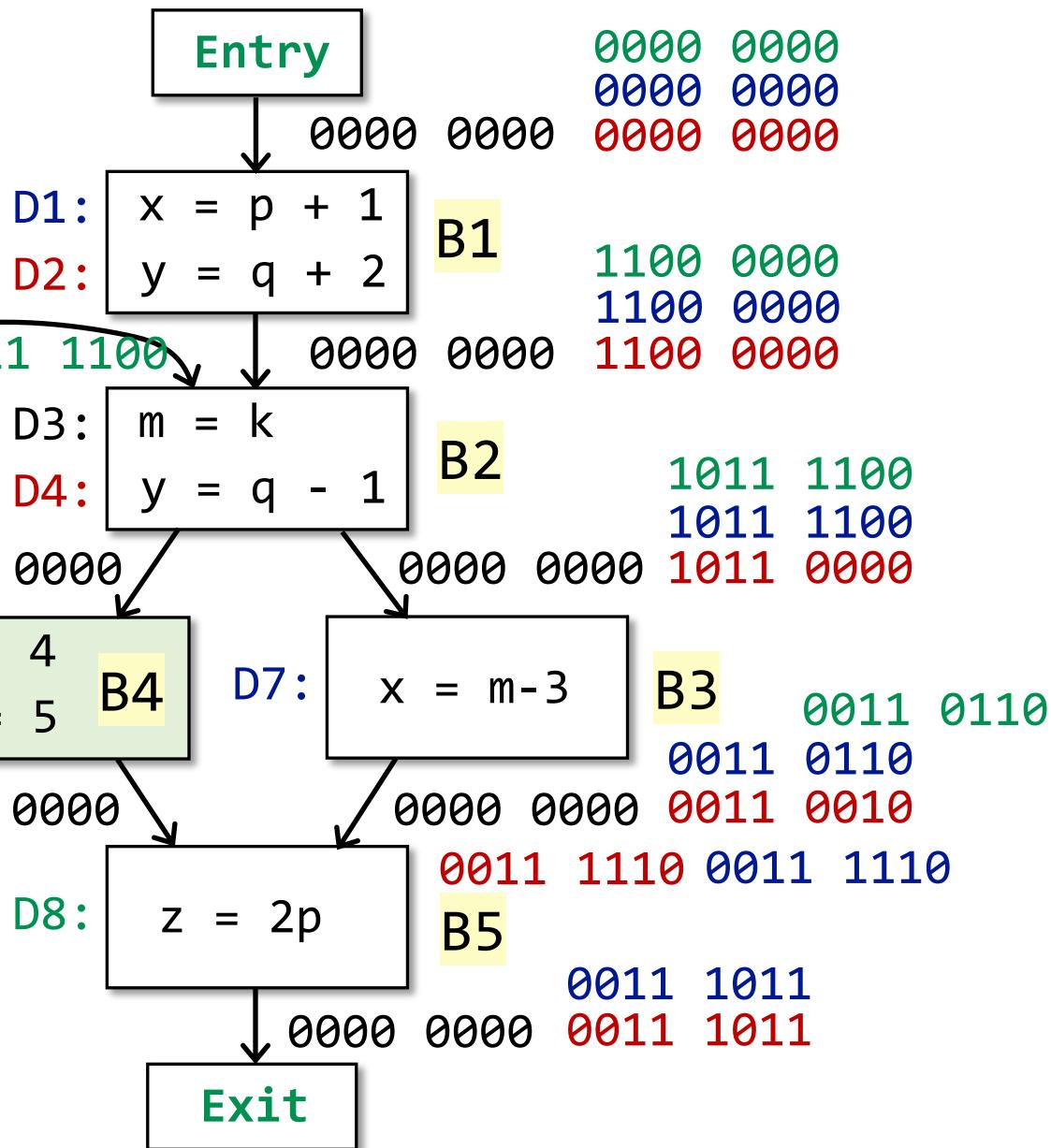
Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100  
0011 1100

Entry

0000

0000

|      |      |
|------|------|
| 0000 | 0000 |
| 0000 | 0000 |
| 0000 | 0000 |

D1:  $x = p + 1$   
D2:  $y = q + 2$

B1

|      |      |
|------|------|
| 1100 | 0000 |
| 1100 | 0000 |
| 1100 | 0000 |

D3:  $m = k$   
D4:  $y = q - 1$

B2

|      |      |
|------|------|
| 1011 | 1100 |
| 1011 | 1100 |
| 1011 | 0000 |

D5:  $x = 4$   
D6:  $z = 5$

B4

D7:  $x = m - 3$

B3

0011 0110

|      |      |
|------|------|
| 0011 | 0110 |
| 0011 | 0010 |

D8:  $z = 2p$

B5

|      |      |      |      |
|------|------|------|------|
| 0011 | 1110 | 0011 | 1110 |
| 0011 | 1011 | 0011 | 1011 |

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3:  $m = k$

D4:  $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5:  $x = 4$

D6:  $z = 5$

B4

B3

0011 0110

0011 0110

0011 0010

D7:  $x = m - 3$

D8:  $z = 2p$

B5

0011 1110 0011 1110

0011 1011

0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3:  $m = k$

D4:  $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5:  $x = 4$

D6:  $z = 5$

B4

0011 0110

0011 0110

0011 0010

D7:  $x = m - 3$

B3

0011 1110

0011 0011

0011 1110

D8:  $z = 2p$

B5

0011 1011

0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3:  $m = k$

D4:  $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5:  $x = 4$

D6:  $z = 5$

B4

B7

B3

0011 0110

0011 0110

0011 0010

D7:  $x = m - 3$

D8:  $z = 2p$

B4

B5

0011 1011

0011 1011

Exit

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100  
0011 1100

0011 1110

0011 0110

0011 0110

0011 0010

0011 1110 0011 1110

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100  
0011 1100

D1:  $x = p + 1$   
D2:  $y = q + 2$

D3:  $m = k$   
D4:  $y = q - 1$

D5:  $x = 4$   
D6:  $z = 5$

D7:  $x = m - 3$

D8:  $z = 2p$

Entry

Exit

0000 0000  
0000 0000  
0000 0000

1100 0000  
1100 0000  
1100 0000

1011 1100  
1011 1100  
1011 0000

0011 0110  
0011 0110  
0011 0010

0011 1110 0011 1110  
B5

0011 1011  
0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100  
1011 1100  
1011 0000

0011 1100  
0011 1100  
0011 1100

1111 1100

0000 0000

0000 0000

0011 1110

D1:  $x = p + 1$   
D2:  $y = q + 2$

D3:  $m = k$   
D4:  $y = q - 1$

D5:  $x = 4$

D6:  $z = 5$

Entry

Exit

B1

B2

B3

B4

B5

0000 0000  
0000 0000  
0000 0000

1100 0000  
1100 0000  
1100 0000

1011 1100  
1011 1100  
1011 0000

0011 0110  
0011 0110  
0011 0010

0011 1011  
0011 1011  
0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3:  $m = k$

D4:  $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5:  $x = 4$

D6:  $z = 5$

B4

0011 0110

0011 0110

0011 0010

D7:  $x = m - 3$

B3

0011 0110

0011 0010

0011 1110

D8:  $z = 2p$

B5

0011 1011

0011 1011

0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3:  $m = k$

D4:  $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5:  $x = 4$

D6:  $z = 5$

B4

0011 0110

0011 0110

0011 0010

D7:  $x = m - 3$

B3

0011 0110

0011 0010

0011 1110

D8:  $z = 2p$

B5

0011 1011

0011 1011

0011 1011

Exit

No changes occur  
in any OUT[ ]

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

0000 0000

|      |      |
|------|------|
| 0000 | 0000 |
| 0000 | 0000 |
| 0000 | 0000 |

D1:  $x = p + 1$

D2:  $y = q + 2$

B1

|      |      |
|------|------|
| 1100 | 0000 |
| 1100 | 0000 |
| 1100 | 0000 |

D3:  $m = k$

D4:  $y = q - 1$

B2

|      |      |
|------|------|
| 1011 | 1100 |
| 1011 | 1100 |
| 1011 | 0000 |

D5:  $x = 4$

D6:  $z = 5$

B4

|      |      |
|------|------|
| 0011 | 1100 |
| 0011 | 1100 |
| 0011 | 1100 |

D7:  $x = m - 3$

B3

|      |      |
|------|------|
| 0011 | 0110 |
| 0011 | 0110 |
| 0011 | 0010 |

D8:  $z = 2p$

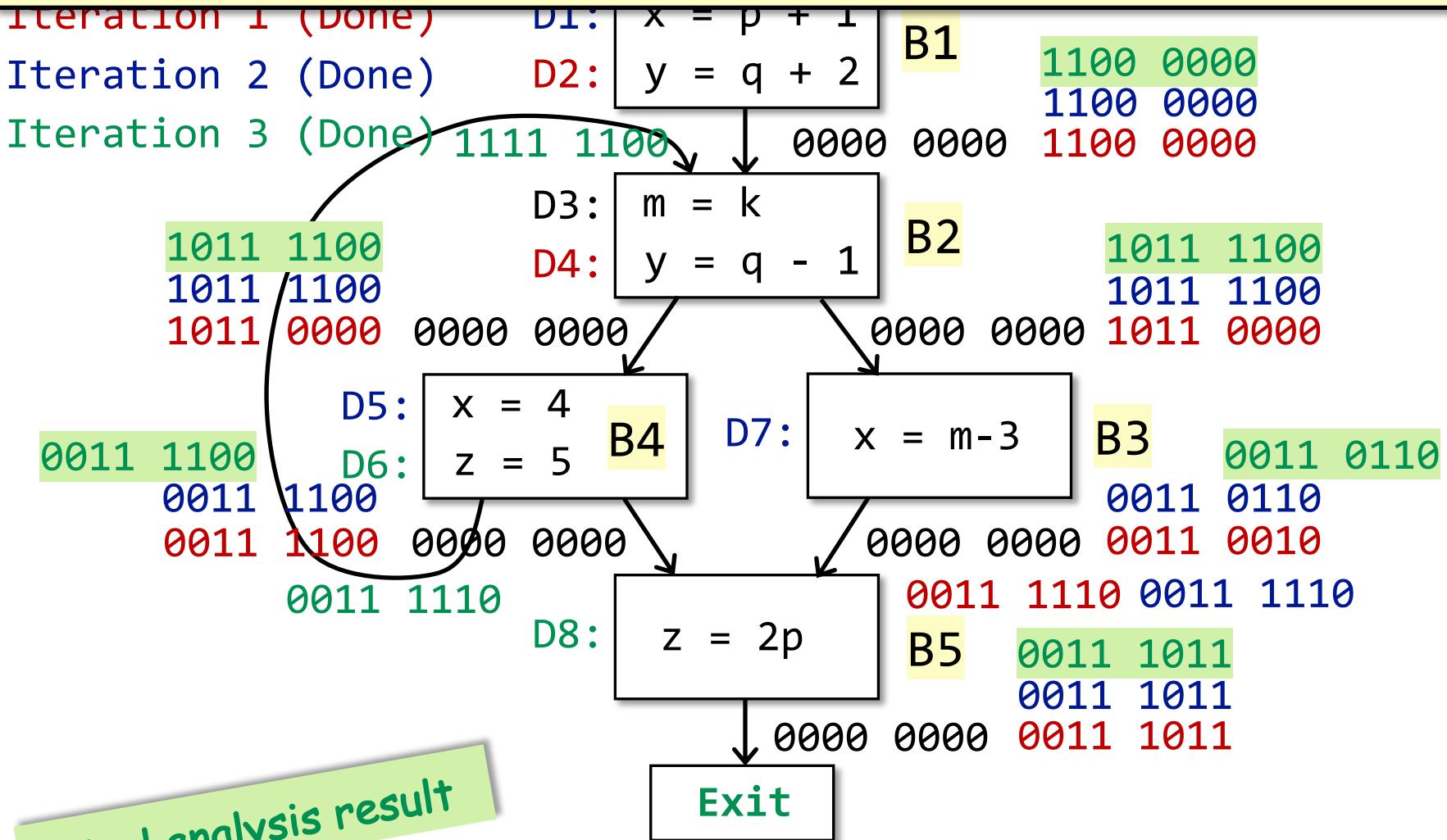
B5

|      |      |
|------|------|
| 0011 | 1011 |
| 0011 | 1011 |
| 0011 | 1011 |

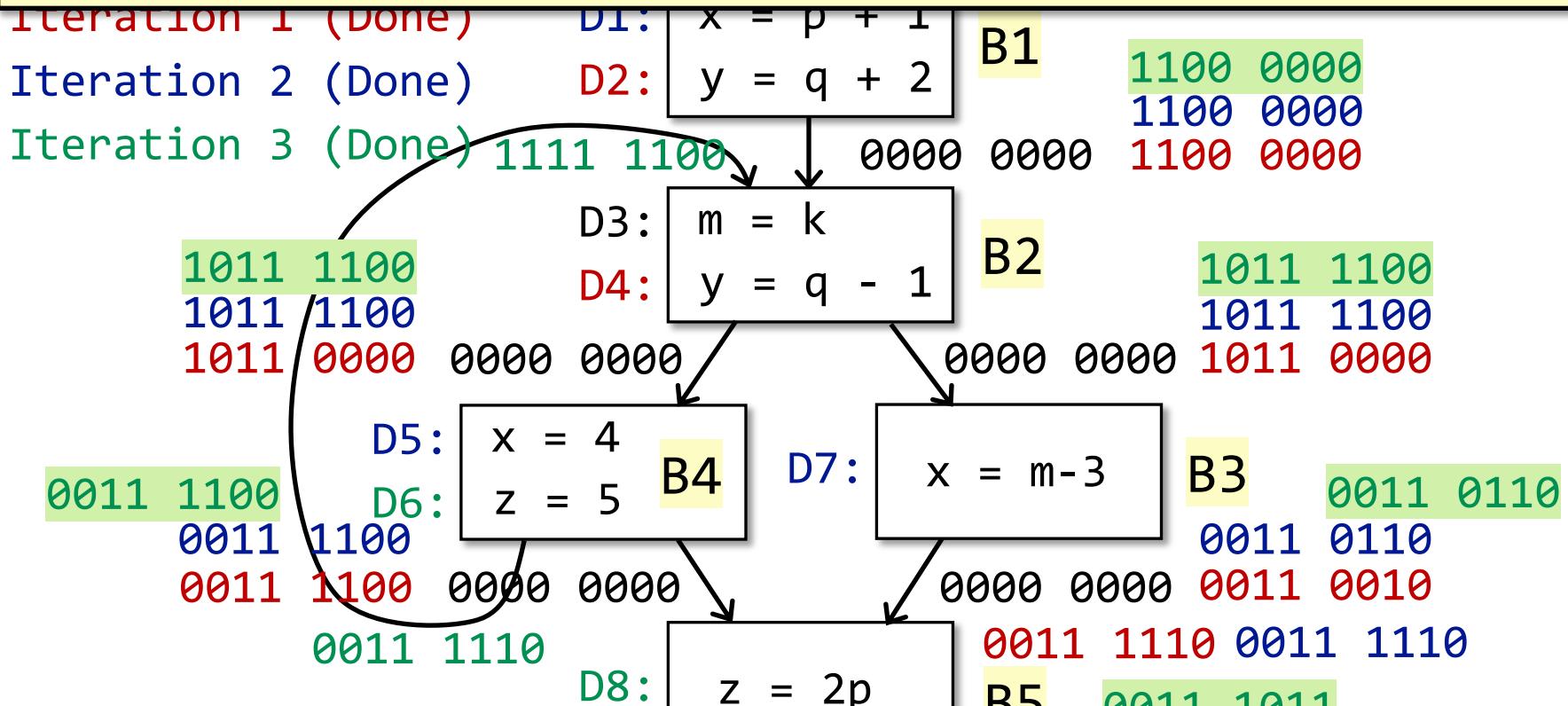
Exit

Final analysis result

In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.



In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.



Data-flow analysis is to find a solution to a set of *safe-approximation-directed constraints* on the  $\text{IN}[s]$ 's and  $\text{OUT}[s]$ 's, for all statements  $s$ .

- constraints based on semantics of statements (*transfer functions*)
- constraints based on the flows of control

# Algorithm of Reaching Definitions Analysis

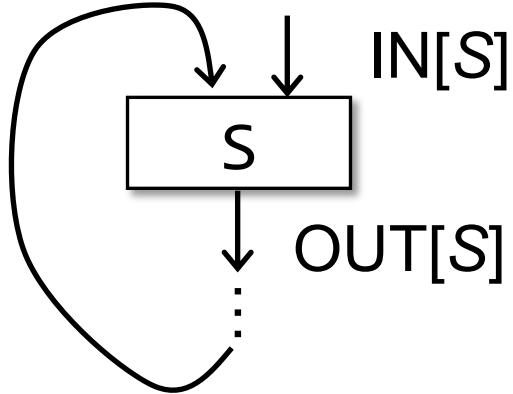
**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

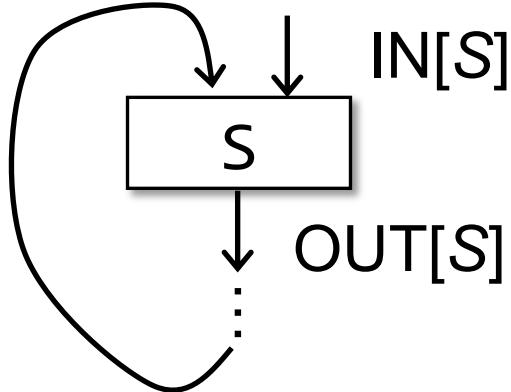
**METHOD:**

```
OUT[entry] = Ø;
for (each basic block B) {
    OUT[B] = Ø;
    while (changes to any OUT occur)
        for (each basic block B\entry) {
            IN[B] =  $\bigcup_{P \text{ a predecessor of } B}$  OUT[P];
            OUT[B] =  $gen_B \cup (IN[B] - kill_B)$ ;
        }
}
```

*Why this iterative algorithm can finally stop?*

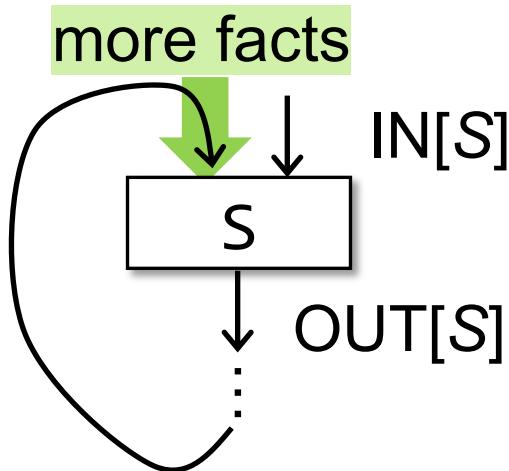


$$\text{OUT}[S] = \text{gen}_S \cup (\text{IN}[S] - \text{kill}_S);$$



$$\text{OUT}[S] = \text{gen}_S \cup (\text{IN}[S] - \text{kill}_S);$$

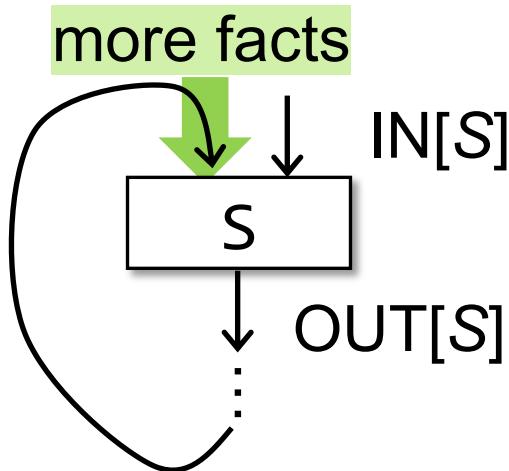
- $\text{gen}_S$  and  $\text{kill}_S$  remain unchanged



more facts

$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

- $gen_S$  and  $kill_S$  remain unchanged
- When more facts flow in  $IN[S]$ , the “more facts” either

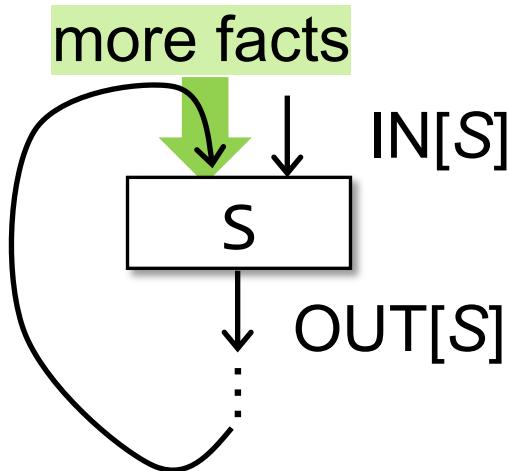


$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

$\underbrace{\phantom{gen_S \cup (}}_{\text{survivors}}$

more facts

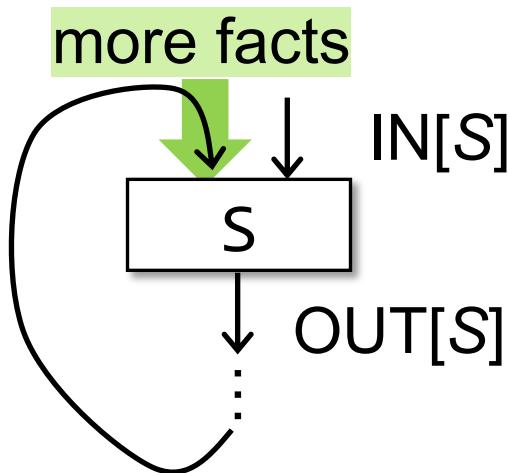
- $gen_S$  and  $kill_S$  remain unchanged
- When **more facts** flow in  $IN[S]$ , the “**more facts**” either
  - is killed, or
  - flows to  $OUT[S]$  ( $\text{survivor}_S$ )



$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

survivors<sub>S</sub>

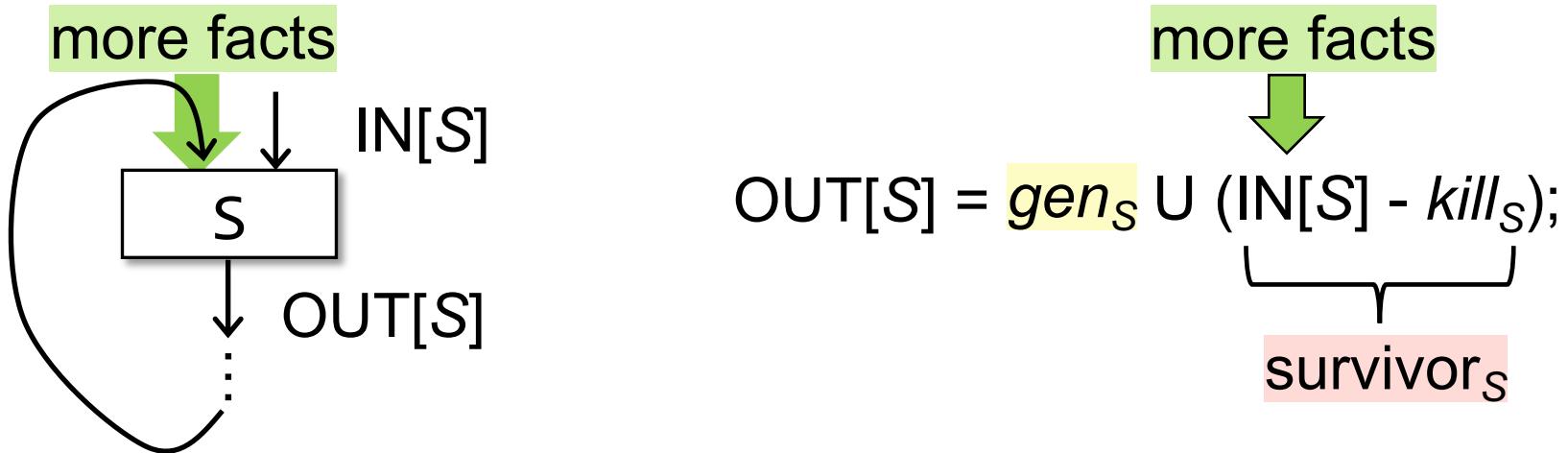
- $gen_S$  and  $kill_S$  remain unchanged
- When "more facts" flow in  $IN[S]$ , the "more facts" either
  - is killed, or
  - flows to  $OUT[S]$  ( $survivors_S$ )
- When a fact is added to  $OUT[S]$ , through either  $gen_S$ , or  $survivors_S$ , it stays there forever



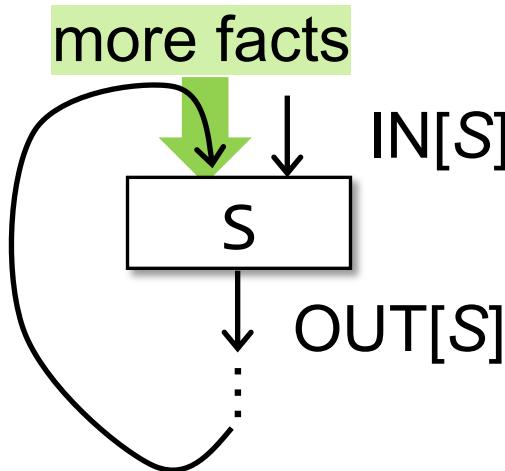
$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

survivors<sub>S</sub>

- $gen_S$  and  $kill_S$  remain unchanged
- When "more facts" flow in  $IN[S]$ , the "more facts" either
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- Thus  $OUT[S]$  never shrinks (e.g.,  $0 \rightarrow 1$ , or  $1 \rightarrow 1$ )



- $gen_S$  and  $kill_S$  remain unchanged
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- When a fact is added to  $OUT[S]$ , through either  $gen_S$ , or  $survivor_S$ , it stays there forever
- Thus  $OUT[S]$  never shrinks (e.g.,  $0 \rightarrow 1$ , or  $1 \rightarrow 1$ )
- As the set of facts is finite (e.g., all definitions in the program),



$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

survivors

- $gen_S$  and  $kill_S$  remain unchanged
- When **more facts** flow in  $IN[S]$ , the “**more facts**” either
  - is killed, or
  - flows to  $OUT[S]$  (**survivor** $_S$ )
- When a fact is added to  $OUT[S]$ , through either  $gen_S$ , or **survivor** $_S$ , it stays there forever
- Thus  $OUT[S]$  never shrinks (e.g.,  $0 \rightarrow 1$ , or  $1 \rightarrow 1$ )
- As the set of facts is finite (e.g., all definitions in the program), there must exist a pass of iteration during which nothing is added to any  $OUT$ , and then the algorithm terminates

# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ ) {  
    OUT[ $B$ ] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[ $B$ ] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P]$ ;  
            OUT[ $B$ ] =  $gen_B \cup (IN[B] - kill_B)$ ;  
        }  
}
```

*Safe to terminate  
by this condition?*

# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

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            OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
        }  
}
```

*Safe to terminate  
by this condition?*

IN's will not change if  
OUT's do not change

# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ ) {  
    OUT[B] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {
```

*Safe to terminate  
by this condition?*

IN's will not change if  
OUT's do not change

$$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P];$$

OUT's will not change  
if IN's do not change

$$OUT[B] = gen_B \cup (IN[B] - kill_B);$$

# Algorithm of Reaching Definitions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ ) {  
    OUT[B] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {
```

*Safe to terminate  
by this condition?*

IN's will not change if  
OUT's do not change

$$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P];$$

OUT's will not change  
if IN's do not change

$$OUT[B] = gen_B \cup (IN[B] - kill_B);$$

Reach a **fixed point**  
Also related with **monotonicity**  
(next lectures)

# Data Flow Analysis Applications

(I) Reaching Definitions Analysis

(II) Live Variables Analysis

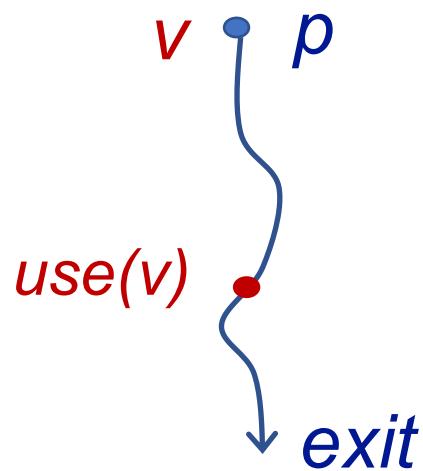
(III) Available Expressions Analysis

# Live Variables Analysis

Live variables analysis tells whether the value of **variable v** at **program point p** could be used along some path in CFG starting at p. If so, v is live at p; otherwise, v is dead at p.

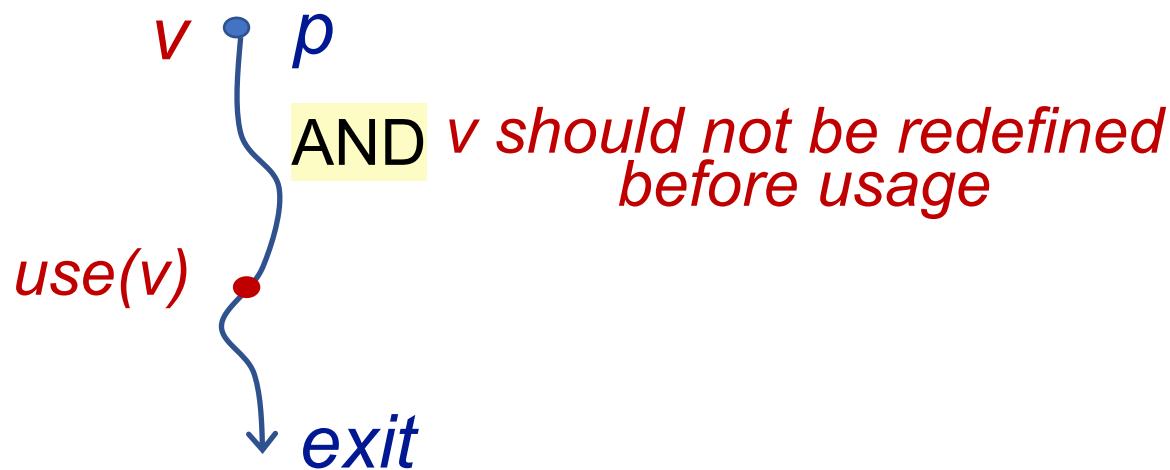
# Live Variables Analysis

Live variables analysis tells whether the value of **variable v** at **program point p** could be used along some path in CFG starting at p. If so, v is live at p; otherwise, v is dead at p.



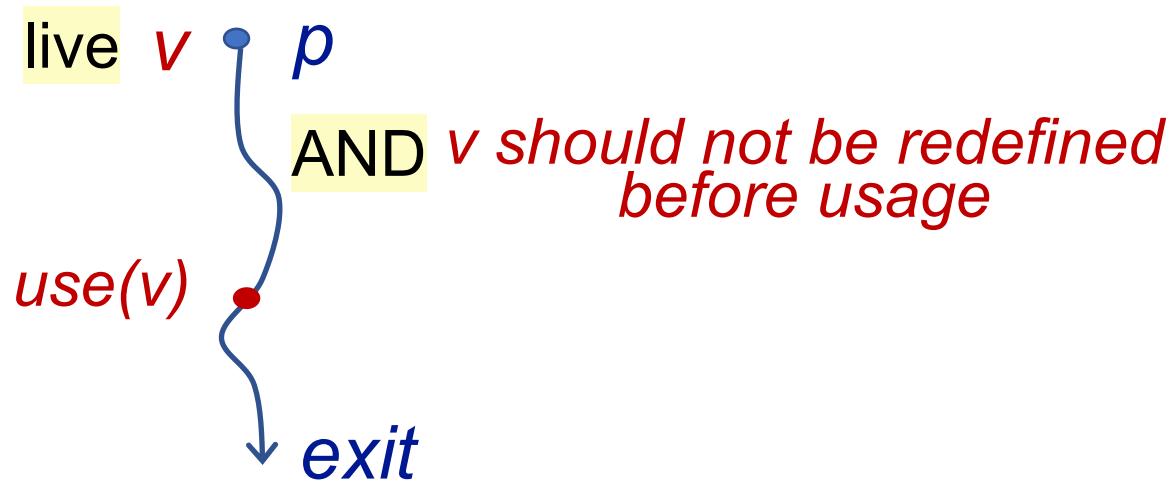
# Live Variables Analysis

Live variables analysis tells whether the value of **variable v** at **program point p** could be used along some path in CFG starting at p. If so, v is live at p; otherwise, v is dead at p.



# Live Variables Analysis

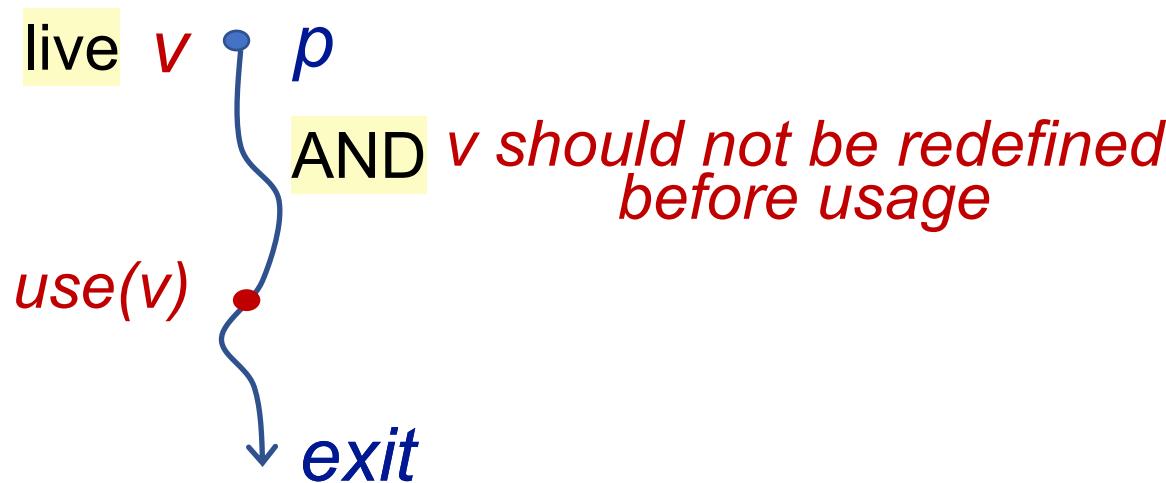
Live variables analysis tells whether the value of **variable v** at **program point p** could be used along some path in CFG starting at p. If so, v is live at p; otherwise, v is dead at p.



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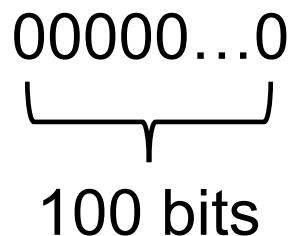
- Information of live variables can be used for register allocations. e.g., at some point all registers are full and we need to use one, then we should favor using a register with a dead value.



# Understanding Live Variables Analysis

- Data Flow Values/Facts
    - All the variables in a program
    - Can be represented by bit vectors
- e.g., V1, V2, V3, V4, ..., V100 (100 variables)

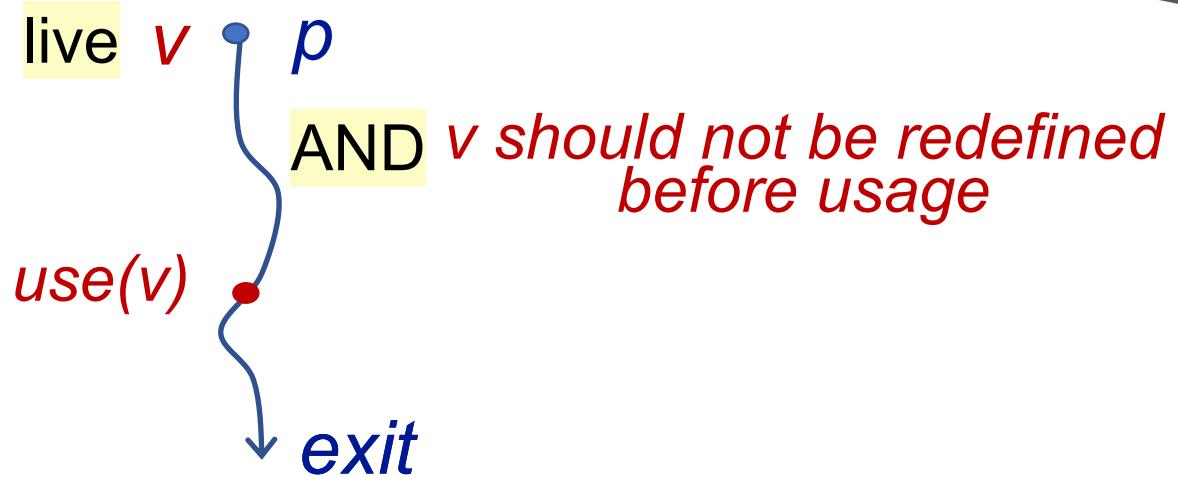
*Abstraction*



Bit i from the left represents variable Vi

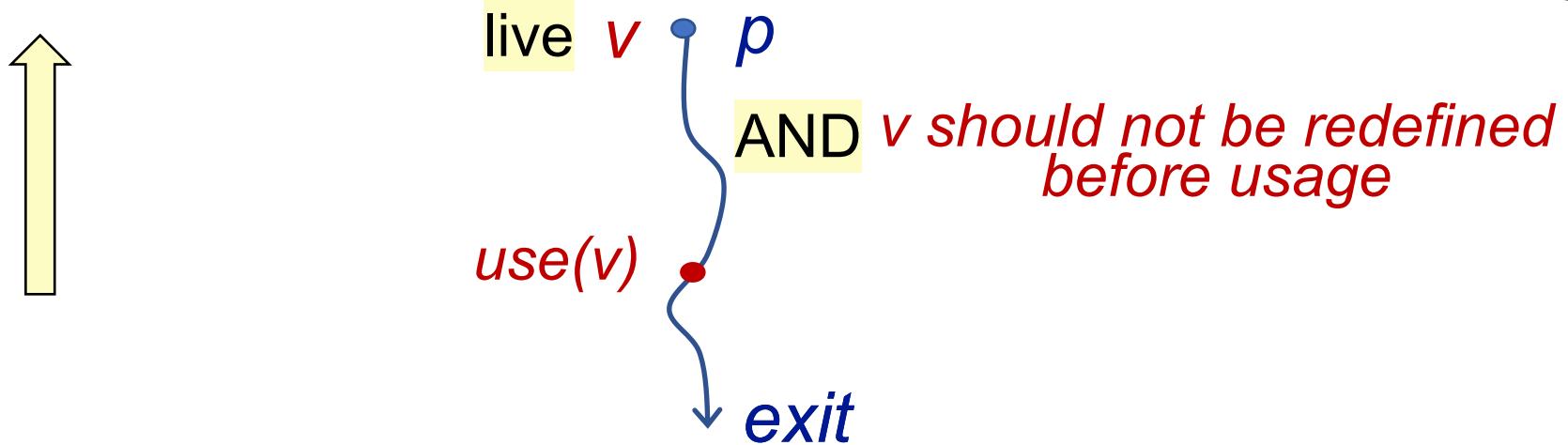
# Understanding Live Variables Analysis

Safe-approximation



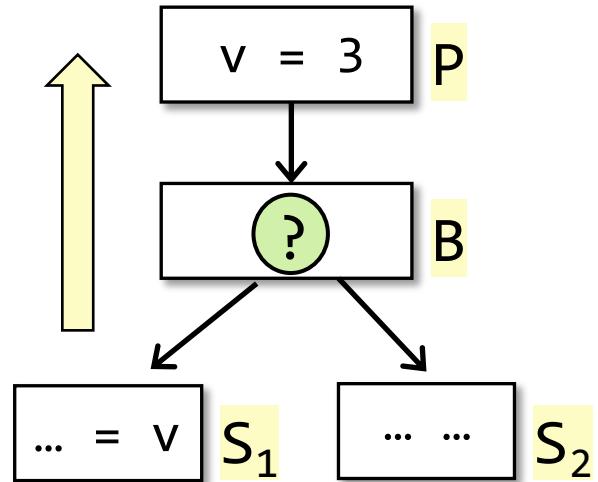
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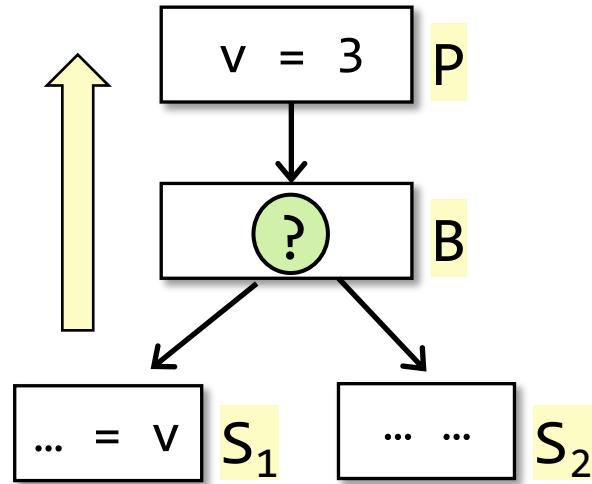


# Understanding Live Variables Analysis

*Safe-approximation*

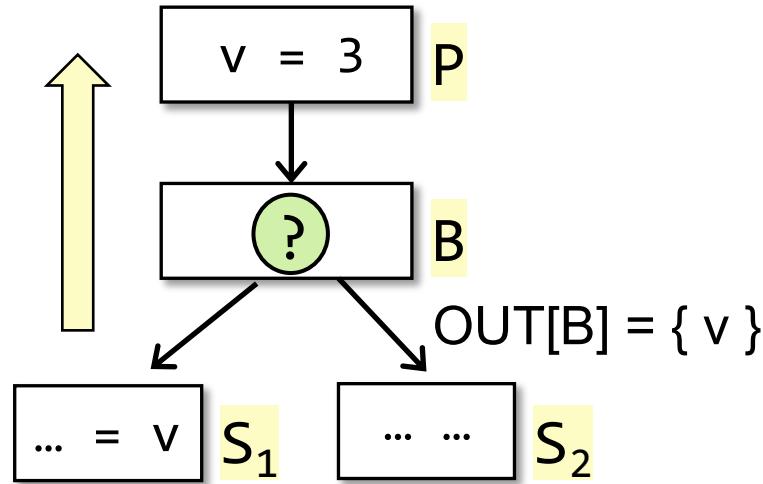


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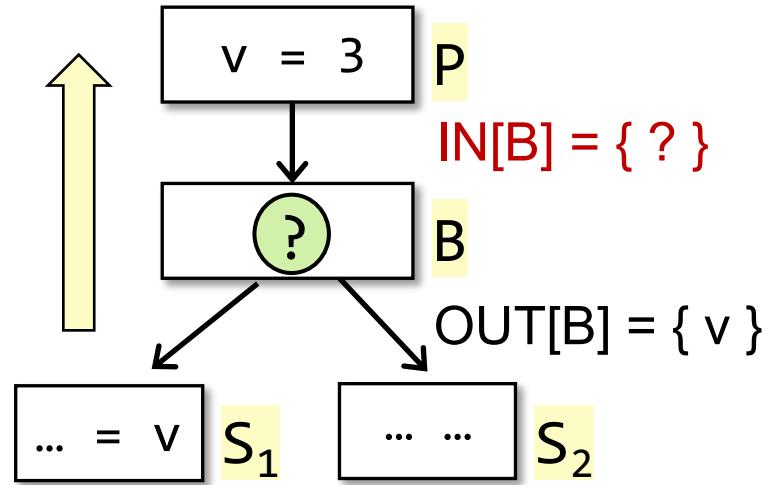
$$\text{OUT}[B] = \bigcup_{S \text{ a successor of } B} \text{IN}[S]$$

# Understanding Live Variables Analysis



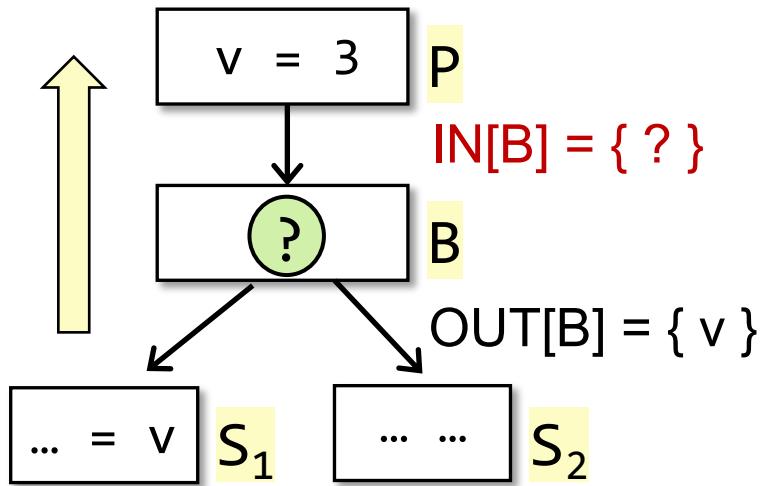
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# Understanding Live Variables Analysis



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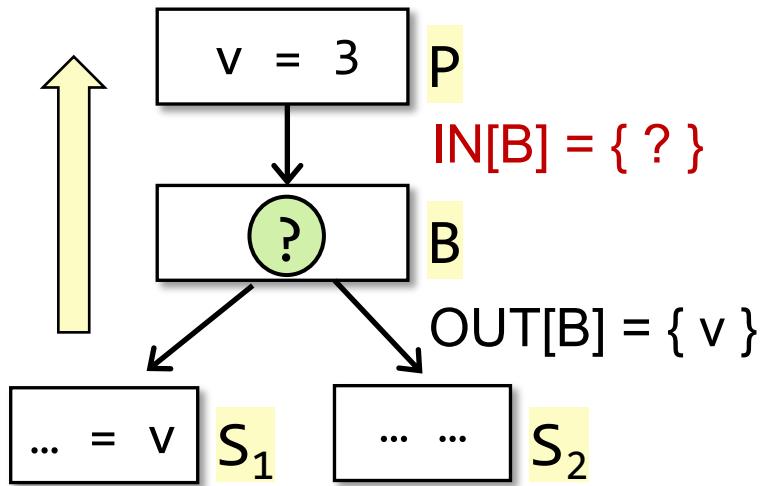


**Tip:** determine **whether** the variable **v** in some register **R** is **live**, or should we delete the value 3 of **v** in **R**, at the point of **IN[B]**?

Yes: **IN[B] = { v }** No: **IN[B] = { }**

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# Understanding Live Variables Analysis



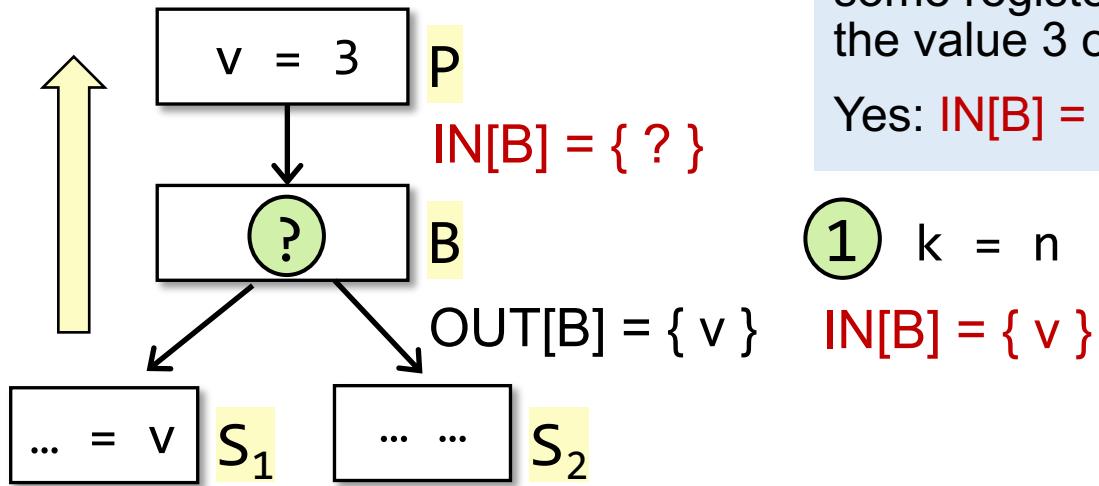
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Yes:  $IN[B] = \{ v \}$  No:  $IN[B] = \{ \}$

1     $k = n$

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

# Understanding Live Variables Analysis

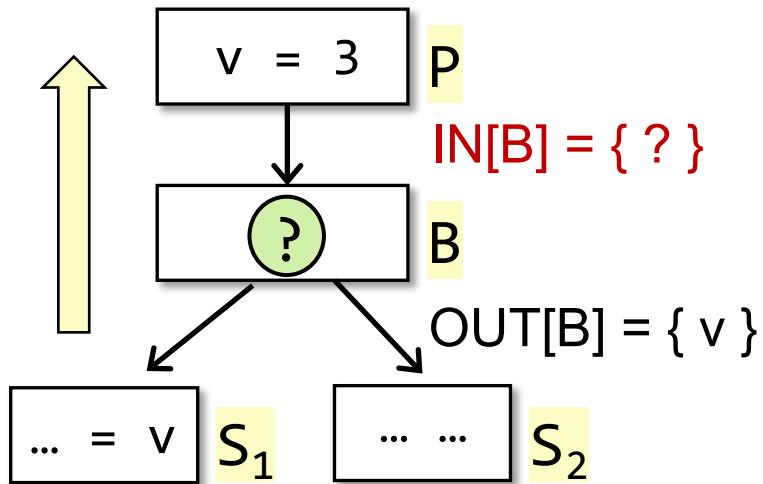


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# Understanding Live Variables Analysis



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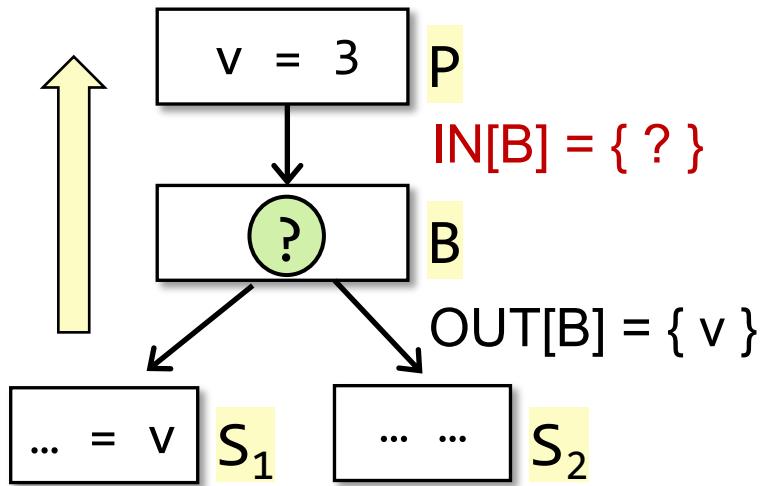
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$$\begin{array}{ll} 1 \quad k = n & 2 \quad k = v \\ \end{array}$$

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# Understanding Live Variables Analysis



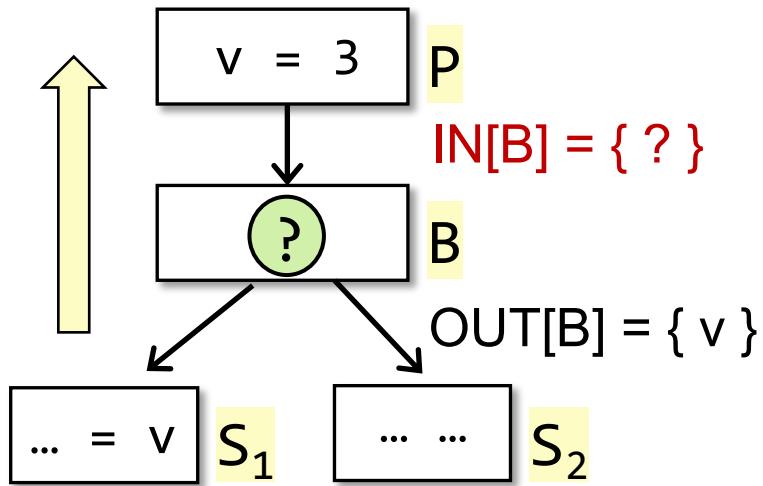
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# Understanding Live Variables Analysis



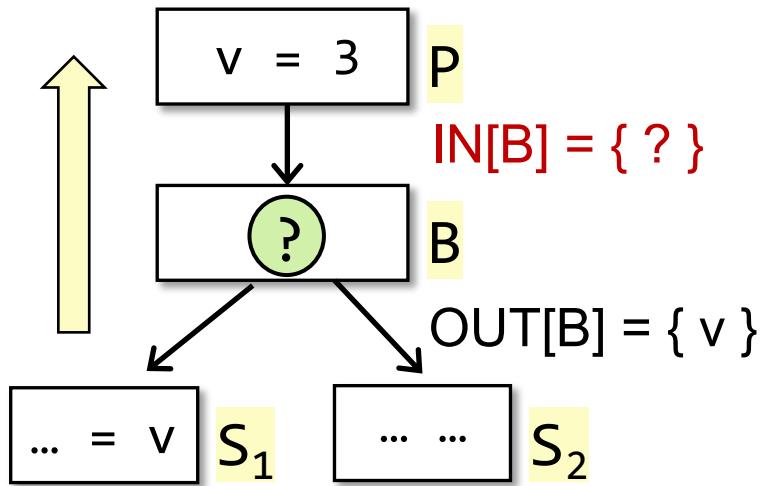
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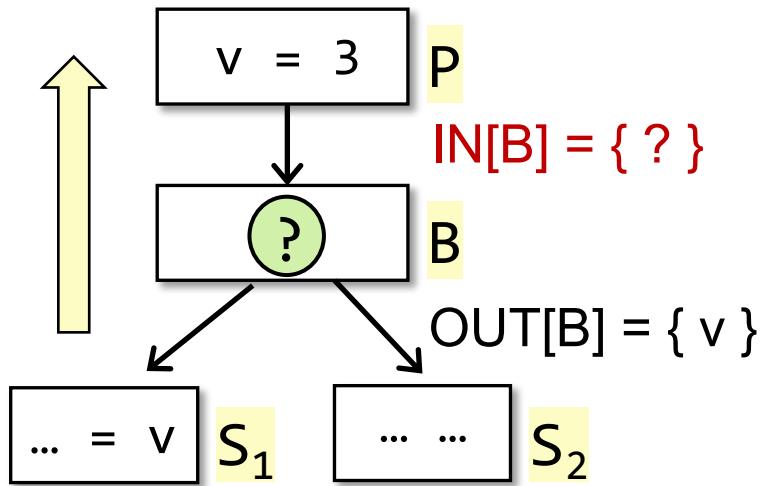
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# Understanding Live Variables Analysis



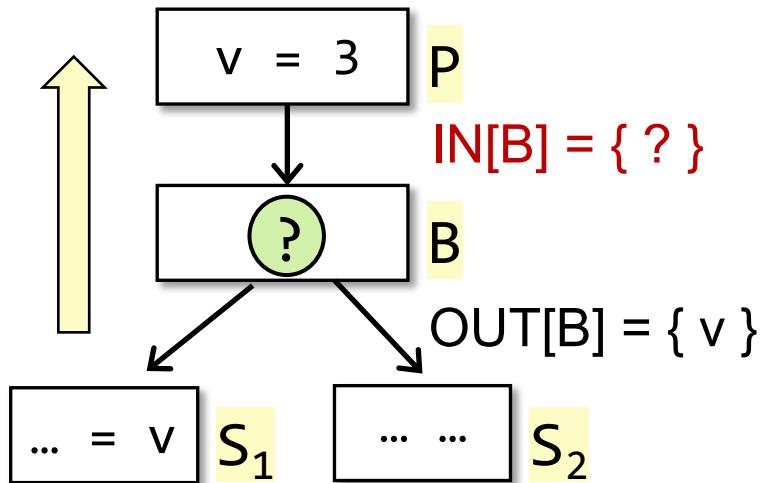
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- (4)  $v = v - 1$

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# Understanding Live Variables Analysis



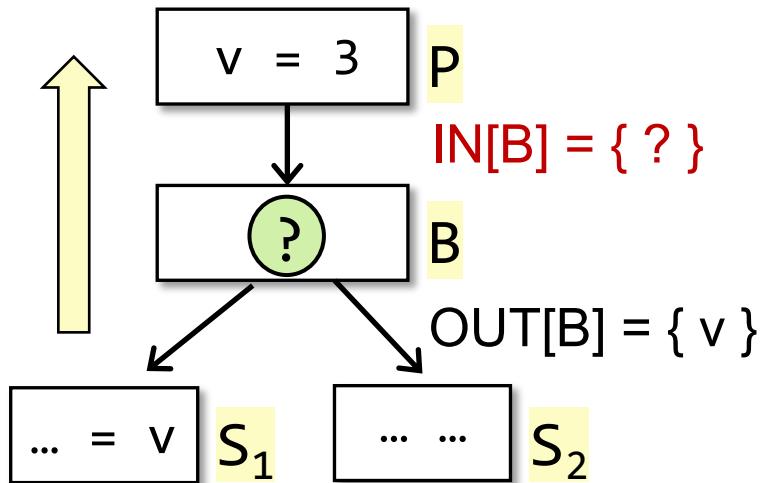
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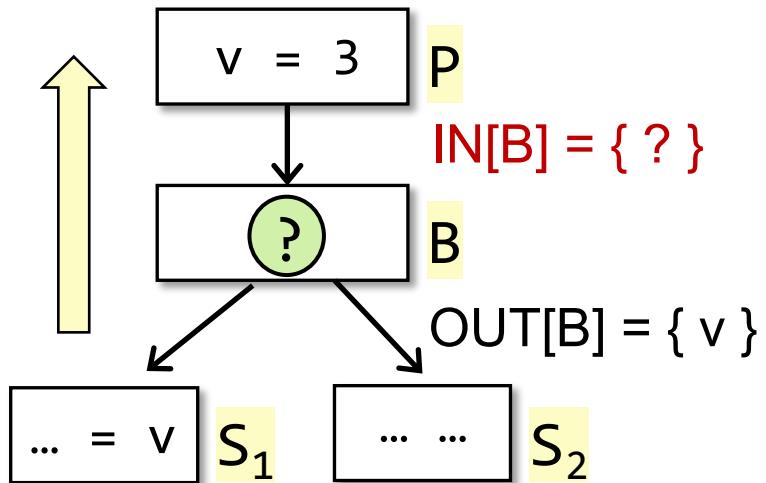
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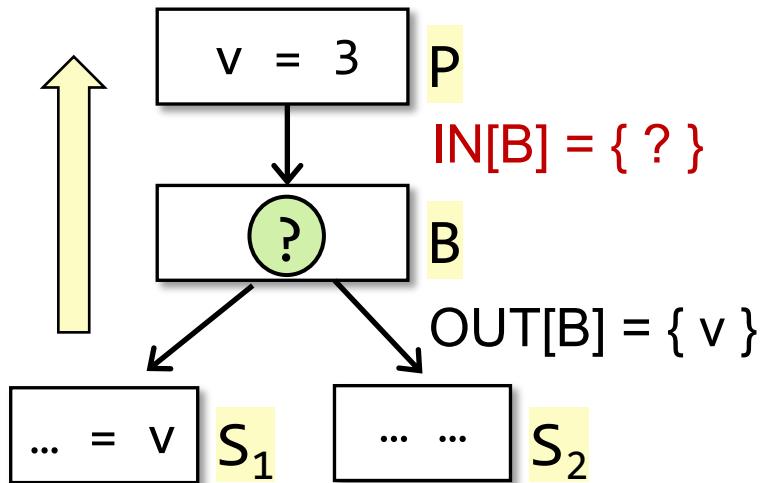
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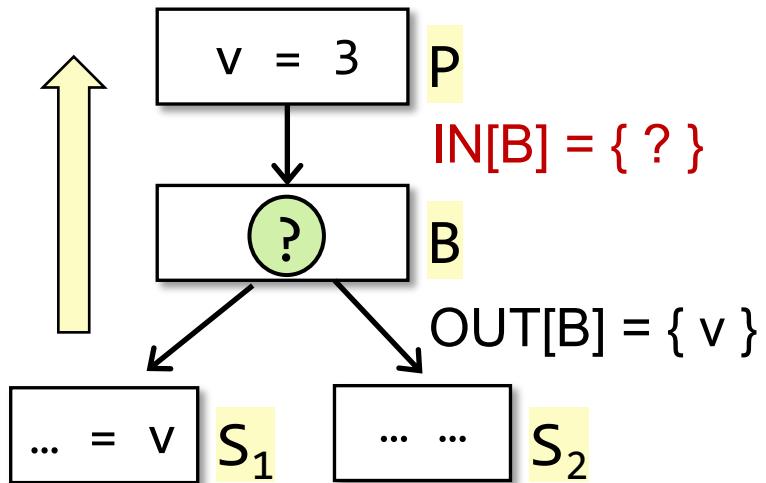
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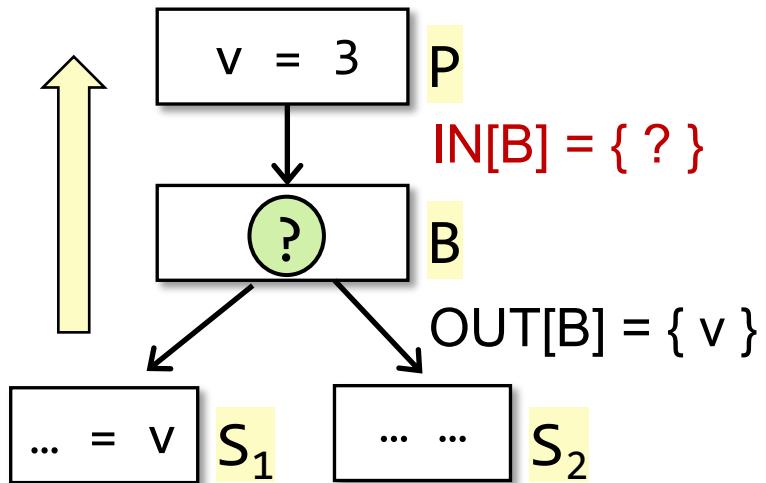
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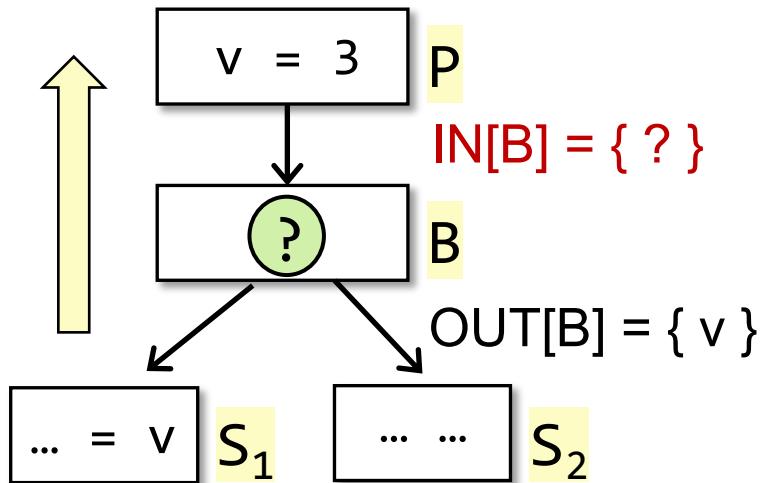
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$$\text{OUT}[B] = \bigcup_{S \text{ a successor of } B} \text{IN}[S]$$

$$\text{IN}[B] = \text{use}_B \cup (\text{OUT}[B] - \text{def}_B)$$

# Understanding Live Variables Analysis



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- |   |                          |   |                          |   |                          |
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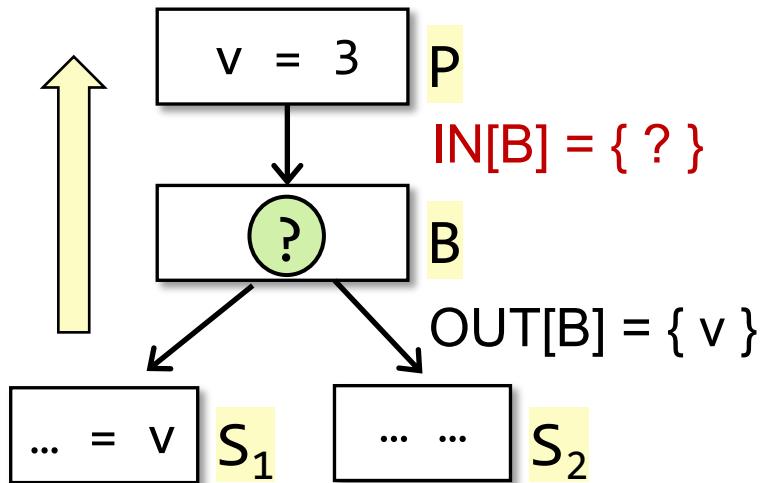
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It is redefined in  $B$

3, 4, 5, 6

# Understanding Live Variables Analysis



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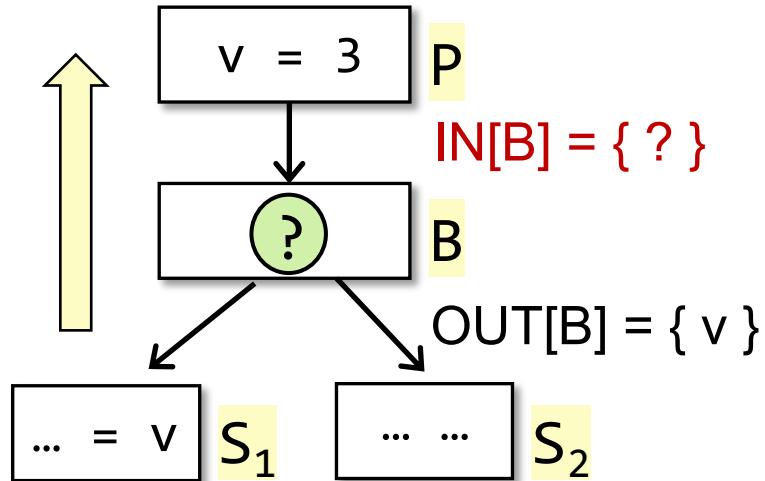
$$IN[B] = use_B \cup (OUT[B] - def_B)$$

It is redefined in B

3, 4, 5, 6

It is live coming out of B and is not redefined in B

# Understanding Live Variables Analysis



**Tip:** determine **whether** the variable  $v$  in some register  $R$  is **live**, or should we delete the value 3 of  $v$  in  $R$ , at the point of  $IN[B]$ ?

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$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

$$IN[B] = use_B \cup (OUT[B] - def_B)$$

It is redefined in B

3, 4, 5, 6

It is used before redefinition in B

It is live coming out of B and is not redefined in B

4, 6

# Algorithm of Live Variables Analysis

**INPUT:** CFG ( $def_B$  and  $use_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

```
IN[exit] = ∅;  
for (each basic block  $B \setminus exit$ )  
    IN[ $B$ ] = ∅;  
    while (changes to any IN occur)  
        for (each basic block  $B \setminus exit$ ) {  
            OUT[ $B$ ] =  $\bigcup_{S \text{ a successor of } B} IN[S]$ ;  
            IN[ $B$ ] =  $use_B \cup (OUT[B] - def_B)$ ;  
        }  
    }
```

# Algorithm of Live Variables Analysis

**INPUT:** CFG ( $def_B$  and  $use_B$  computed for each basic block  $B$ )

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# Algorithm of Live Variables Analysis

**INPUT:** CFG ( $def_B$  and  $use_B$  computed for each basic block  $B$ )

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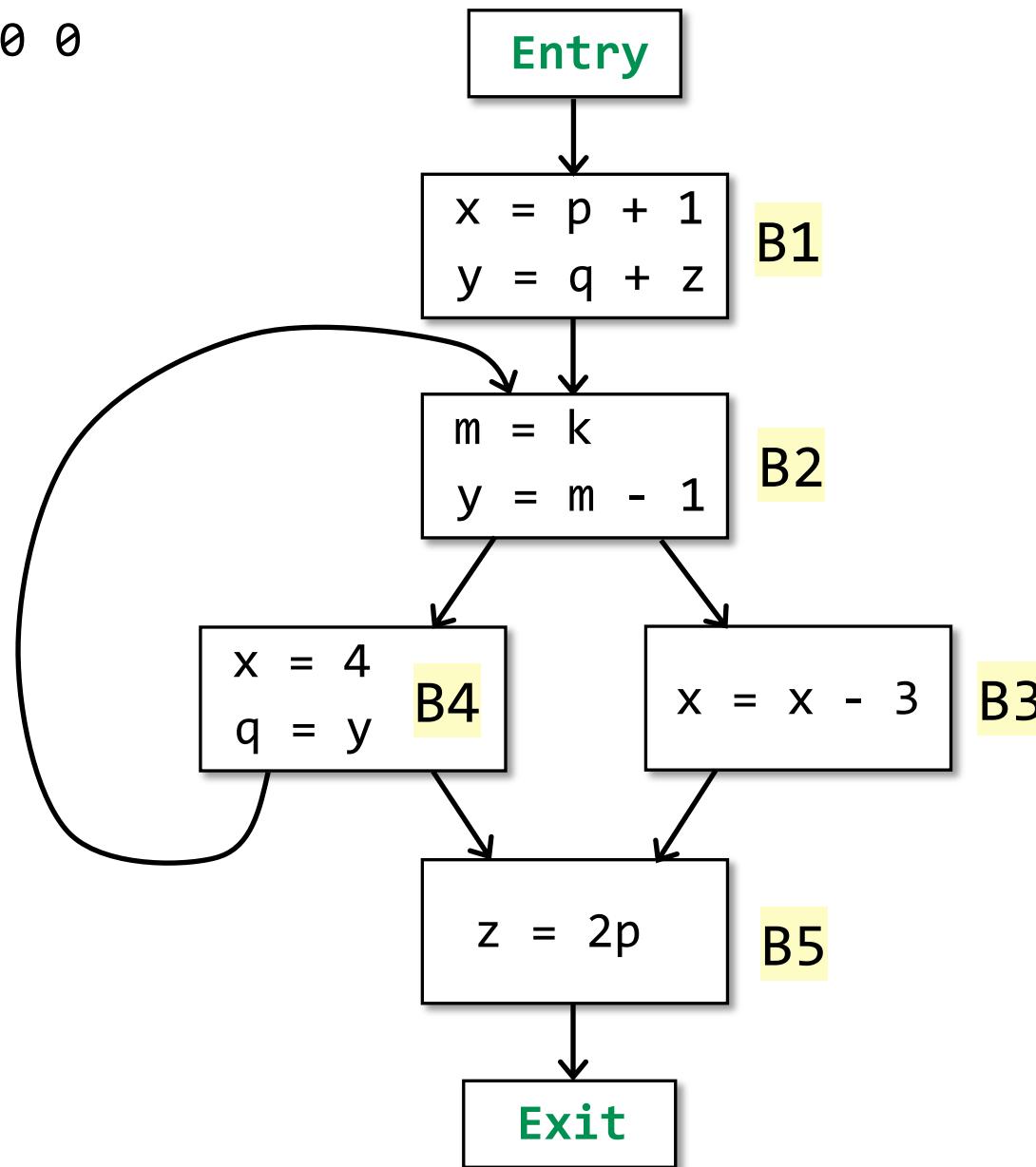
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```

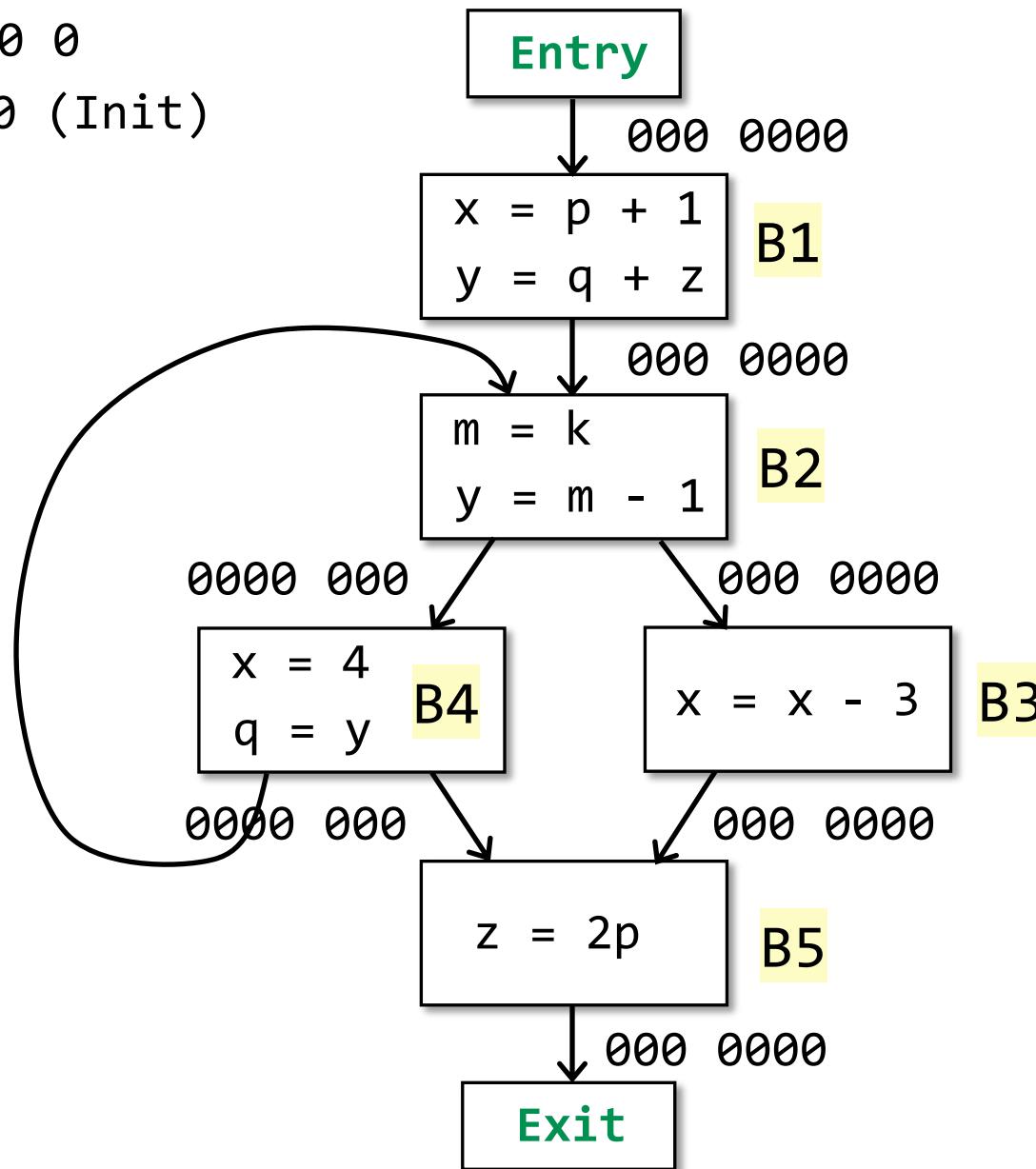


| x | y | z | p | q | m | k |
|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 |



x y z p q m k  
0 0 0 0 0 0 0

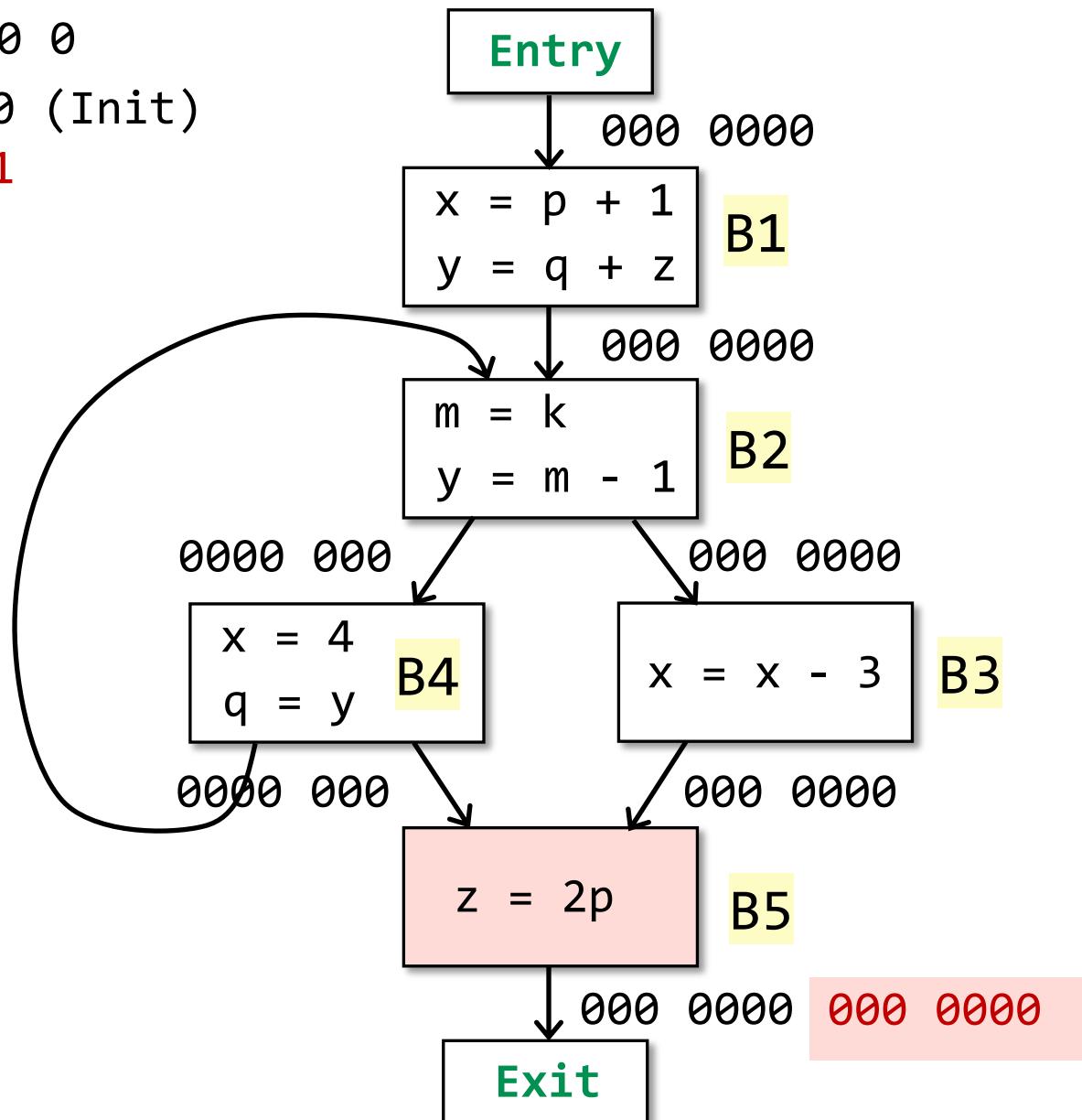
Iteration 0 (Init)



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

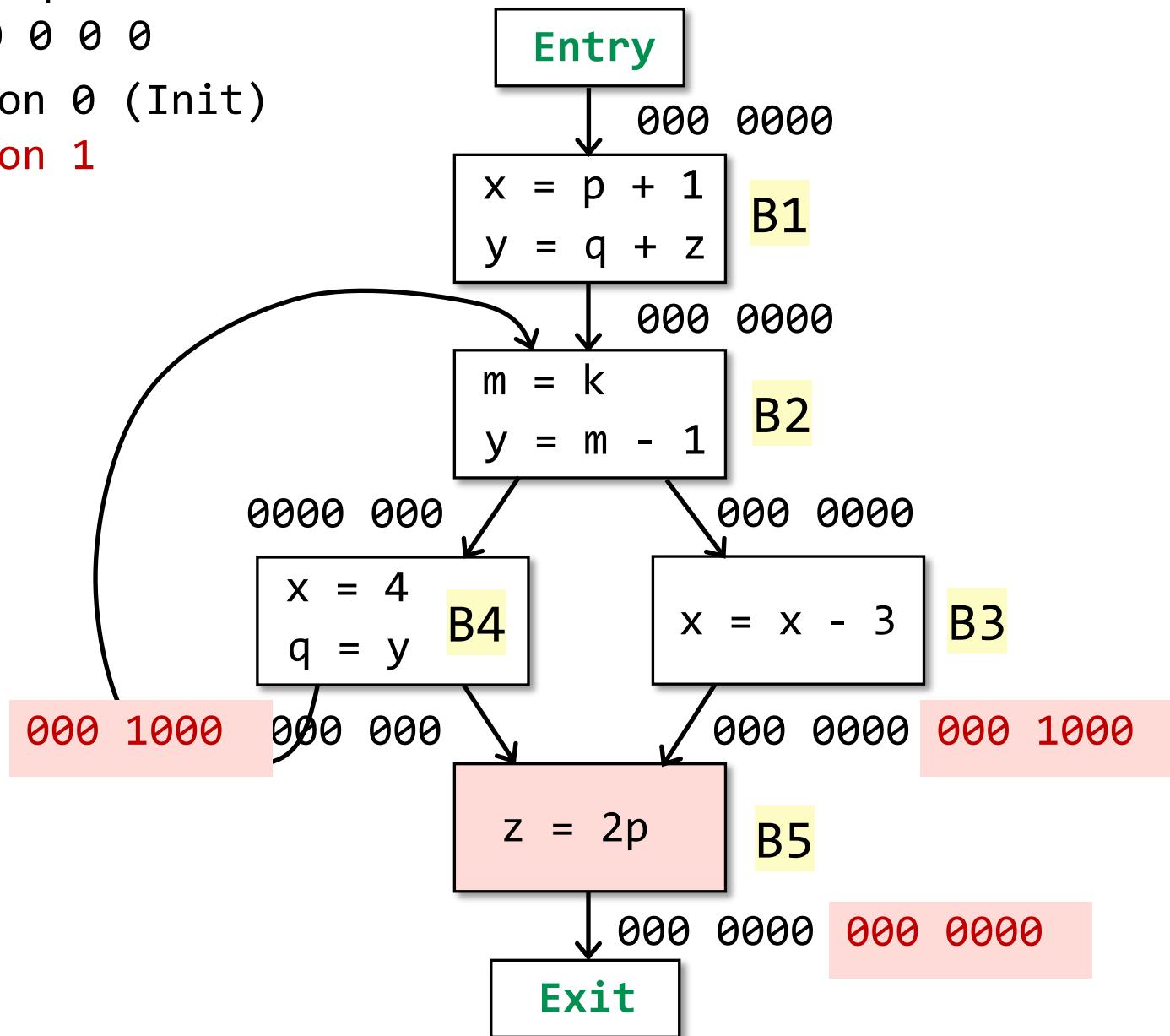


$$IN[B] = use_B \cup (OUT[B] - def_B)$$

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

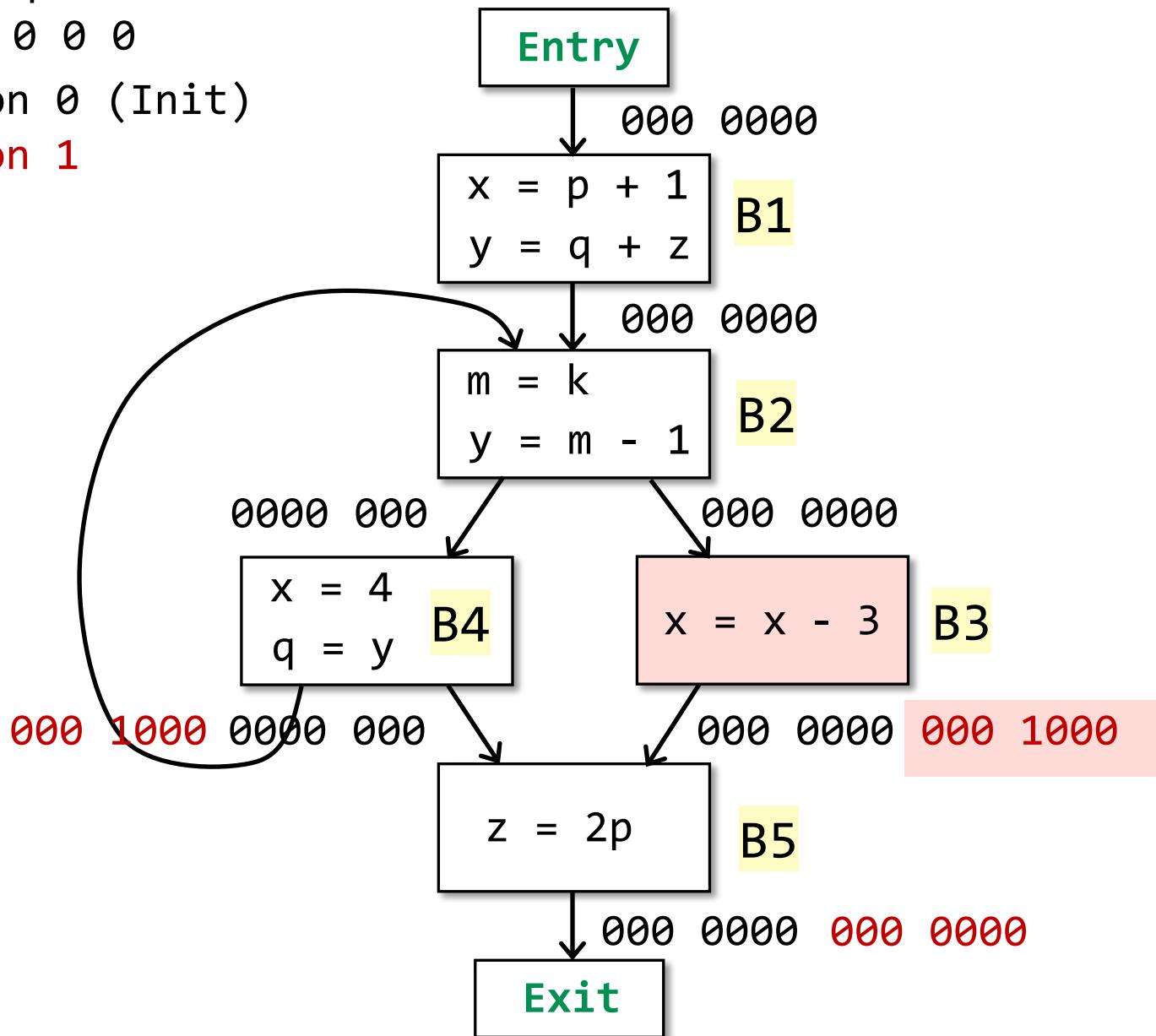
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

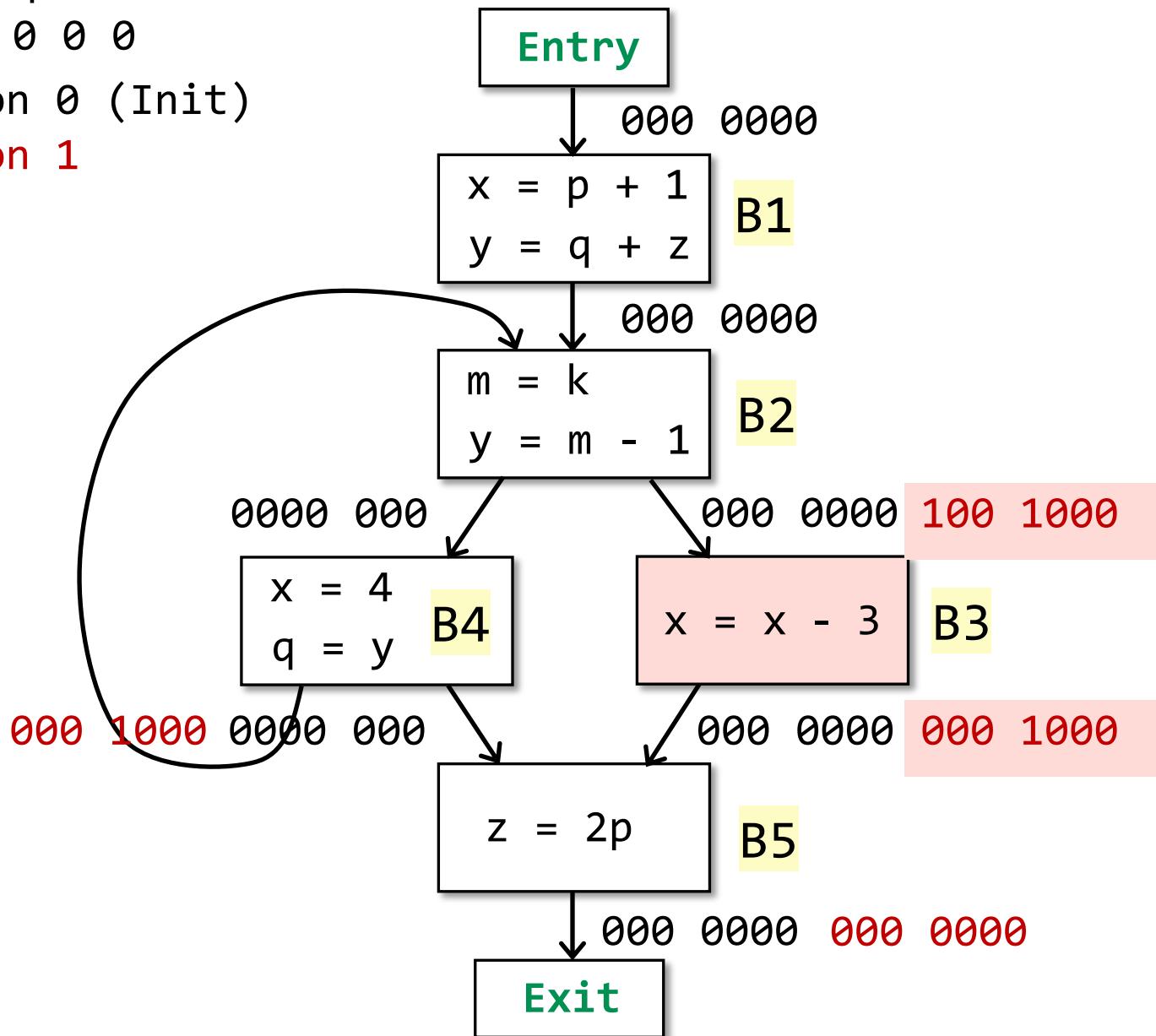
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

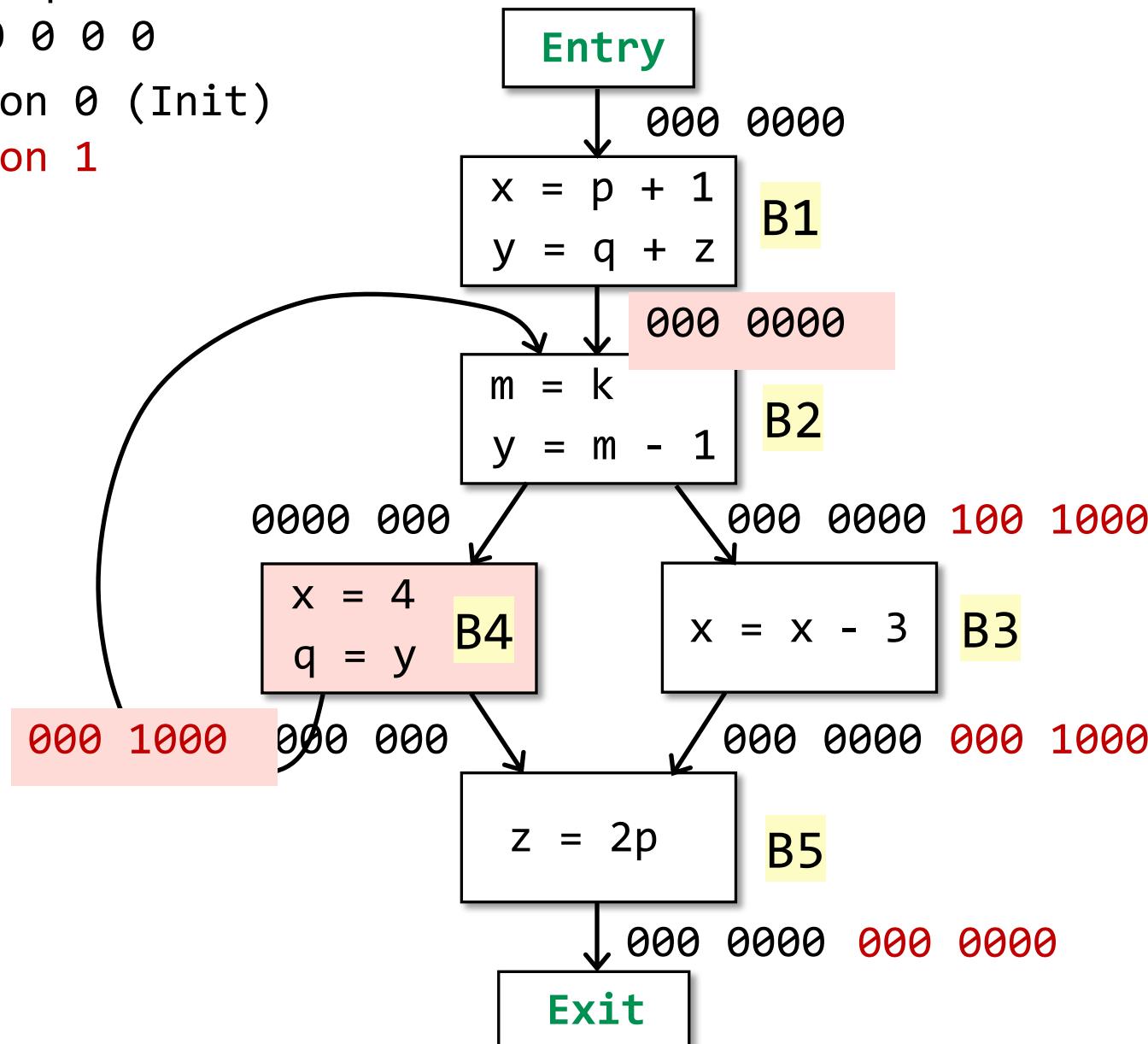
Iteration 1



x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

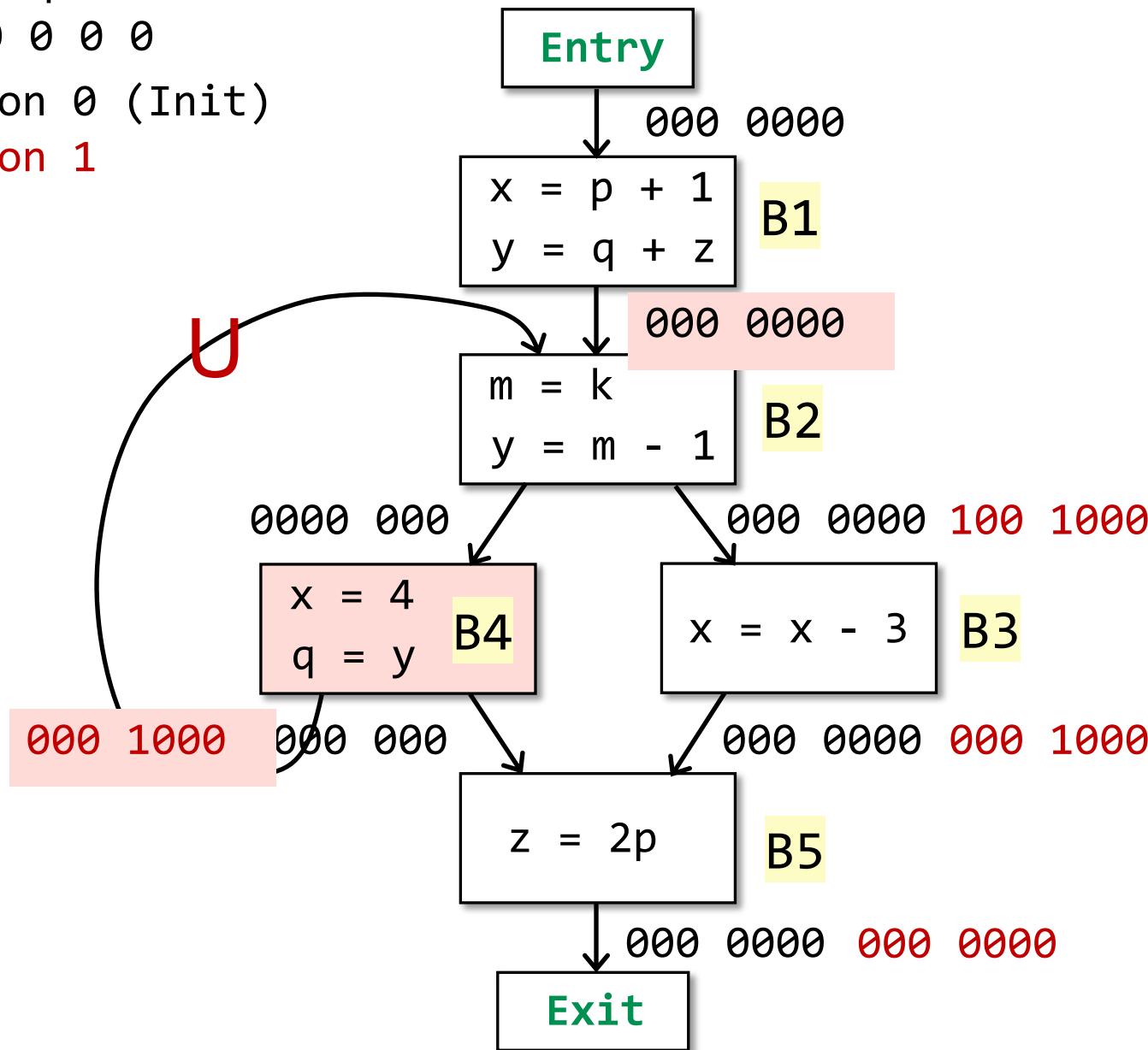


$$\text{OUT}[B] = \bigcup_{S \text{ a successor of } B} \text{IN}[B]$$

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

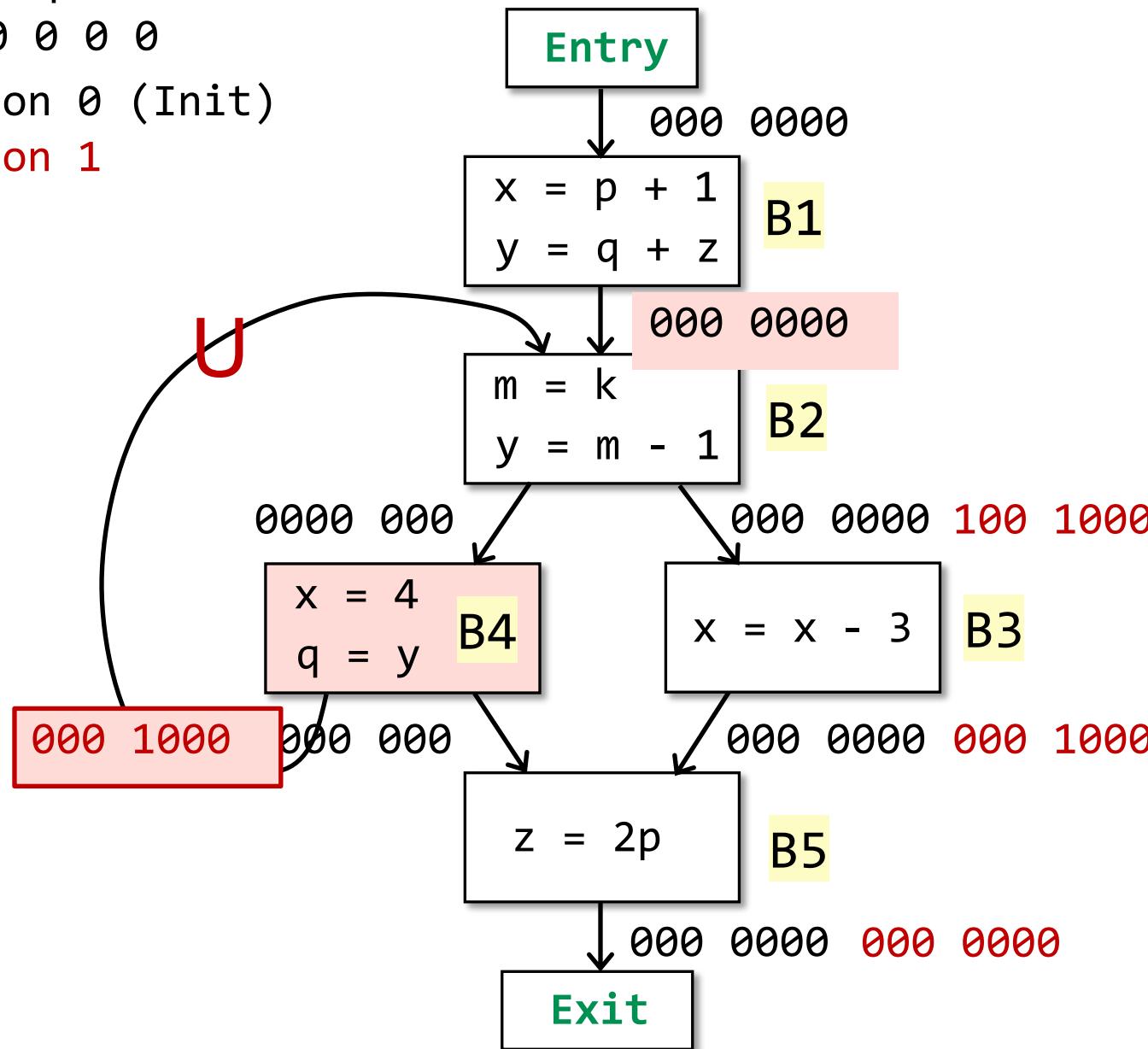
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

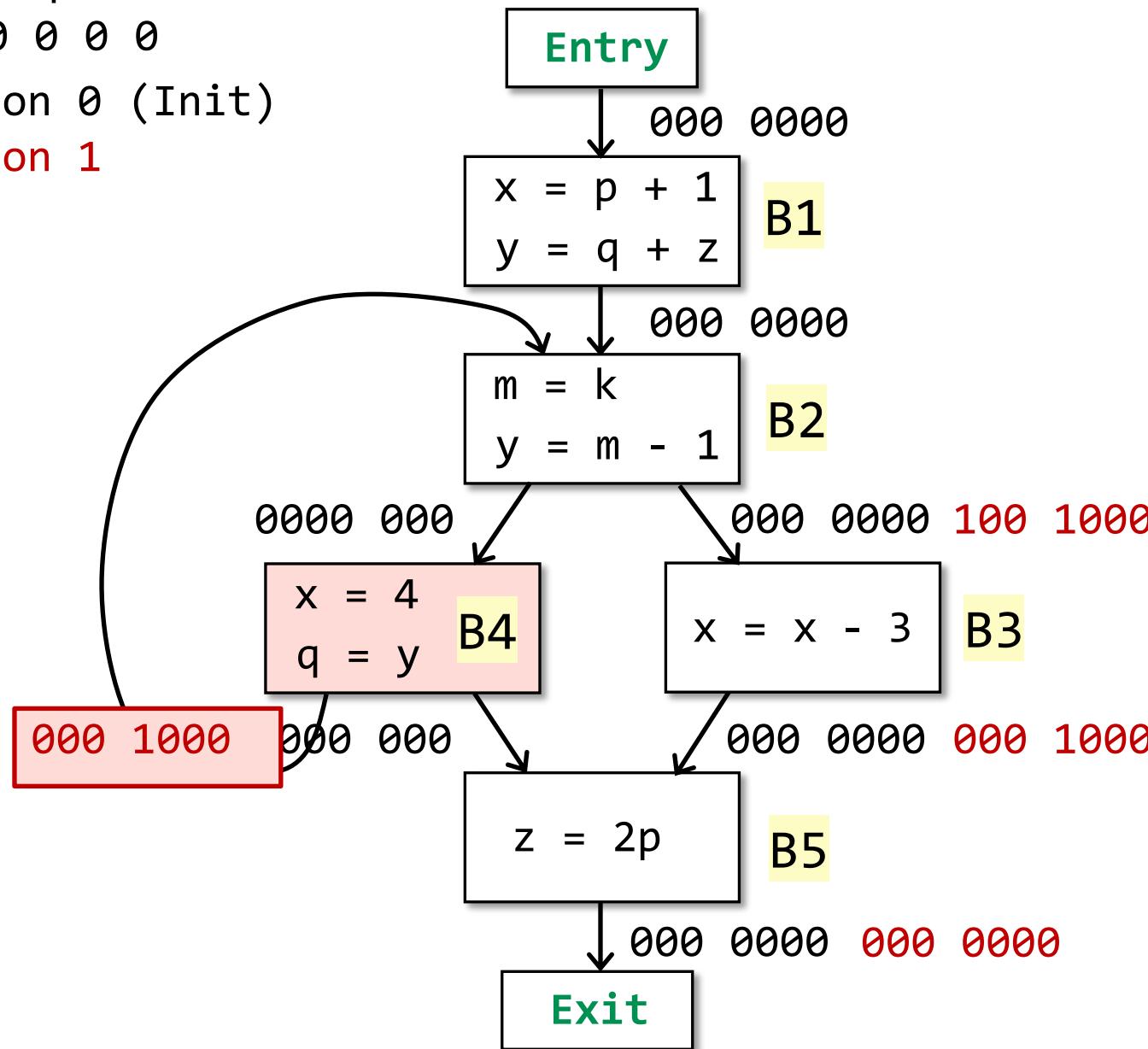
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

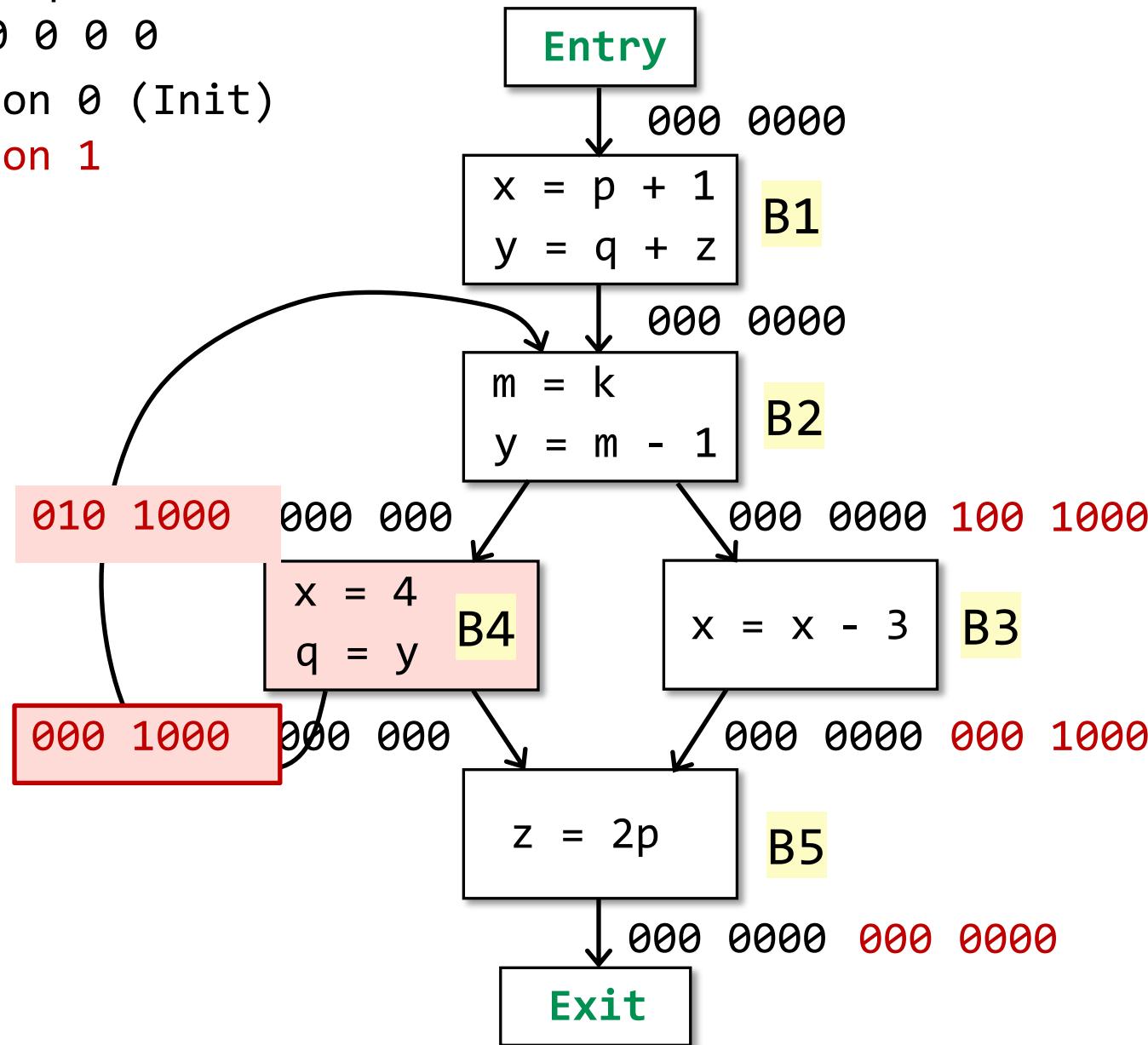
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

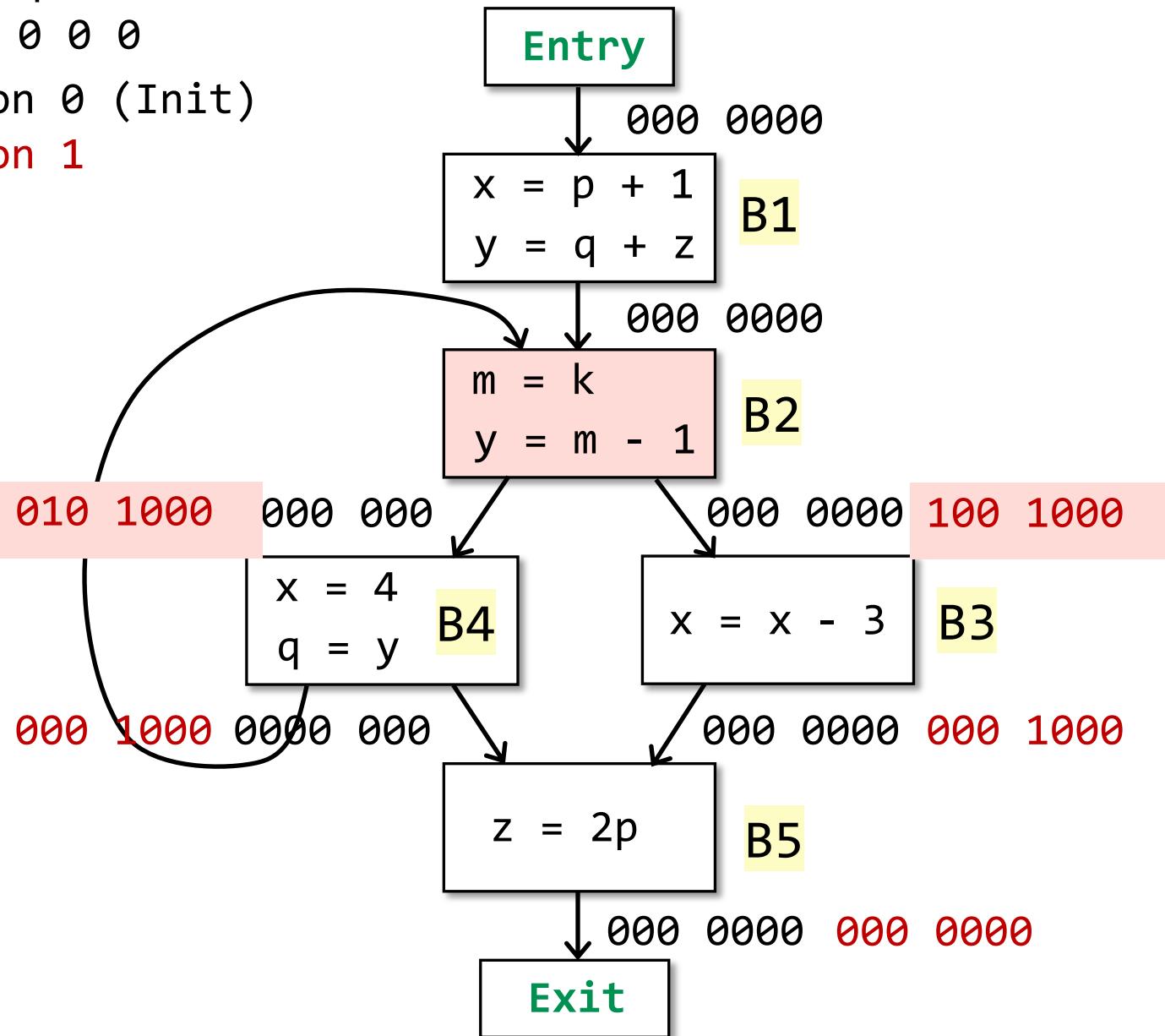
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

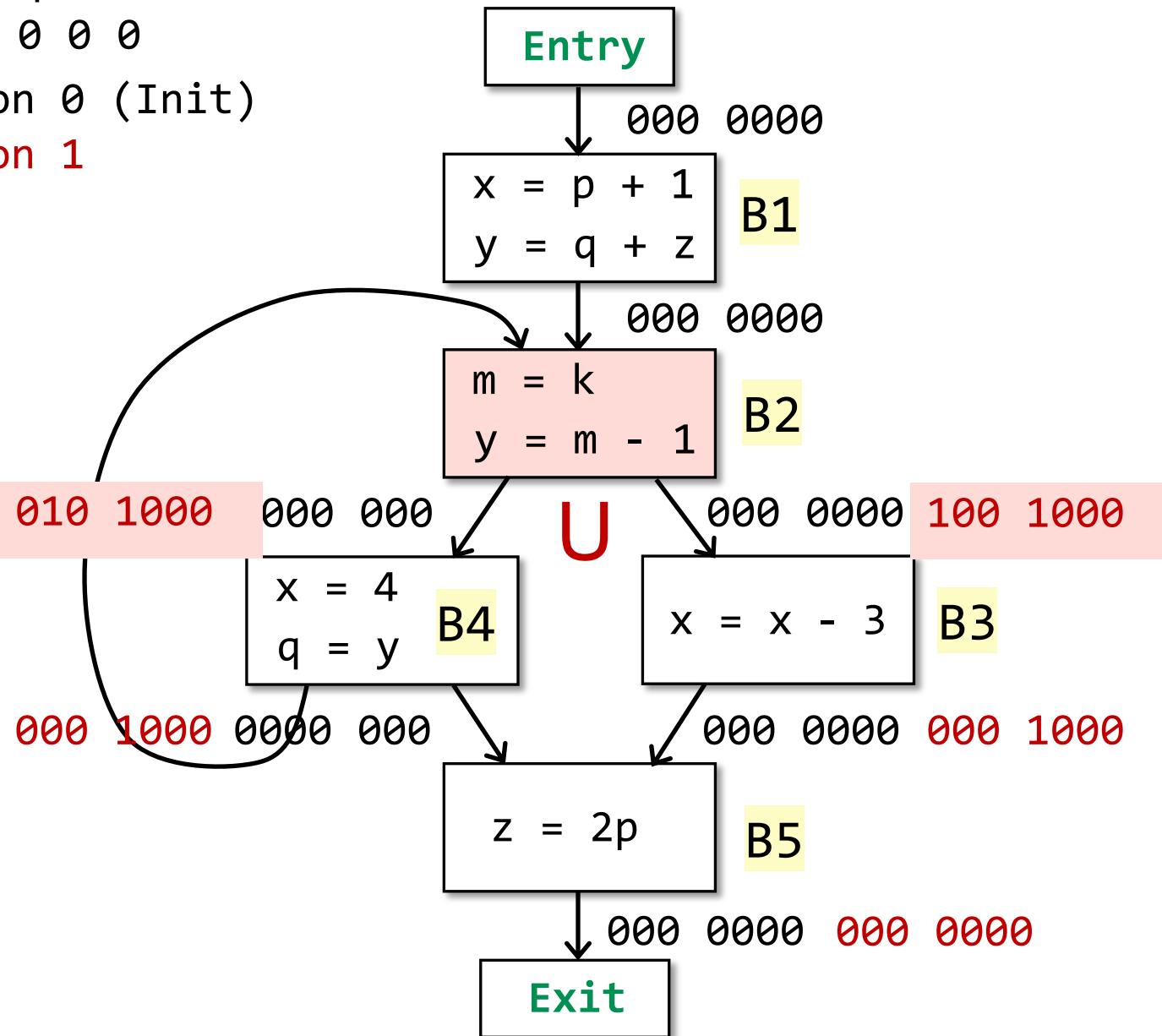
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

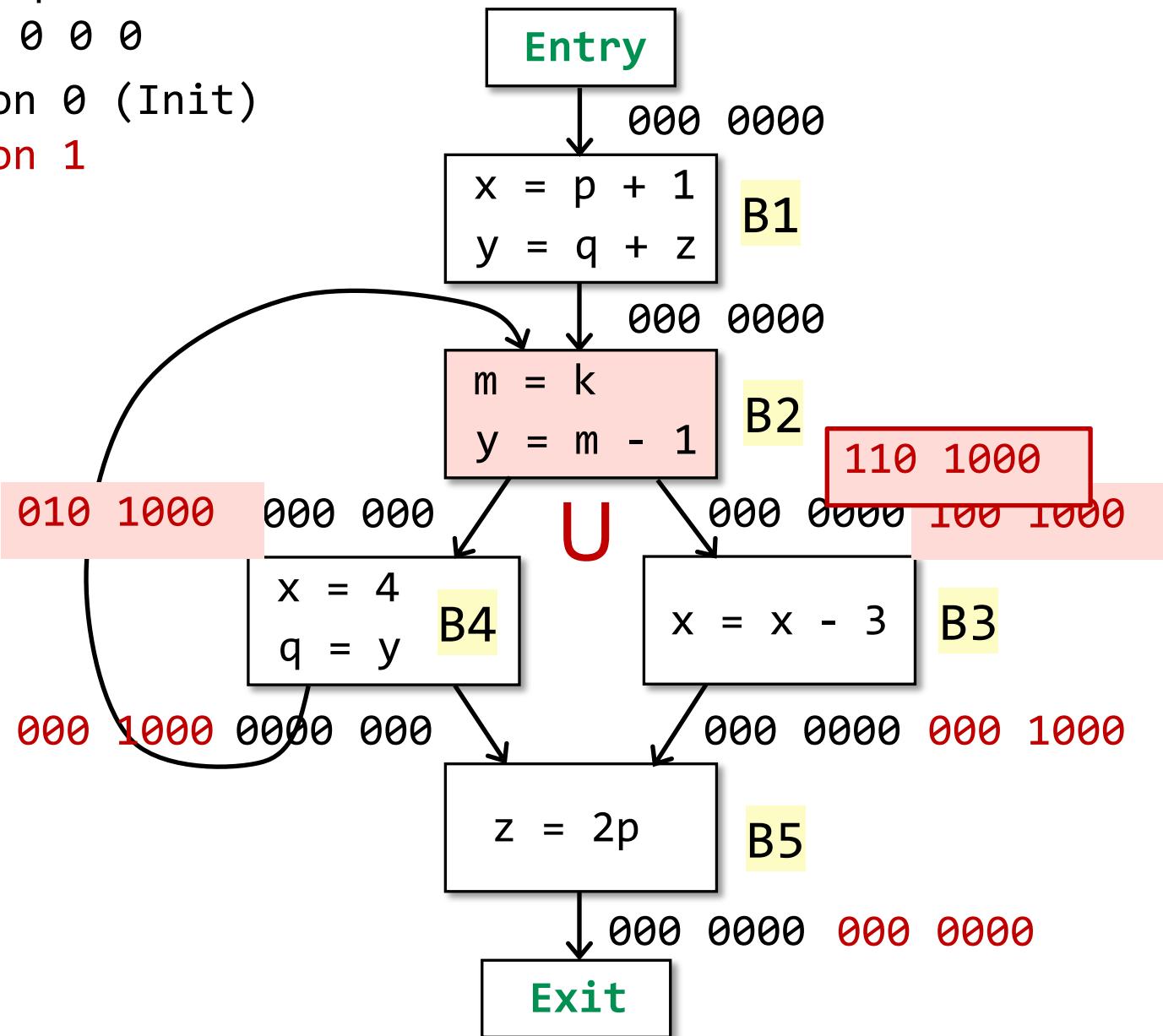
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

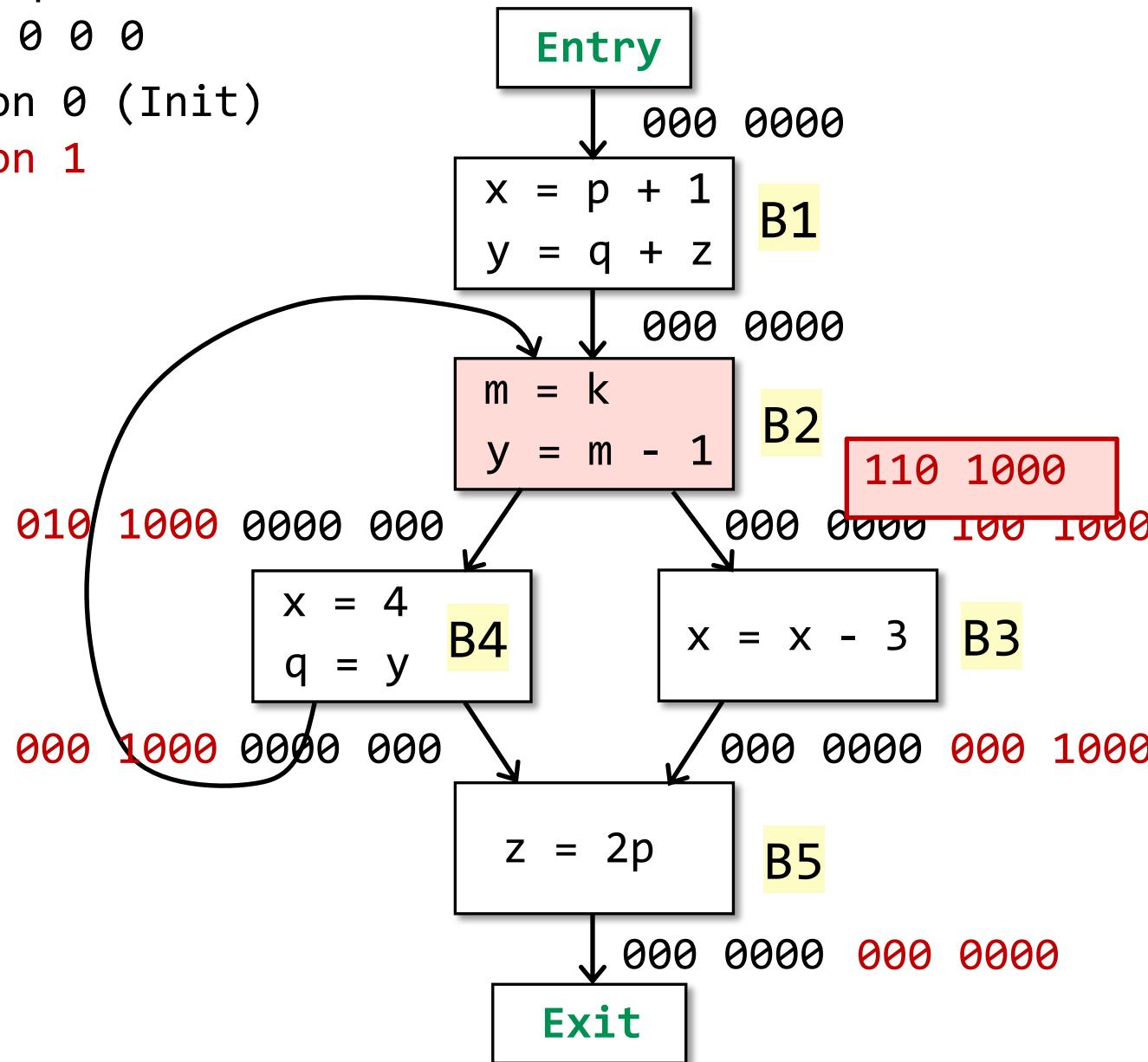
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

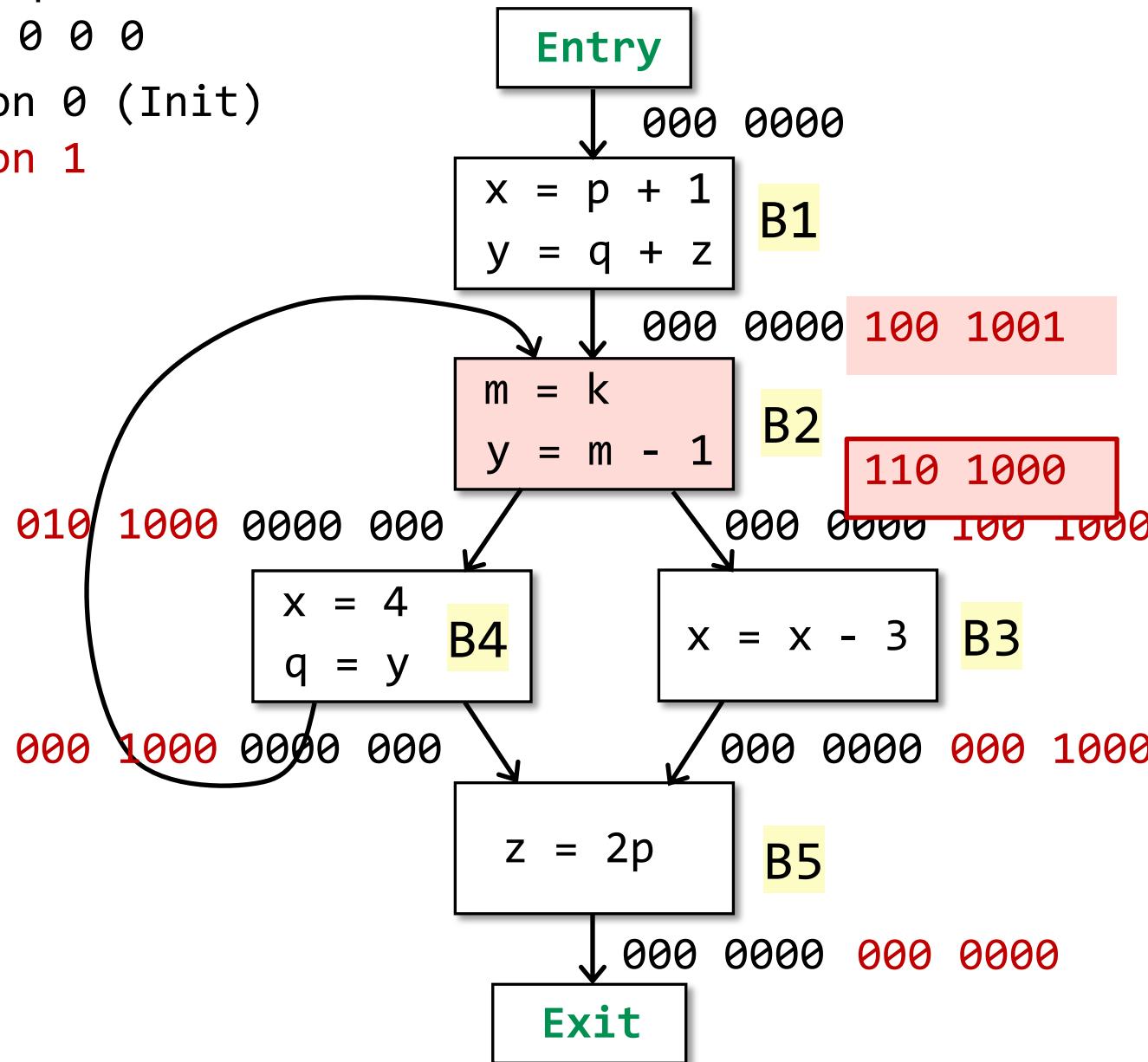
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

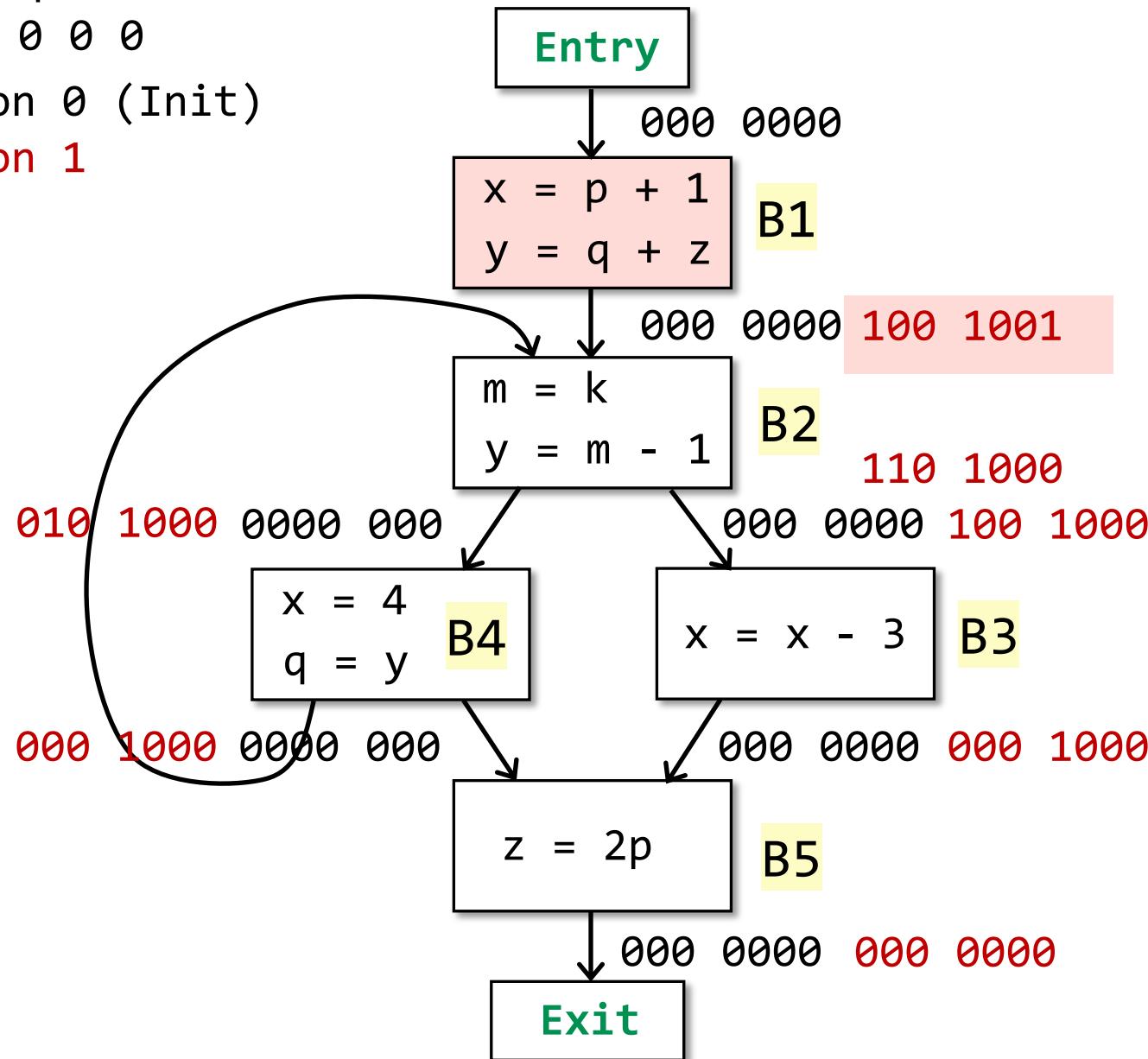
Iteration 1



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

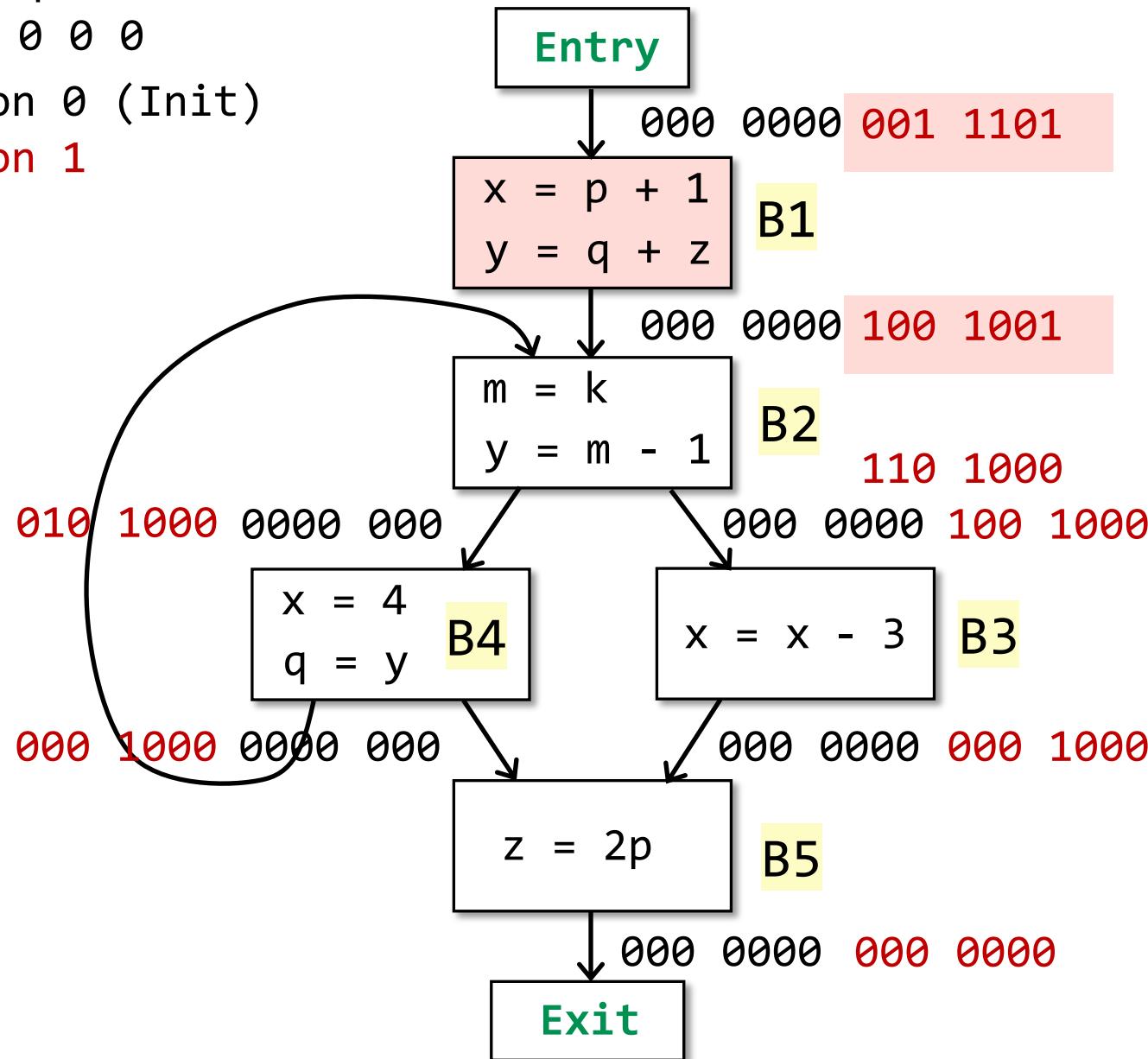
Iteration 1



x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

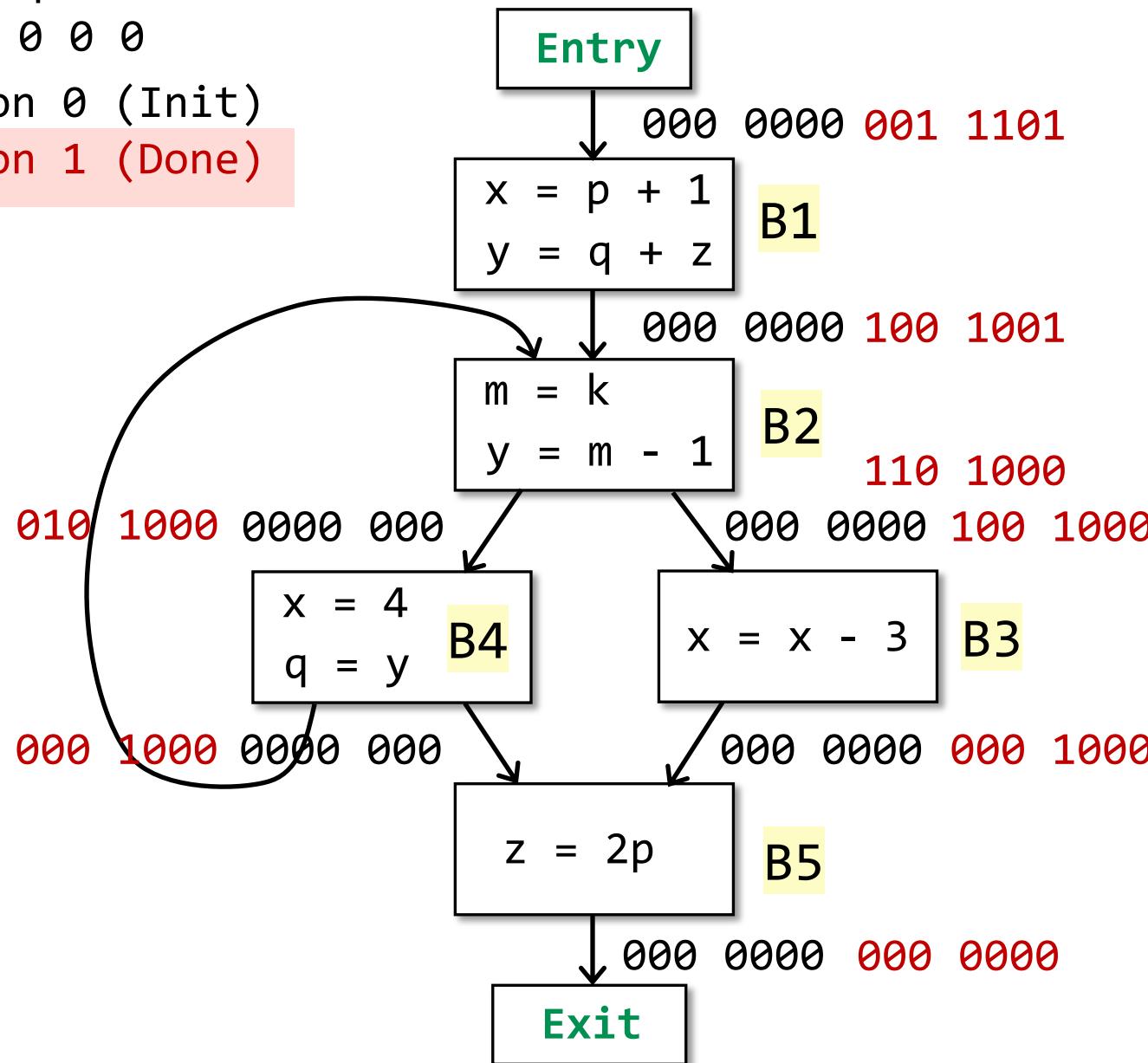
Iteration 1



|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| x | y | z | p | q | m | k |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Iteration 0 (Init)

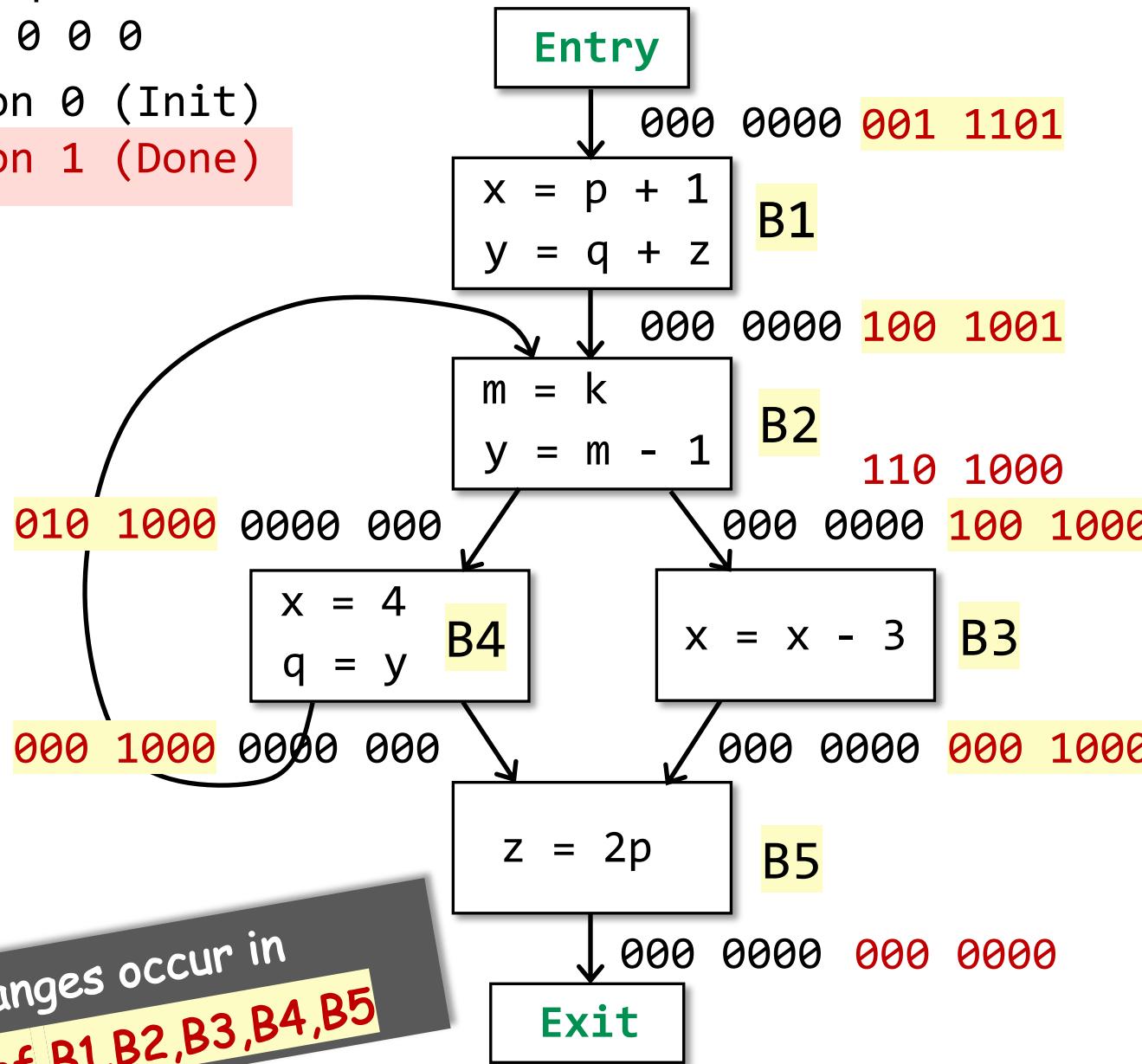
Iteration 1 (Done)



x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)



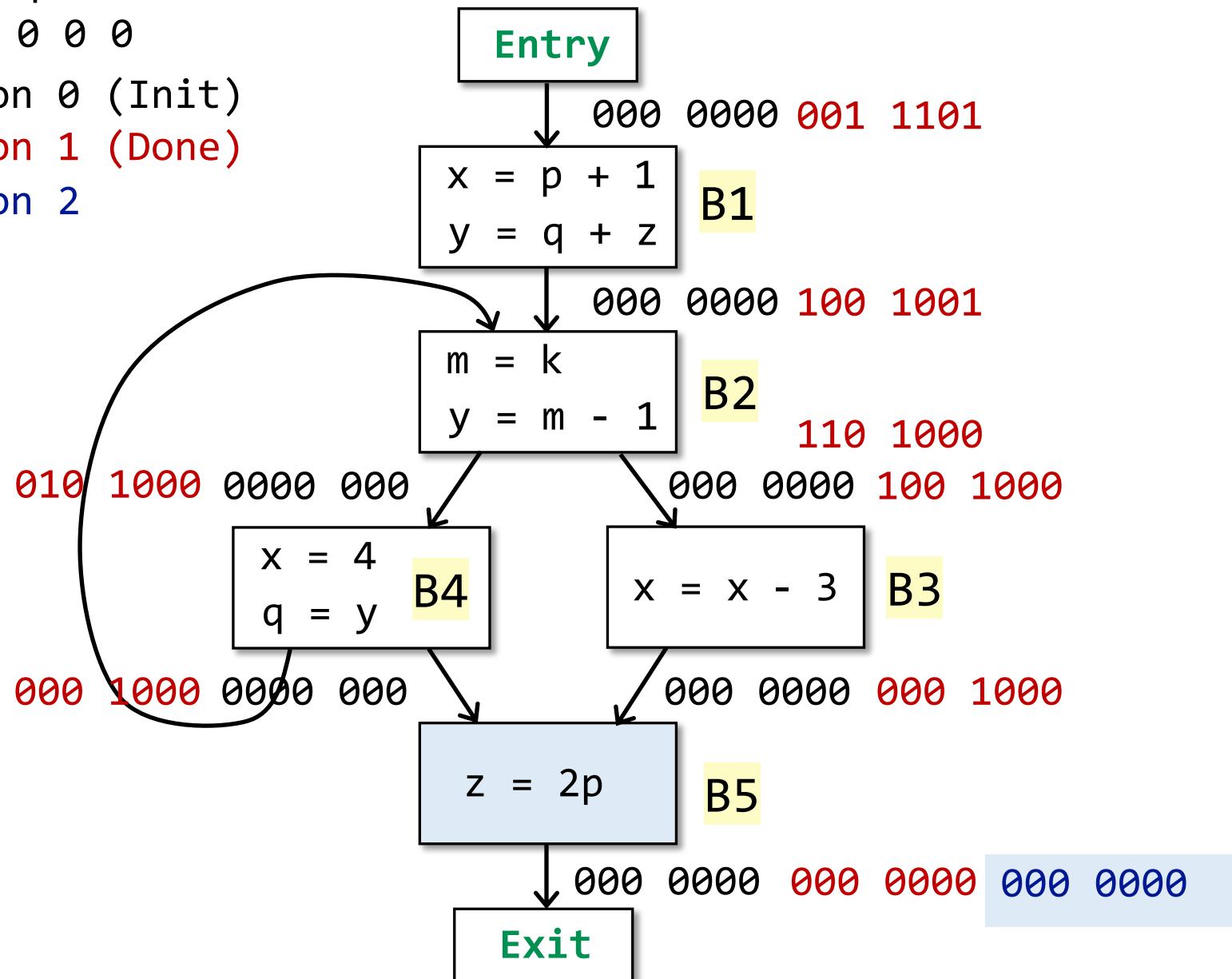
Changes occur in  
 IN[] of B1,B2,B3,B4,B5

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| x | y | z | p | q | m | k |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



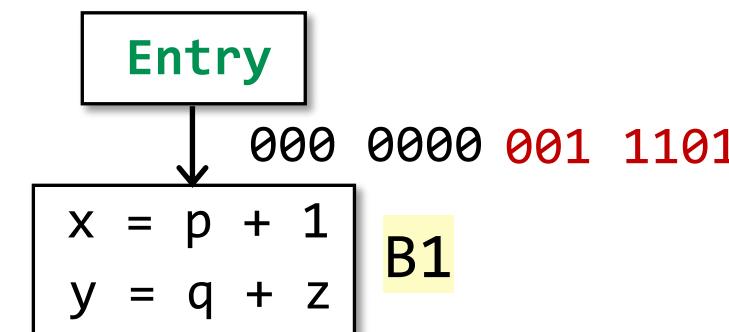
x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



000 0000 001 1101

B1

000 0000 100 1001

B2

110 1000

000 0000 100 1000

B3

010 1000 0000 000

x = 4  
q = y B4

x = x - 3

000 1000 000 1000 0000 000

000 0000 000 1000 000 1000

000 0000 000 0000 000 0000

000 0000 000 0000 000 0000

000 0000 000 0000 000 0000

000 0000 000 0000 000 0000

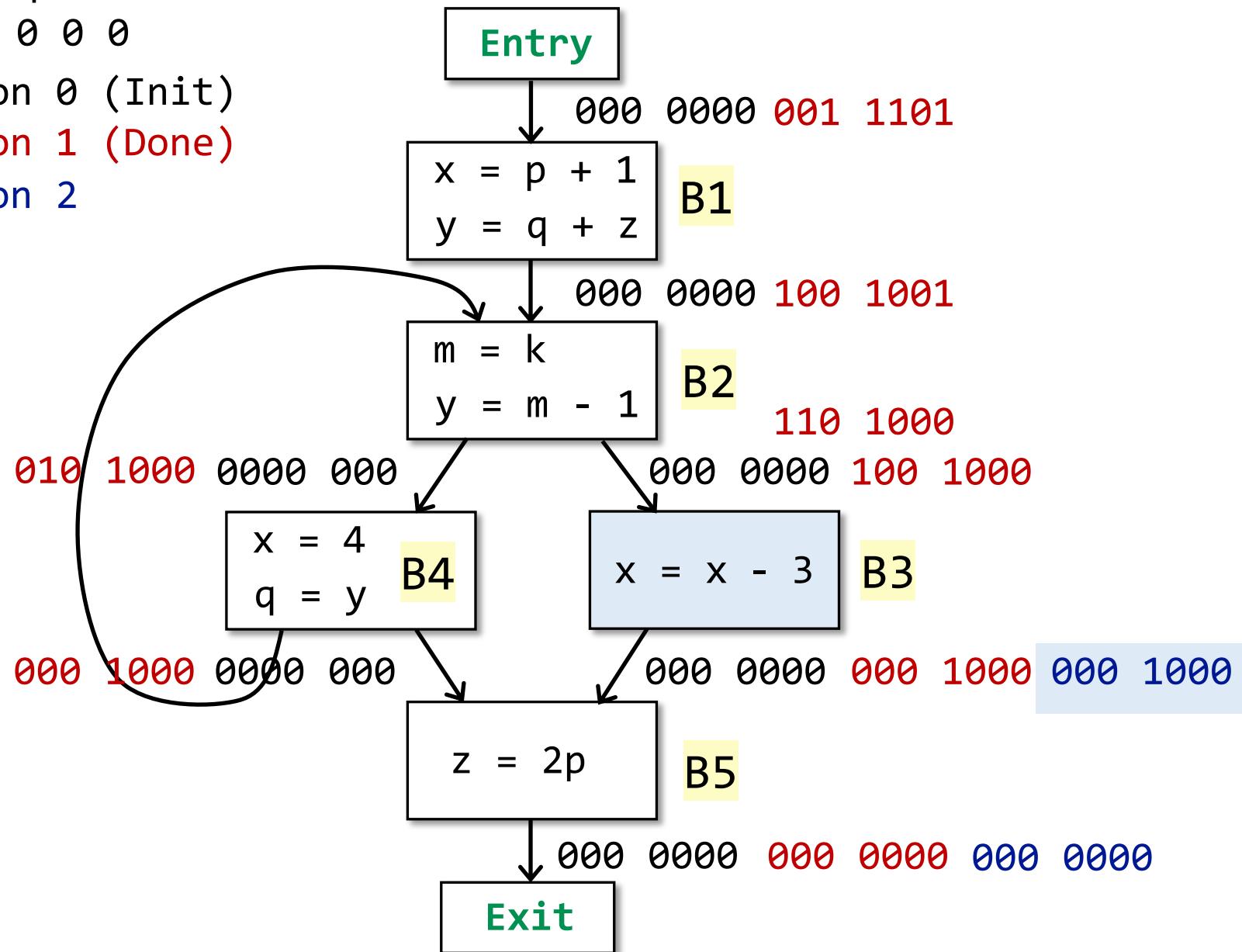
B5

x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

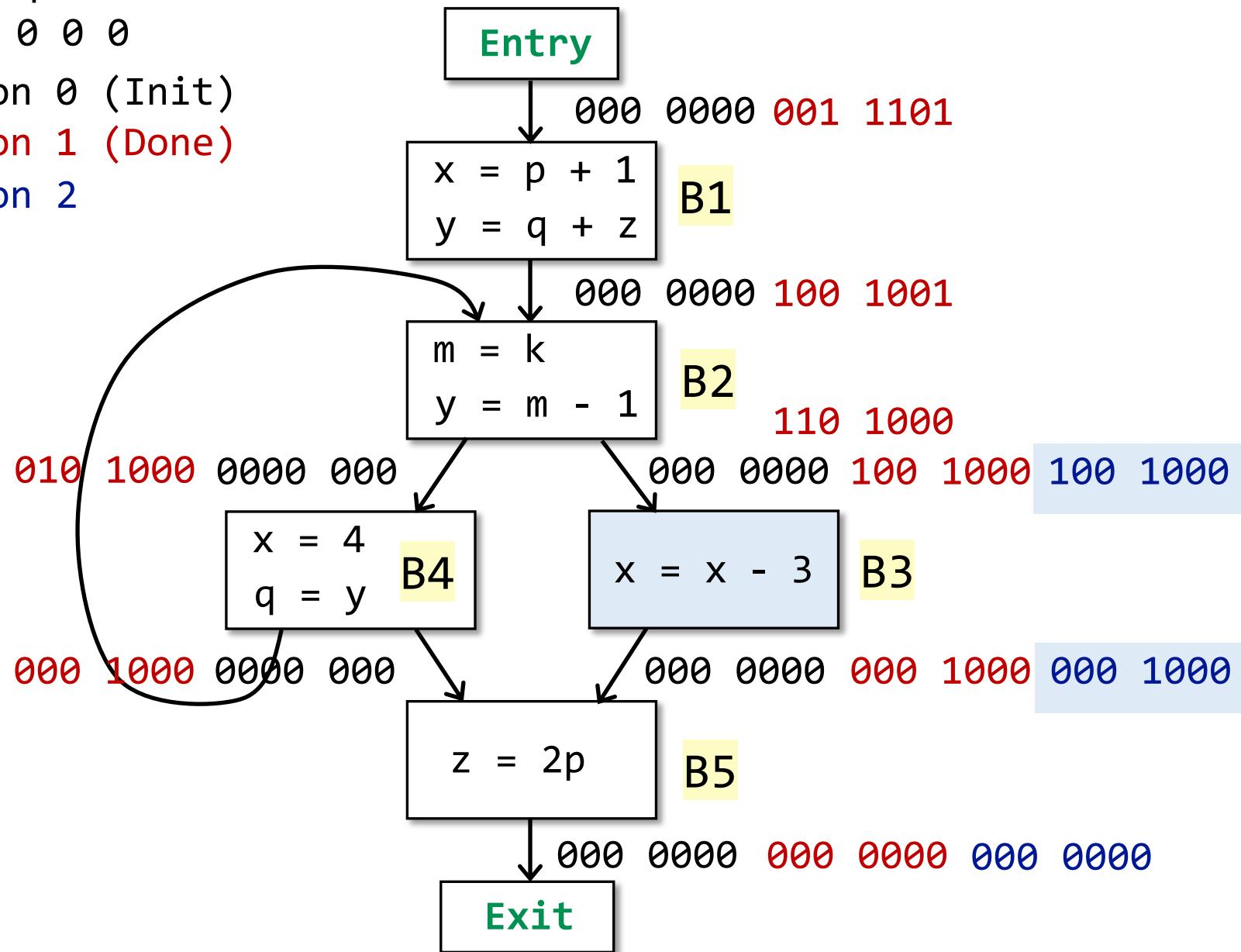


x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

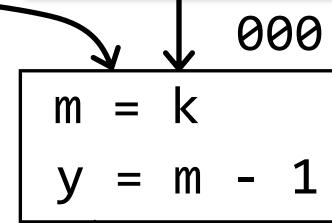
Iteration 2

Entry



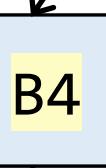
x = p + 1  
y = q + z

B1



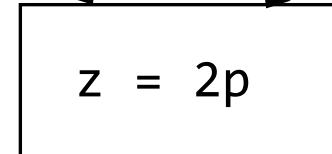
m = k  
y = m - 1

B2



x = 4  
q = y

B4



z = 2p

B5

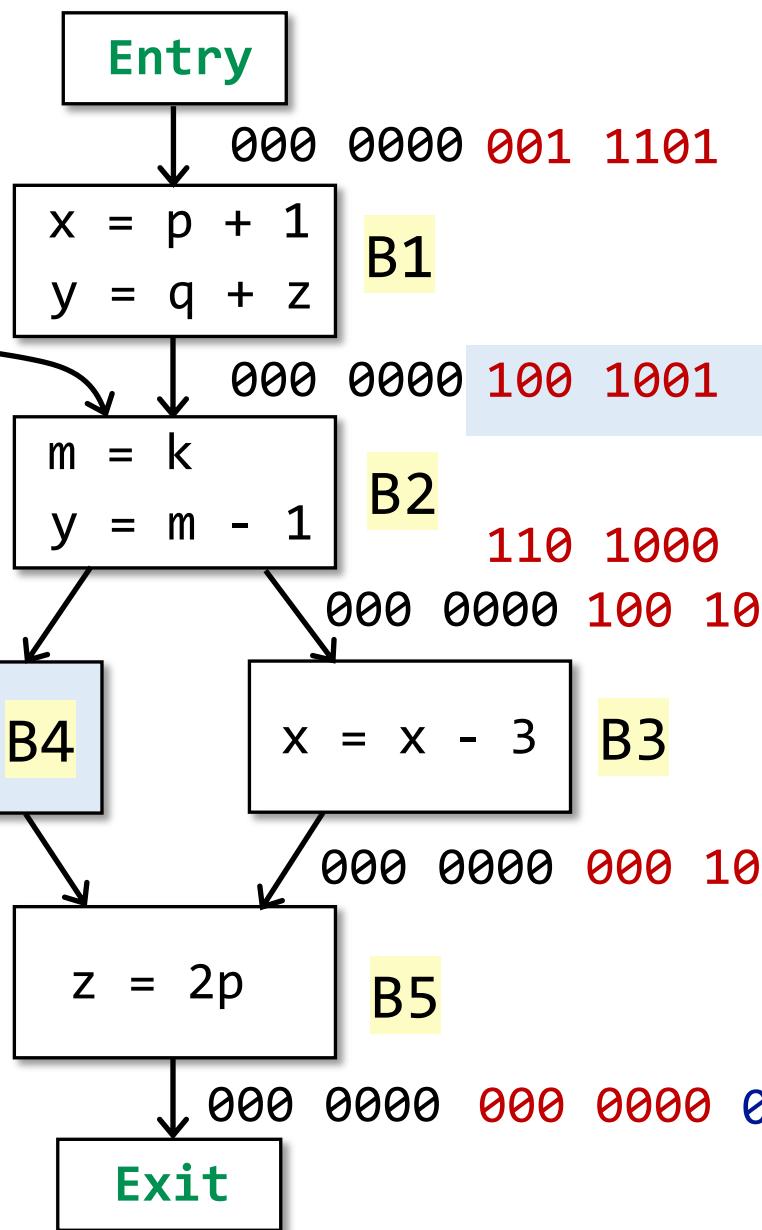
x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry

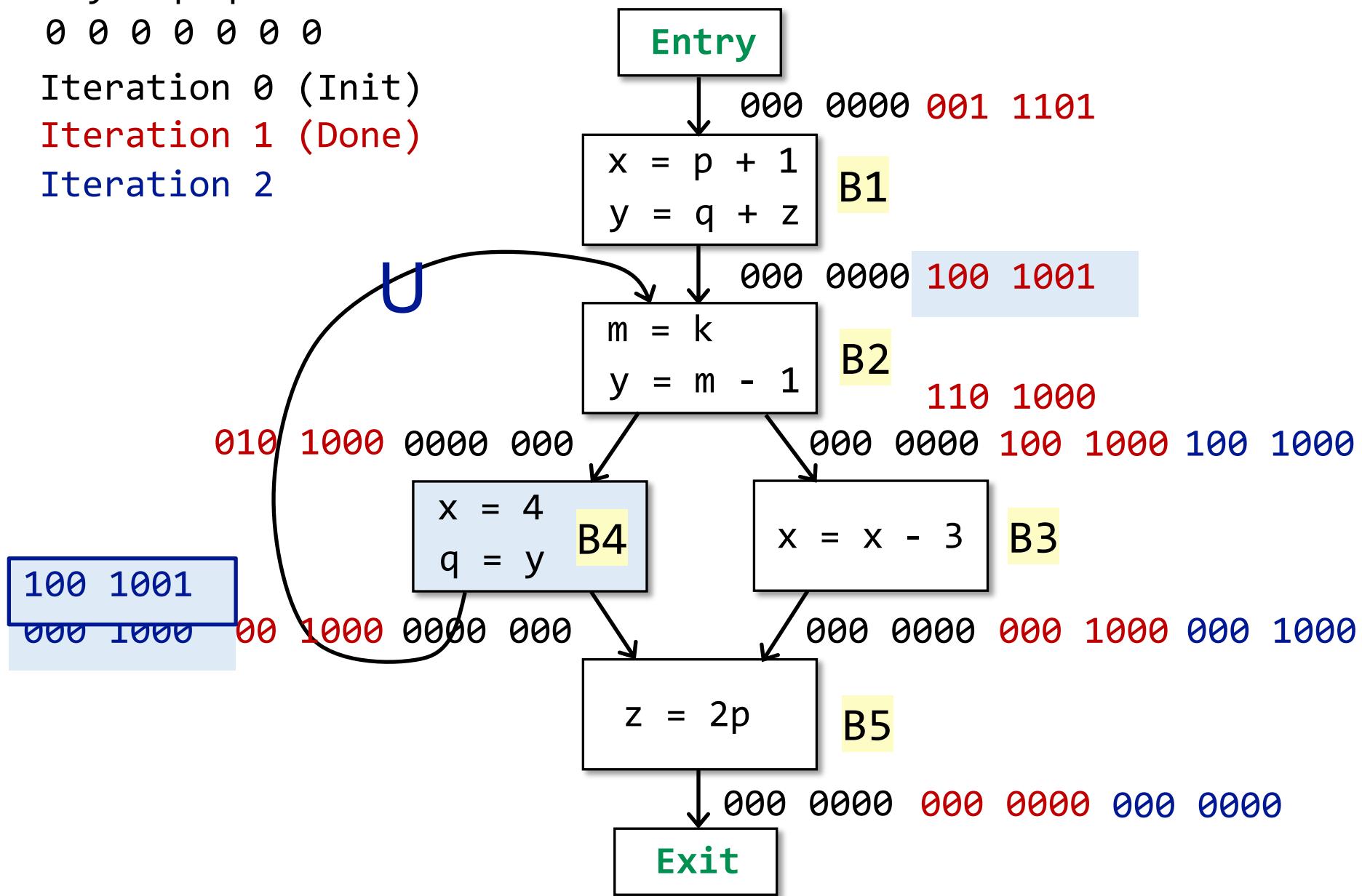


x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



000 0000 001 1101

B1

000 0000 100 1001

B2

110 1000

010 1000 0000 000

x = 4  
q = y B4

000 0000 100 1000 100 1000

x = x - 3 B3

100 1001

000 1000 000

1000 0000 000

z = 2p B5

000 0000 000 0000 000 0000

Exit

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

010 1001      10 1000 0000 000  
100 1001  
000 1000 000 1000 000 000

Entry

000 0000 001 1101

x = p + 1  
y = q + z

B1

m = k  
y = m - 1

B2

x = 4  
q = y

B4

x = x - 3

B3

z = 2p

B5

Exit

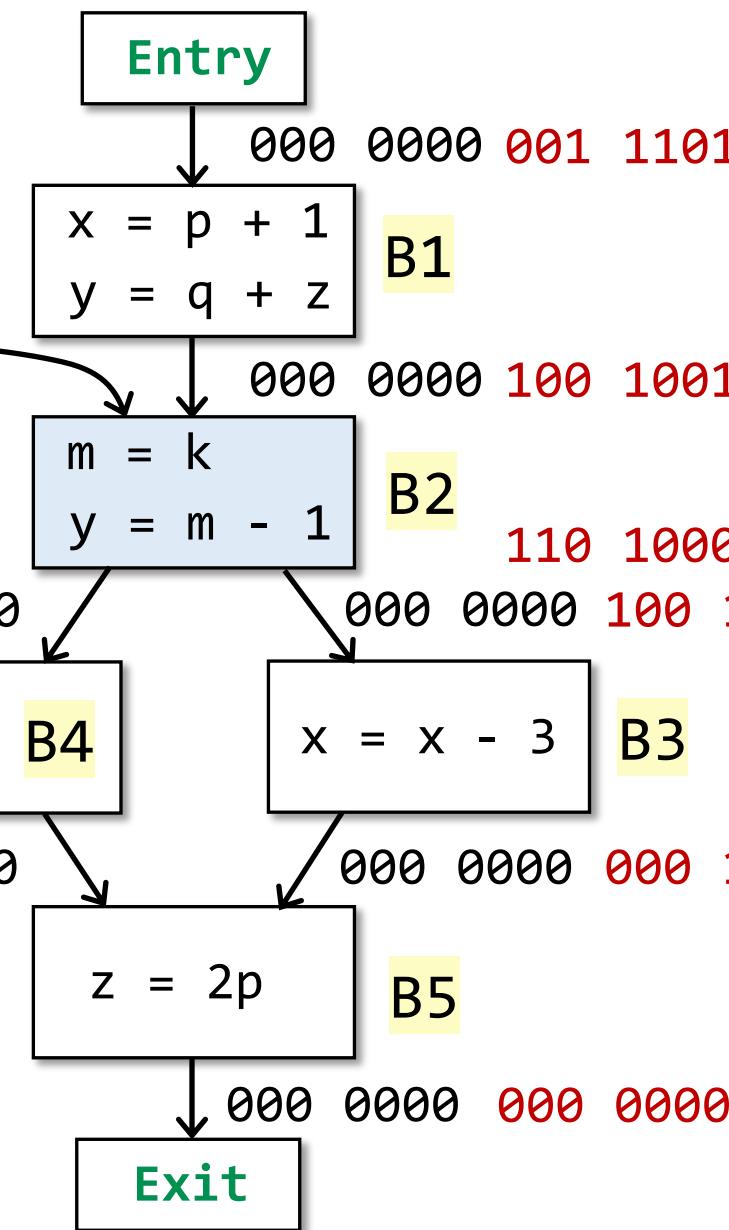
x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



|          |          |              |    |  |                   |                   |                   |
|----------|----------|--------------|----|--|-------------------|-------------------|-------------------|
| 010 1001 | 10 1000  | 0000 000     | B1 |  | 000 0000 100 1001 | 110 1000          | 100 1000 100 1000 |
| 100 1001 | 000 1000 | 000 1000     | B2 |  | 000 0000 100 1000 | 100 1000 100 1000 | 100 1000          |
| 000 1000 | 000 1000 | 000 0000 000 | B3 |  | 000 0000 000 1000 | 000 1000 000 1000 | 000 1000          |
| 100 1001 | 000 1000 | 000 1000     | B4 |  | 000 0000 000 0000 | 000 0000 000 0000 | 000 0000          |
| 000 1000 | 000 1000 | 000 0000 000 | B5 |  | 000 0000 000 0000 | 000 0000 000 0000 | 000 0000          |

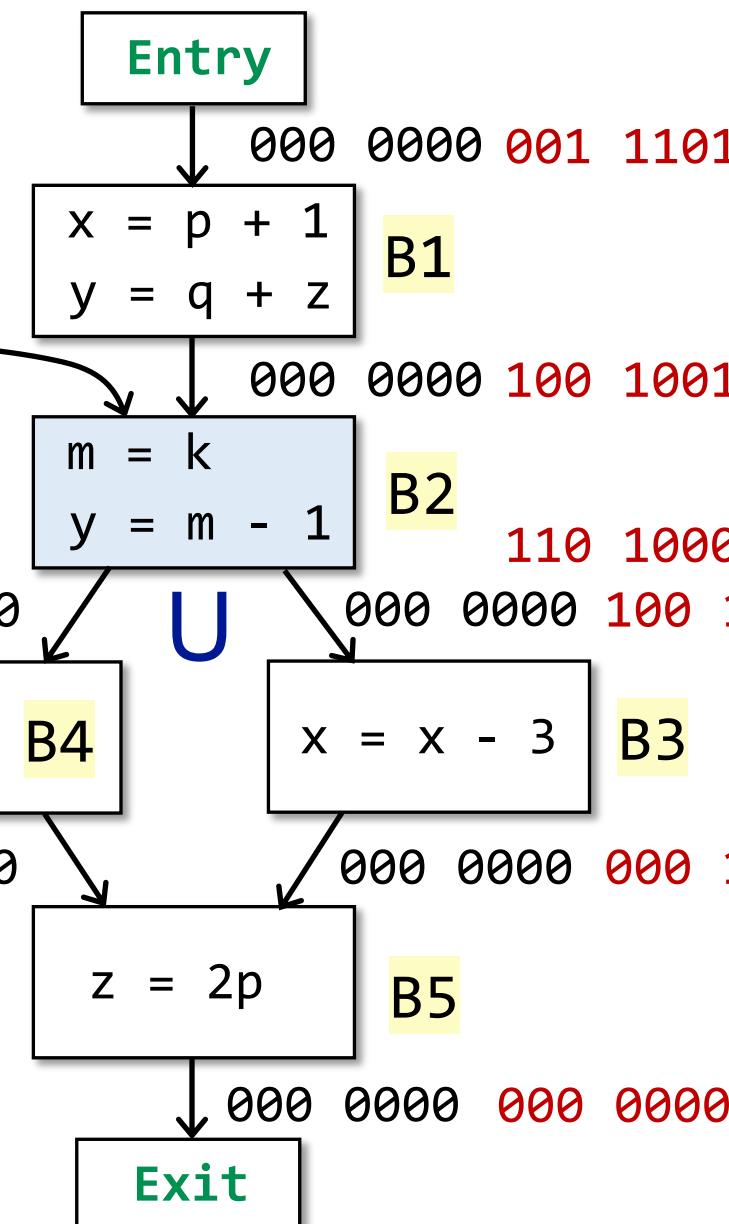
x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



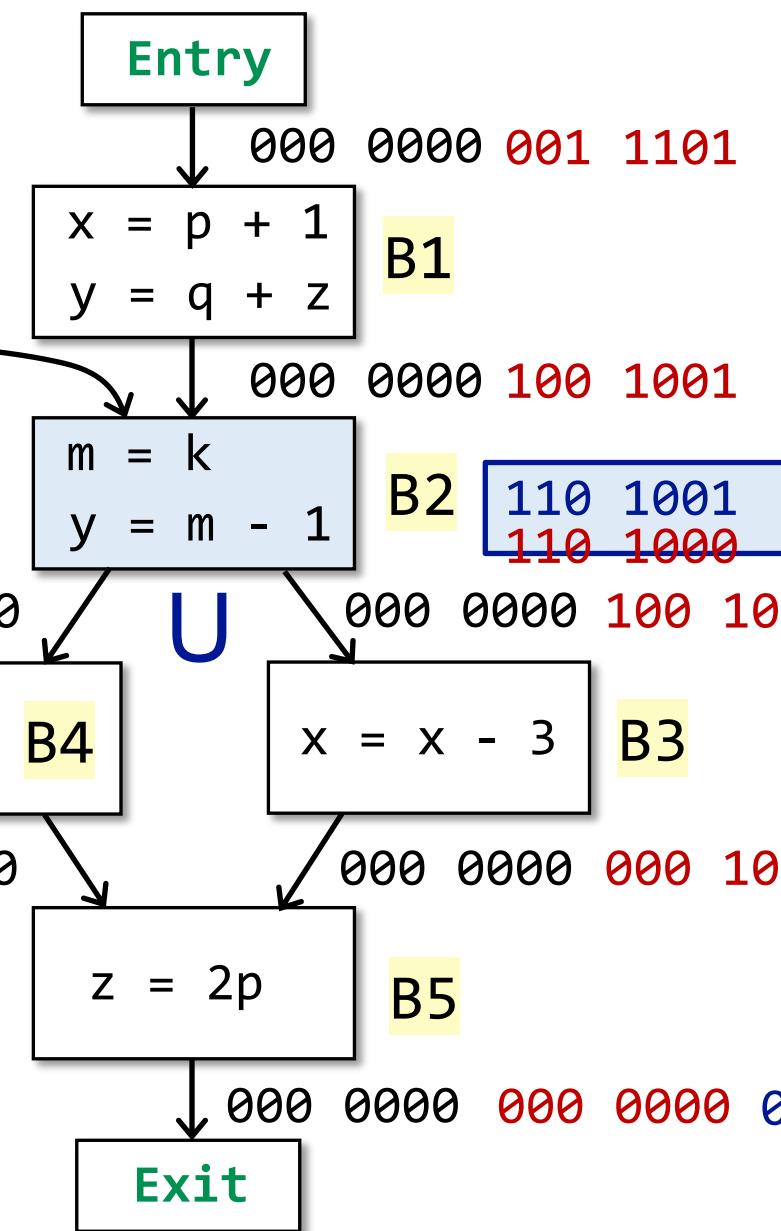
x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



|          |          |          |  |  |  |  |  |  |  |  |  |  |  |
|----------|----------|----------|--|--|--|--|--|--|--|--|--|--|--|
| 010 1001 | 10 1000  | 0000 000 |  |  |  |  |  |  |  |  |  |  |  |
| 100 1001 |          |          |  |  |  |  |  |  |  |  |  |  |  |
| 000 1000 | 000 1000 | 0000 000 |  |  |  |  |  |  |  |  |  |  |  |

x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

010 1001 010 1000 0000 000  
 100 1001  
 000 1000 000 1000 000 000

Entry

000 0000 001 1101

x = p + 1  
 y = q + z

B1

m = k  
 y = m - 1

B2

110 1001  
 110 1000

x = 4  
 q = y

B4

x = x - 3

B3

z = 2p

B5

x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

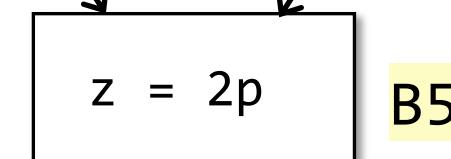
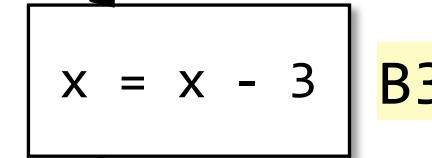
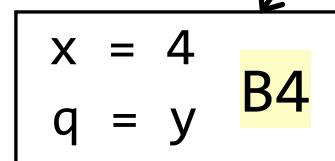
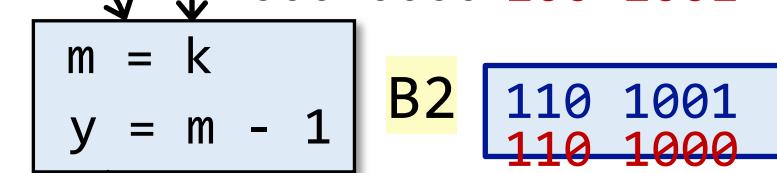
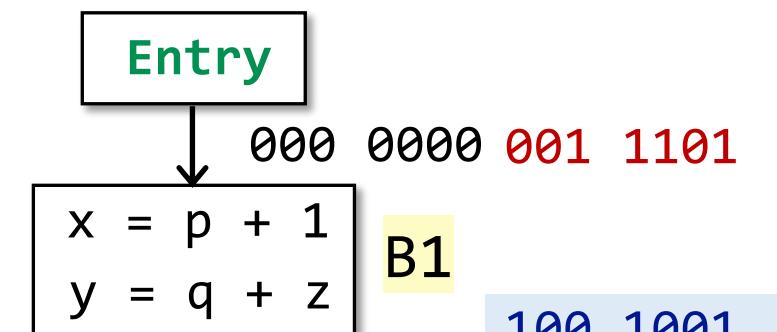
010 1001 010 1000 0000 000

100 1001

000 1000 000 1000 000 000 1000



Entry



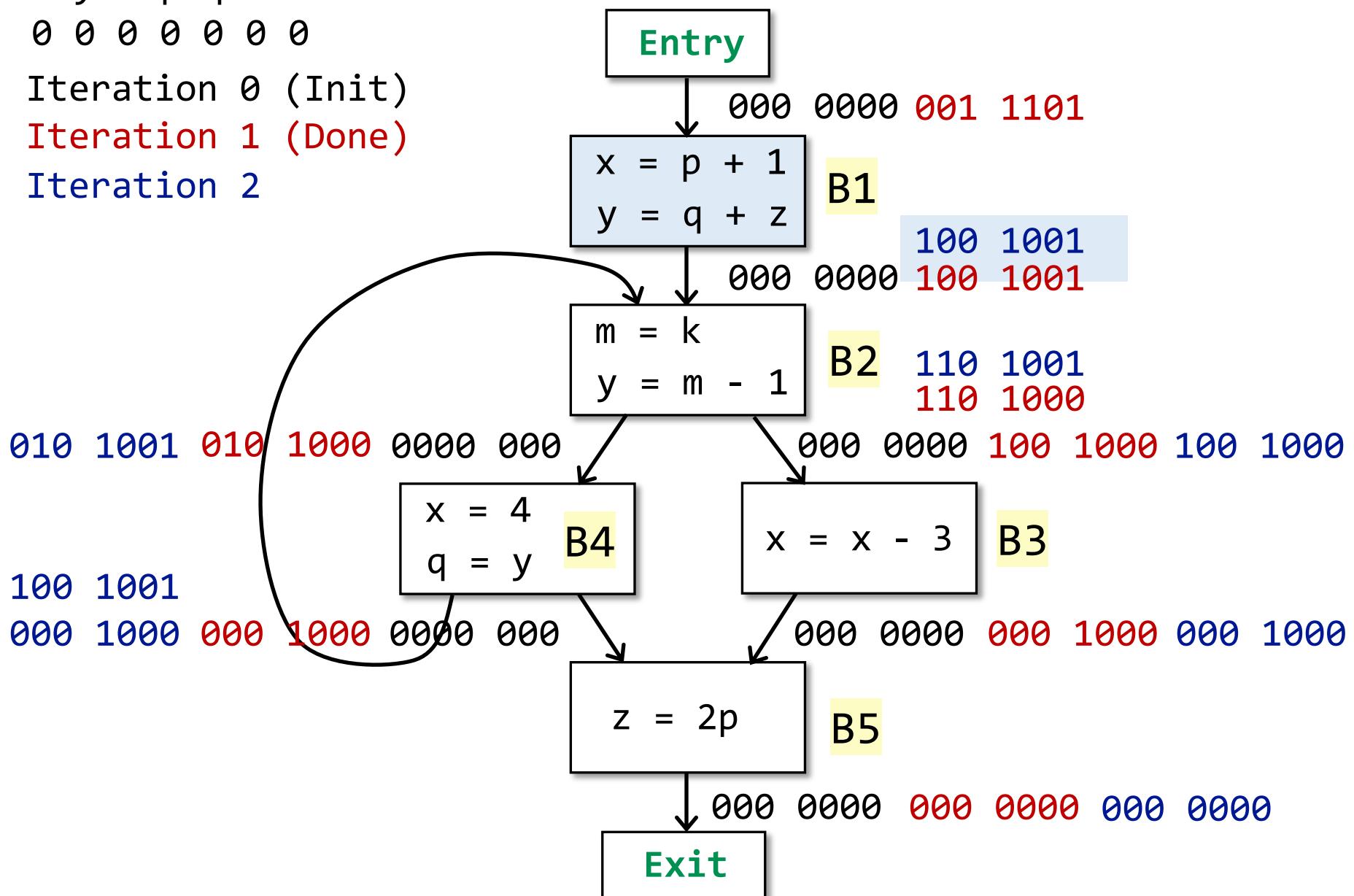
Exit

x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

010 1001 010 1000 0000 000

100 1001  
000 1000 000 1000 000 000

Entry

$x = p + 1$   
 $y = q + z$

001 1101  
001 1101

$m = k$   
 $y = m - 1$

100 1001  
100 1001

$x = 4$   
 $q = y$

110 1001  
110 1000

B1

B2

B3

$x = x - 3$   
 $z = 2p$

100 1000 100 1000 100 1000

B4

B5

Exit

x y z p q m k  
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Entry

$x = p + 1$   
 $y = q + z$

B1

$m = k$   
 $y = m - 1$

B2

$x = 4$   
 $q = y$

B4

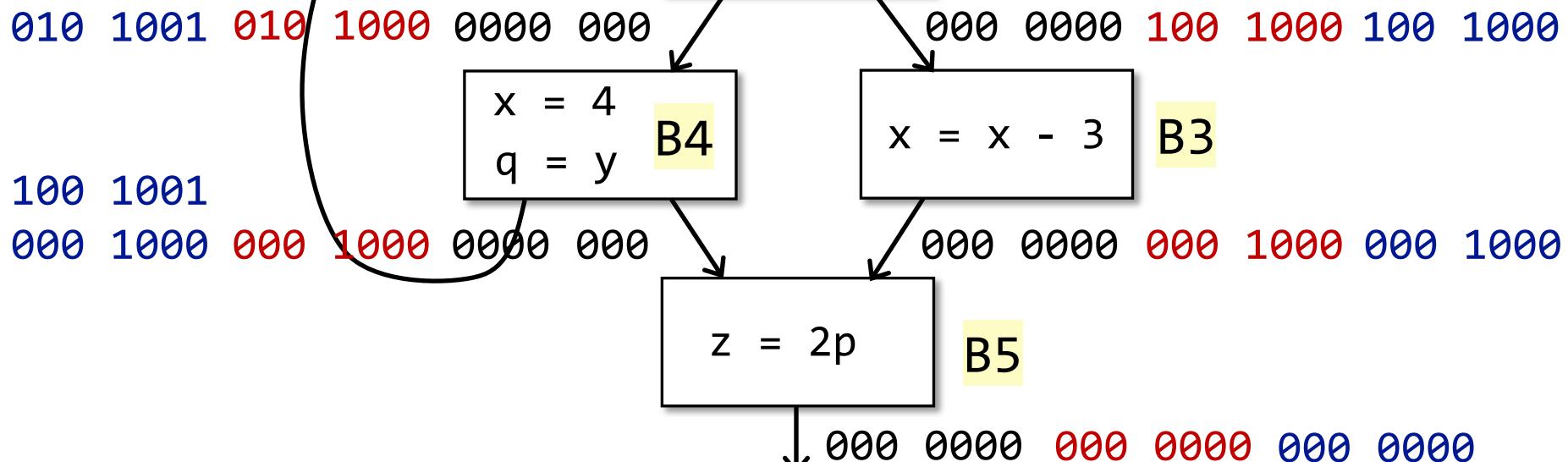
$x = x - 3$

B3

$z = 2p$

B5

Exit



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Entry

$x = p + 1$   
 $y = q + z$

B1

$m = k$   
 $y = m - 1$

B2

$x = 4$   
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

Changes occur in  
IN[] of B4

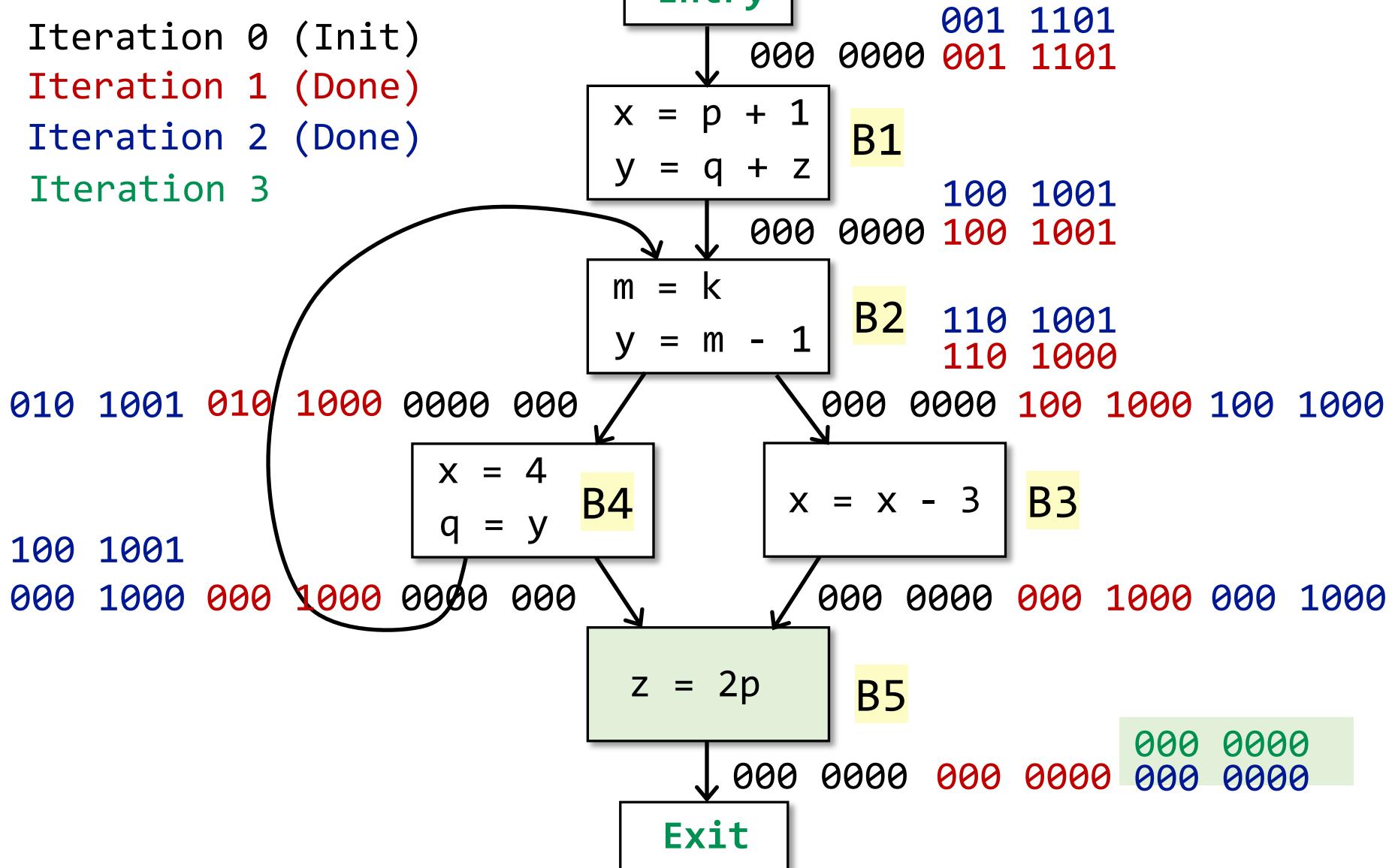
x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001 010 1000 0000 000  
100 1001  
000 1000 000 1000 0000 000  
000 1000

Entry

000 0000

001 1101  
001 1101

x = p + 1  
y = q + z

B1

100 1001  
100 1001

m = k  
y = m - 1

B2

110 1001  
110 1000

x = 4  
q = y

B4

x = x - 3

B3

z = 2p

B5

Exit

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001 010 1000 0000 000  
100 1001  
000 1000 000 1000 0000 000  
000 1000

Entry

000 0000

001 1101  
001 1101

x = p + 1  
y = q + z

B1

100 1001  
100 1001

m = k  
y = m - 1

B2

110 1001  
110 1000

x = 4  
q = y

B4

x = x - 3

B3

000 1000  
000 1000

z = 2p

B5

000 0000  
000 0000

Exit

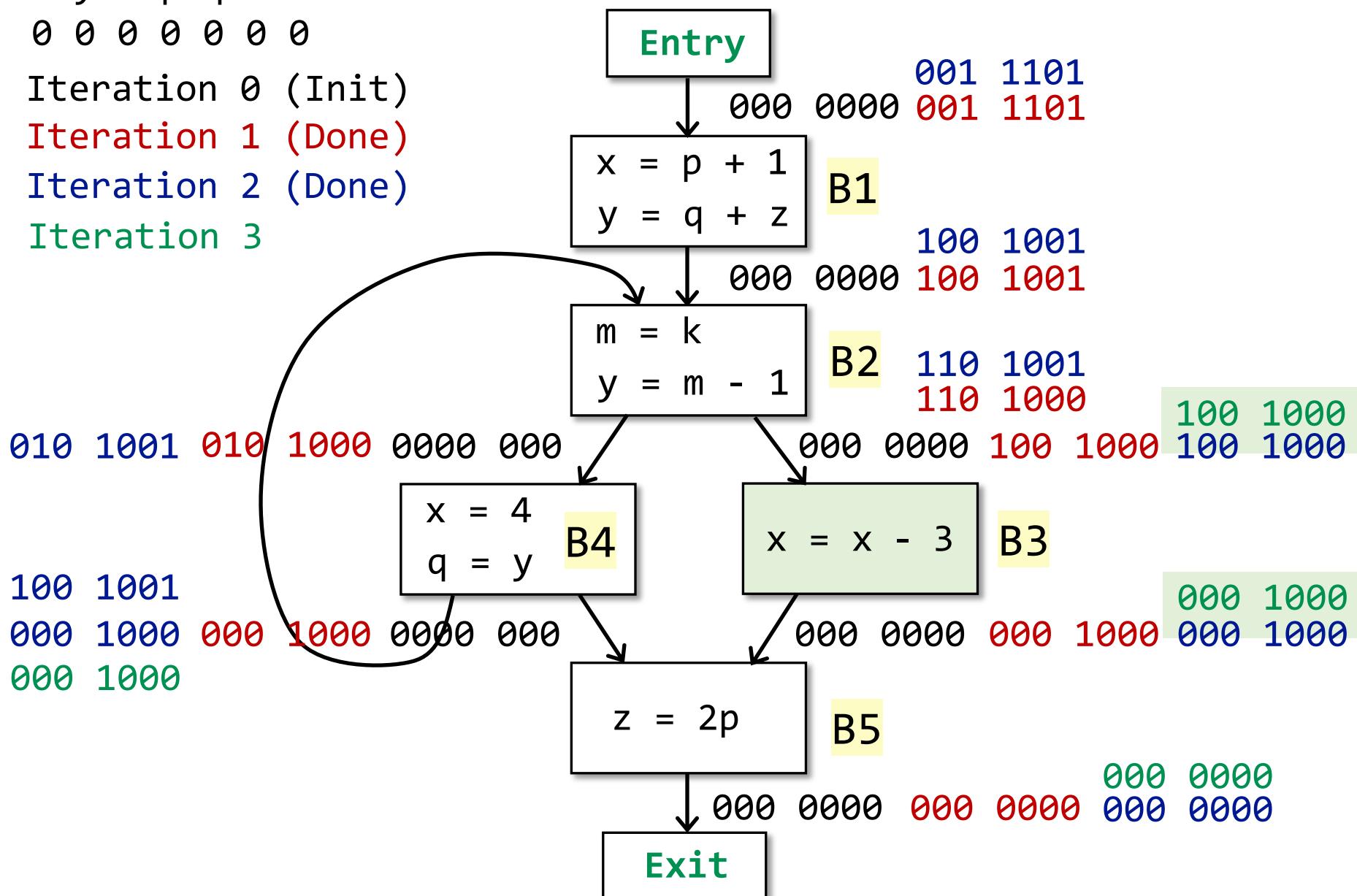
x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



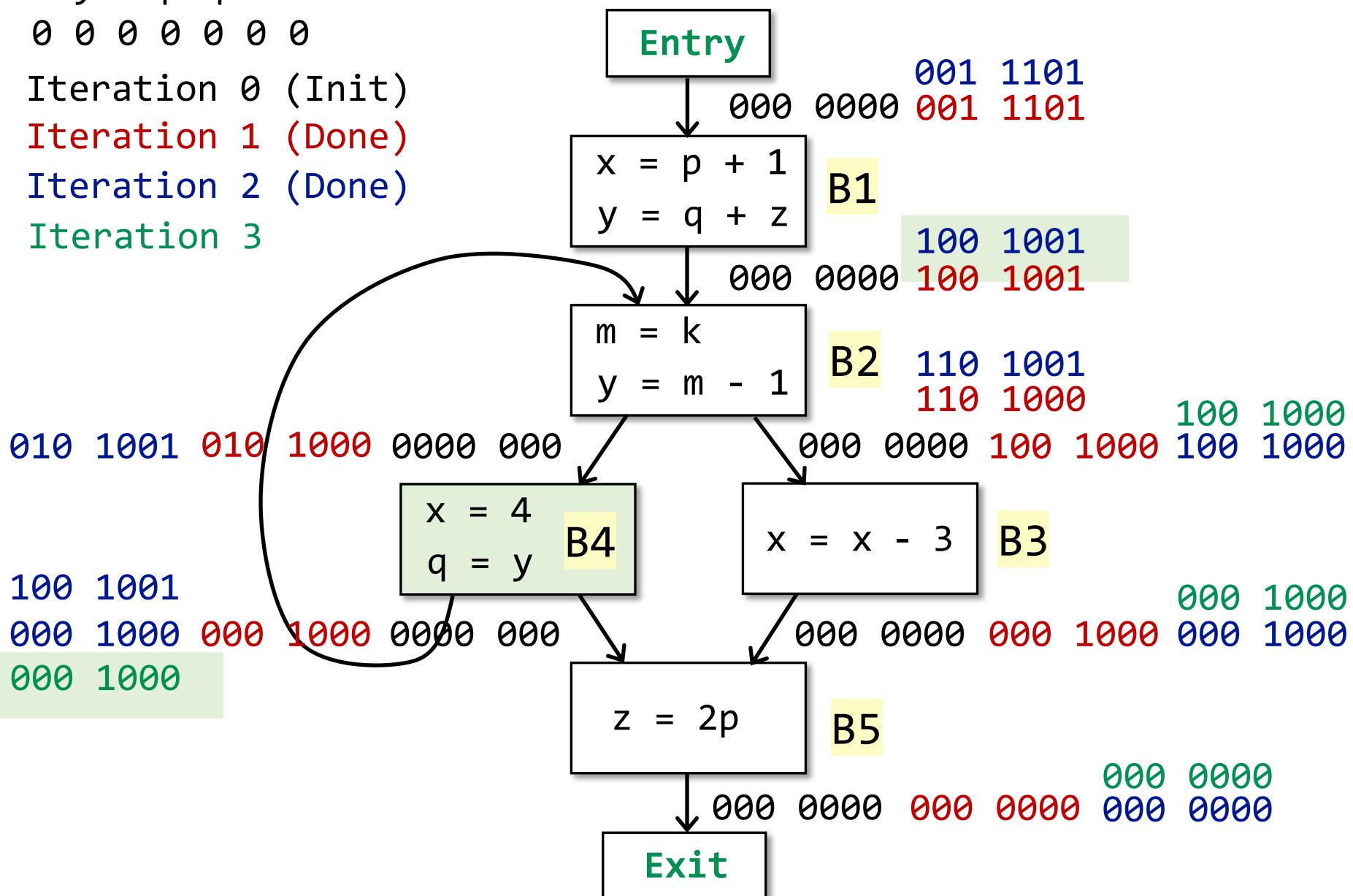
x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



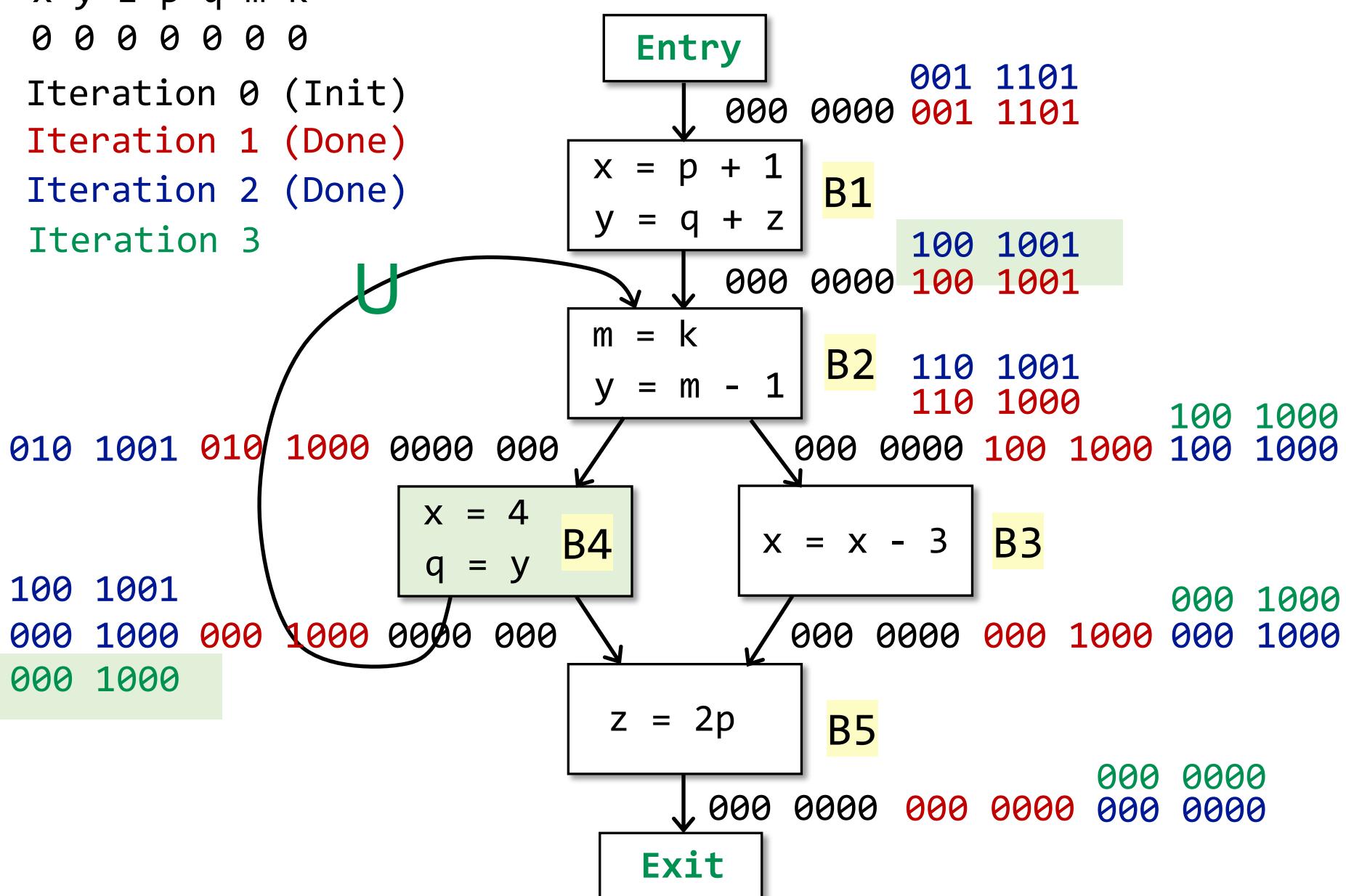
x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



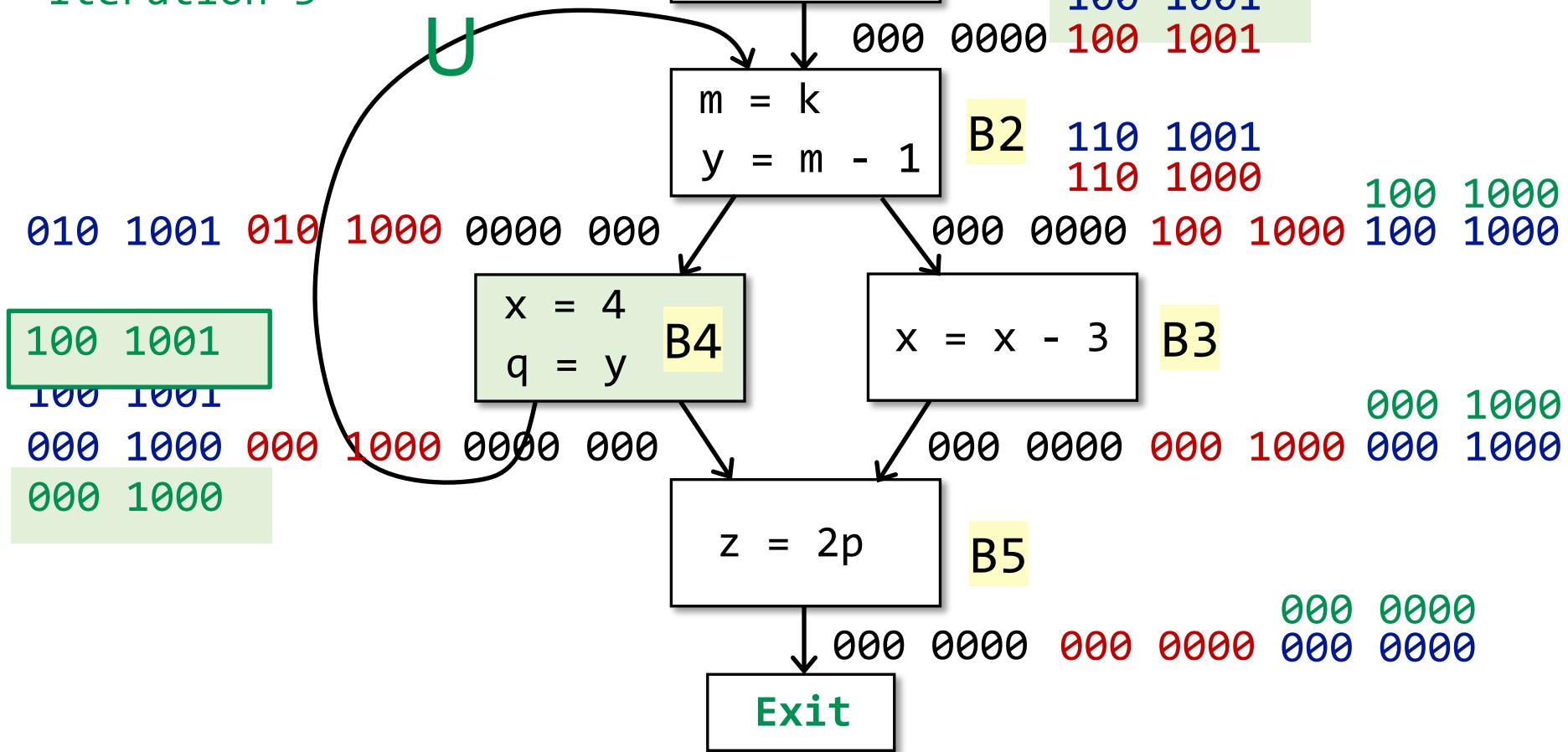
x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



x y z p q m k  
0 0 0 0 0 0 0

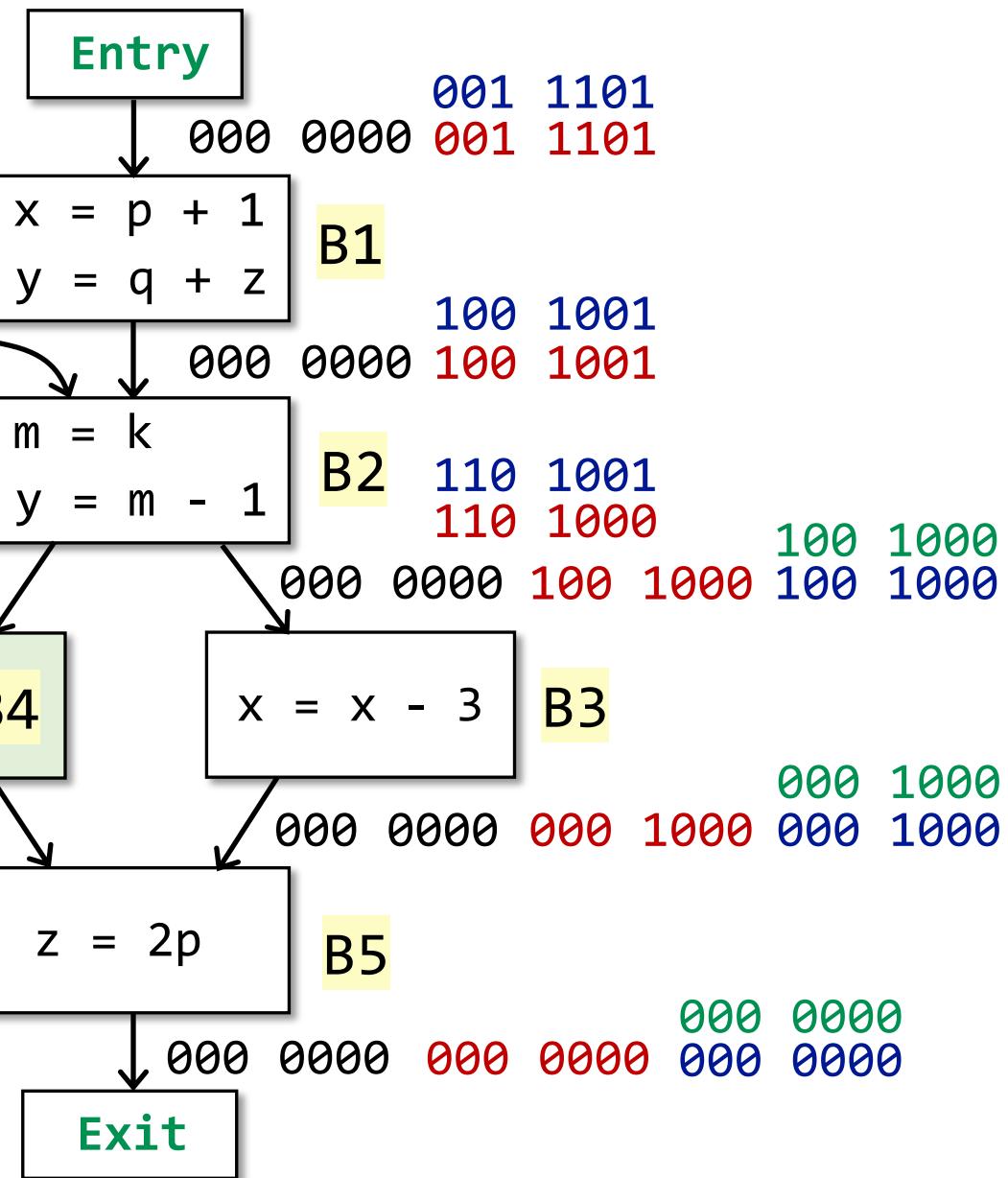
Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

|              |               |               |               |               |               |               |
|--------------|---------------|---------------|---------------|---------------|---------------|---------------|
| 010 1001 010 | 1000 0000 000 | 000 1000 000  | 000 1000 000  | 000 1000 000  | 000 1000 000  | 000 1000 000  |
| 100 1001     |               | 1000 0000 000 | 1000 0000 000 | 1000 0000 000 | 1000 0000 000 | 1000 0000 000 |
| 100 1001     |               | 000 1000 000  | 000 1000 000  | 000 1000 000  | 000 1000 000  | 000 1000 000  |
| 000 1000     |               | 000 1000 000  | 000 1000 000  | 000 1000 000  | 000 1000 000  | 000 1000 000  |



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

|     |      |
|-----|------|
| 010 | 1001 |
| 010 | 1001 |
| 010 | 1001 |
| 100 | 1001 |
| 100 | 1001 |
| 000 | 1000 |
| 000 | 1000 |

Entry

x = p + 1  
y = q + z

B1

m = k  
y = m - 1

B2

x = 4  
q = y

B4

x = x - 3

B3

z = 2p

B5

Exit

001 1101  
001 1101

100 1001  
100 1001

110 1001  
110 1000

100 1000  
100 1000

000 1000  
000 1000

000 0000  
000 0000

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

|          |          |              |          |
|----------|----------|--------------|----------|
| 010 1001 | 010 1001 | 000 0000 000 | 001 1101 |
| 010 1001 | 010 1001 | 000 0000 000 | 001 1101 |
| 100 1001 | 100 1001 | 000 0000 000 | 100 1001 |
| 100 1001 | 100 1001 | 000 0000 000 | 100 1001 |
| 000 1000 | 000 1000 | 000 0000 000 | 110 1001 |
| 000 1000 | 000 1000 | 000 0000 000 | 110 1000 |
| 000 1000 | 000 1000 | 000 0000 000 | 100 1000 |
| 000 1000 | 000 1000 | 000 0000 000 | 100 1000 |

Entry

$x = p + 1$   
 $y = q + z$

B1

$m = k$   
 $y = m - 1$

B2

$x = 4$   
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

|          |          |          |          |           |          |          |          |
|----------|----------|----------|----------|-----------|----------|----------|----------|
| 010 1001 | 010 1001 | 010 1001 | 010 1000 | 0000 0000 | 000 0000 | 001 1101 | 001 1101 |
| 010 1001 | 010 1001 | 010 1001 | 010 1000 | 0000 0000 | 000 0000 | 100 1001 | 100 1001 |
| 100 1001 | 100 1001 | 100 1001 | 100 1000 | 0000 0000 | 000 0000 | 110 1001 | 110 1000 |
| 100 1001 | 100 1001 | 100 1001 | 100 1000 | 0000 0000 | 000 0000 | 100 1000 | 100 1000 |
| 000 1000 | 000 1000 | 000 1000 | 000 1000 | 0000 0000 | 000 0000 | 000 1000 | 000 1000 |
| 000 1000 | 000 1000 | 000 1000 | 000 1000 | 0000 0000 | 000 0000 | 000 1000 | 000 1000 |

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

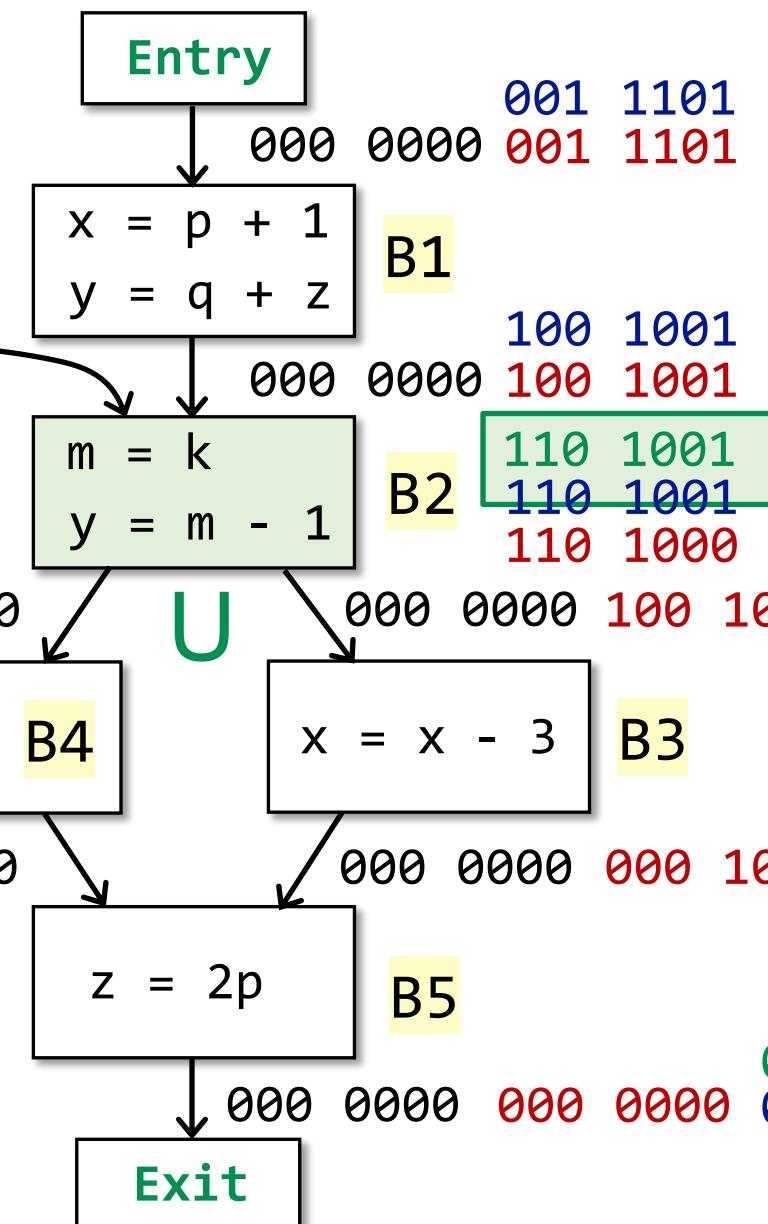
Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

|          |          |          |          |         |         |         |
|----------|----------|----------|----------|---------|---------|---------|
| 010 1001 | 010 1001 | 010 1000 | 0000 000 | 000 000 | 000 000 | 000 000 |
| 010 1001 | 010 1001 | 010 1000 | 0000 000 | 000 000 | 000 000 | 000 000 |
| 100 1001 | 100 1001 | 100 1000 | 000 000  | 000 000 | 000 000 | 000 000 |
| 100 1001 | 100 1001 | 100 1000 | 000 000  | 000 000 | 000 000 | 000 000 |
| 000 1000 | 000 1000 | 000 1000 | 000 000  | 000 000 | 000 000 | 000 000 |
| 000 1000 | 000 1000 | 000 1000 | 000 000  | 000 000 | 000 000 | 000 000 |

Entry



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

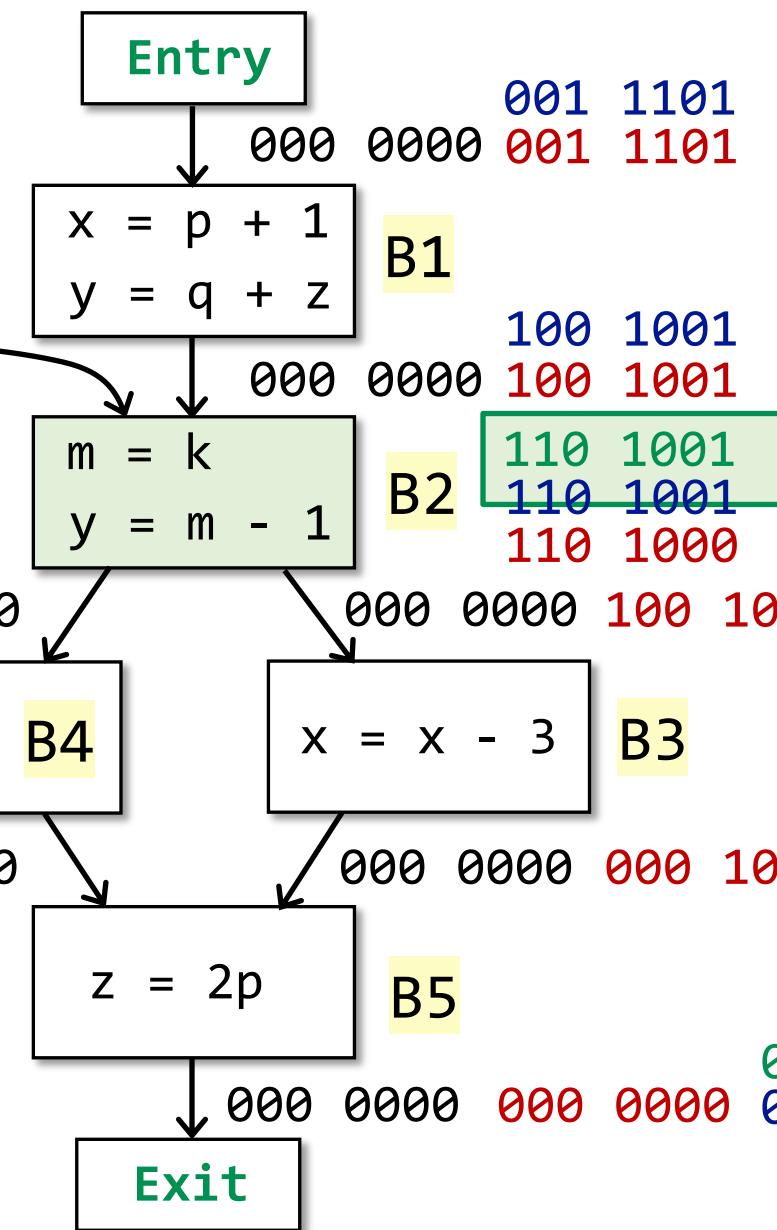
Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001  
010 1001 010 1000 0000 000  
100 1001  
100 1001  
000 1000 000 1000 0000 000  
000 1000

Entry



x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001

010 1001

010 1000

0000 000

100 1001

100 1001

000 1000

000 1000

Entry

x = p + 1  
y = q + z

B1

001 1101  
001 1101

100 1001  
100 1001  
100 1001

m = k  
y = m - 1

B2

110 1001  
110 1001  
110 1000

x = 4  
q = y

B4

x = x - 3

B3

z = 2p

B5

Exit

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001

010 1001 010 1000 0000 000

100 1001

100 1001

000 1000

000 1000

Entry

$x = p + 1$   
 $y = q + z$

B1

001 1101  
001 1101

100 1001  
100 1001  
100 1001

$m = k$   
 $y = m - 1$

B2

110 1001  
110 1001  
110 1000

$x = 4$   
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| x | y | z | p | q | m | k |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001

010 1001

010 1000 0000 000

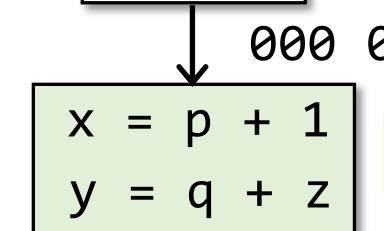
100 1001

100 1001

000 1000

000 1000

Entry



|     |      |
|-----|------|
| 001 | 1101 |
| 001 | 1101 |
| 001 | 1101 |

|     |      |
|-----|------|
| 100 | 1001 |
| 100 | 1001 |
| 100 | 1001 |

|     |      |
|-----|------|
| 110 | 1001 |
| 110 | 1001 |
| 110 | 1000 |

|     |      |
|-----|------|
| 100 | 1000 |
| 100 | 1000 |

|     |      |
|-----|------|
| 000 | 1000 |
| 000 | 1000 |

|     |      |
|-----|------|
| 000 | 1000 |
| 000 | 1000 |

|     |      |
|-----|------|
| 000 | 0000 |
| 000 | 0000 |

x y z p q m k

0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

$x = p + 1$   
 $y = q + z$

B1

$m = k$   
 $y = m - 1$

B2

$x = 4$   
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

001 1101

001 1101

001 1101

100 1001

100 1001

100 1001

110 1001

110 1001

110 1000

100 1000

100 1000

000 1000

000 1000

000 0000

000 0000

x y z p q m k

0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

010 1001

010 1001

010 1000

0000 000

100 1001

100 1001

000 1000

000 1000

0000 000

000 1000

Entry

$x = p + 1$   
 $y = q + z$

B1

$m = k$   
 $y = m - 1$

B2

$x = 4$   
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

001 1101

001 1101

001 1101

100 1001

100 1001

100 1001

110 1001

110 1001

110 1000

100 1000

100 1000

000 1000

000 1000

000 0000

000 0000

No changes occur  
in any IN[]

x y z p q m k  
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

010 1001

010 1001

100 1001

100 1001

000 1000

000 1000

010 1000

1000 0000 000

000 1000 000

1000 0000 000

000 0000 000

000 0000 000

000 0000 000

000 0000 000

001 1101

001 1101

001 1101

100 1001

100 1001

100 1001

110 1001

110 1001

110 1000

100 1000

100 1000

000 1000

000 1000

000 0000

000 0000

Entry

x = p + 1  
y = q + z

B1

m = k  
y = m - 1

B2

x = 4  
q = y

B4

x = x - 3

B3

z = 2p

B5

Exit

Final analysis result

# Data Flow Analysis Applications

(I) Reaching Definitions Analysis

(II) Live Variables Analysis

(III) Available Expressions Analysis

# Available Expressions Analysis

An expression  $x \ op \ y$  is available at program point p if (1) **all** paths from the entry to p **must** pass through the evaluation of  $x \ op \ y$ , and (2) after the last evaluation of  $x \ op \ y$ , there is no redefinition of x or y

# Available Expressions Analysis

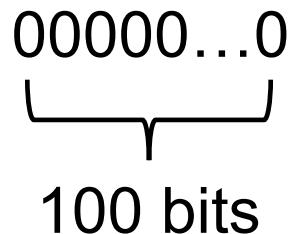
An expression  $x \ op \ y$  is available at program point p if (1) **all** paths from the entry to p **must** pass through the evaluation of  $x \ op \ y$ , and (2) after the last evaluation of  $x \ op \ y$ , there is no redefinition of x or y

- This definition means at program p, we can replace expression  $x \ op \ y$  by the result of its last evaluation
- The information of available expressions can be used for detecting global common subexpressions.

# Understanding Available Expressions Analysis

- Data Flow Values/Facts
    - All the expressions in a program
    - Can be represented by bit vectors
- e.g., E1, E2, E3, E4, ..., E100 (100 expressions)

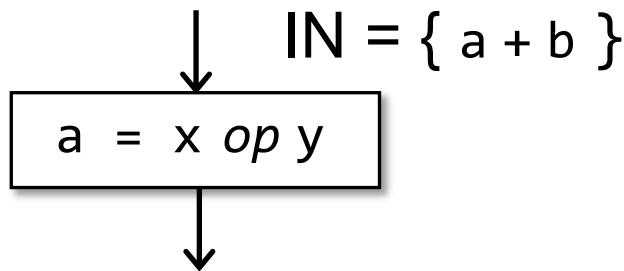
*Abstraction*



Bit i from the left represents expression  $E_i$

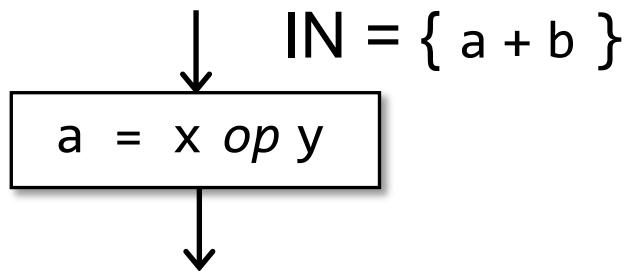
# Understanding Available Expressions Analysis

*Safe-approximation*



# Understanding Available Expressions Analysis

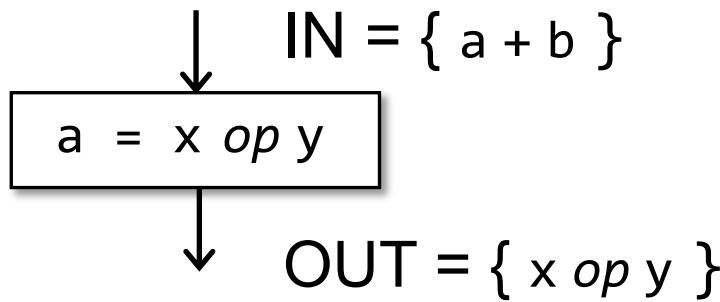
*Safe-approximation*



- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

# Understanding Available Expressions Analysis

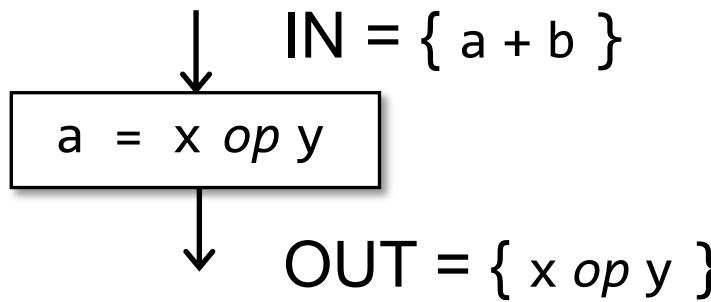
*Safe-approximation*



- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

# Understanding Available Expressions Analysis

*Safe-approximation*



- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

# Understanding Available Expressions Analysis

*Safe-approximation*

$$\downarrow \quad \text{IN} = \{ a + b \}$$

$$a = x \ op \ y$$

$$\downarrow \quad \text{OUT} = \{ x \ op \ y \ }$$

- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$

$$a = e^{16} * x$$

$$\begin{aligned} x &= \dots \\ b &= e^{16} * x \end{aligned}$$

$$c = e^{16} * x$$

# Understanding Available Expressions Analysis

*Safe-approximation*

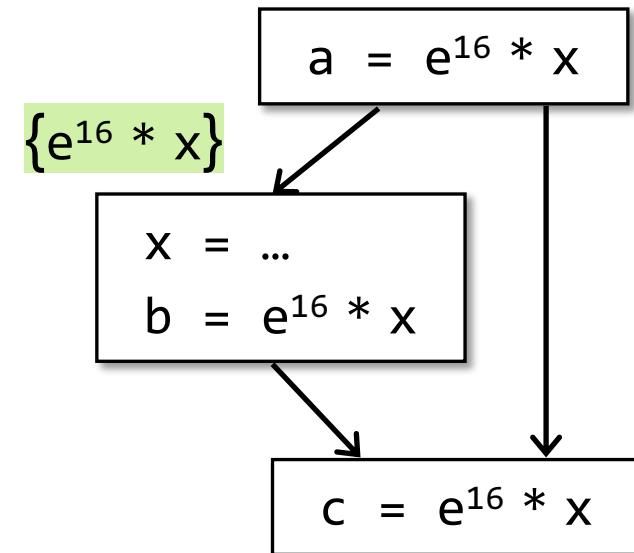
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- Add to OUT the expression  $x \ op \ y$  (gen)
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# Understanding Available Expressions Analysis

*Safe-approximation*

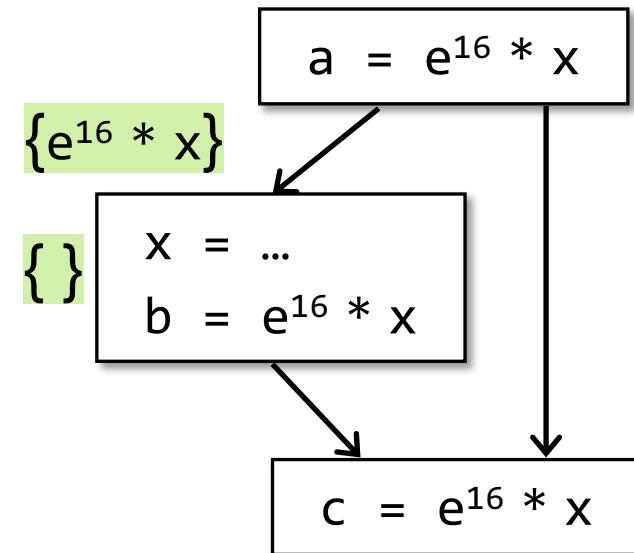
$$\downarrow \quad \text{IN} = \{ a + b \}$$

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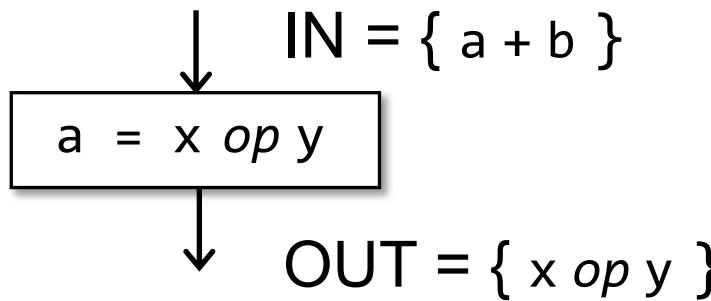
- Add to OUT the expression  $x \ op \ y$  (gen)
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$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$



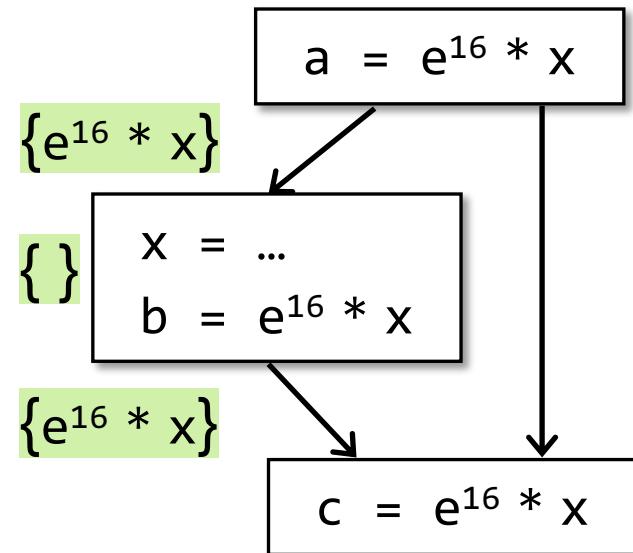
# Understanding Available Expressions Analysis

*Safe-approximation*



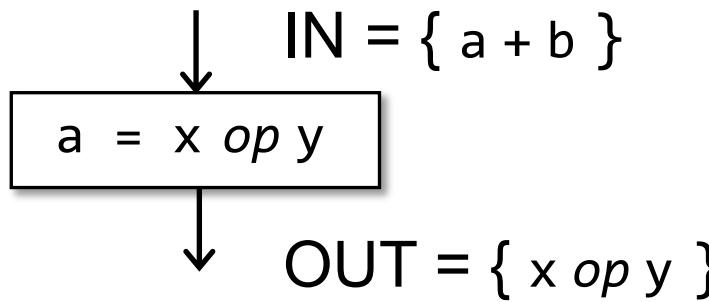
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- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$



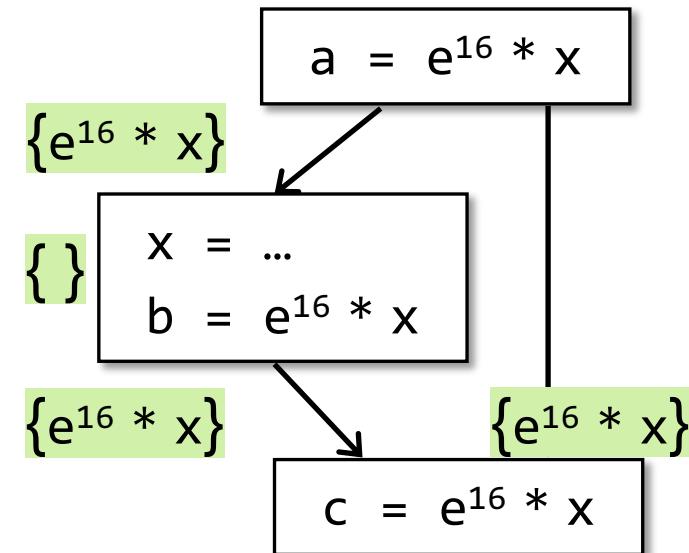
# Understanding Available Expressions Analysis

*Safe-approximation*



- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$



# Understanding Available Expressions Analysis

*Safe-approximation*

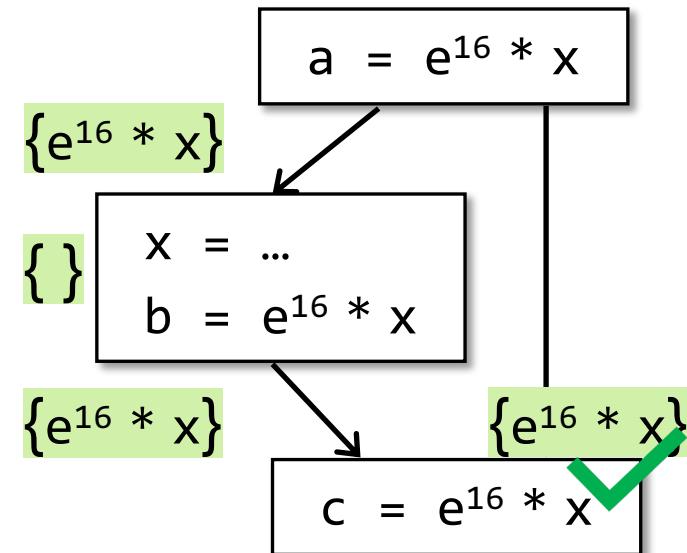
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$$\downarrow \quad \text{OUT} = \{ x \ op \ y \ }$$

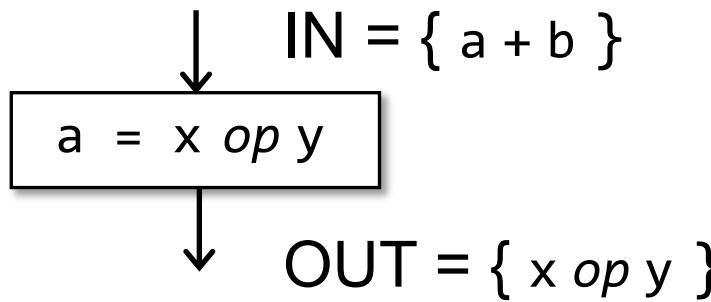
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$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$



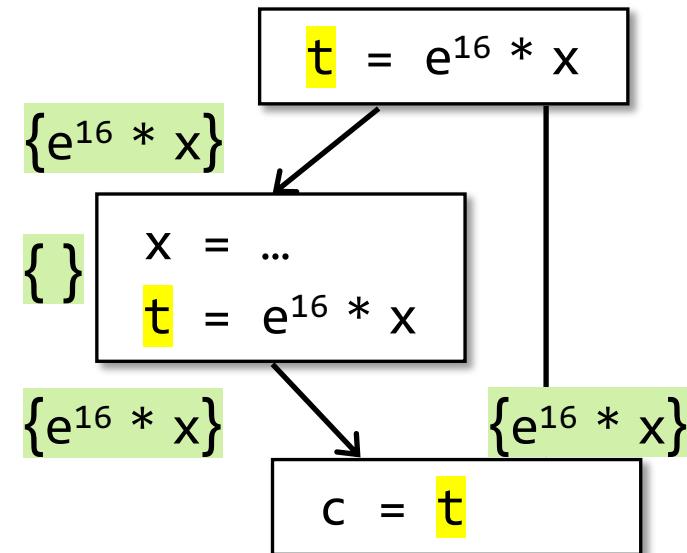
# Understanding Available Expressions Analysis

*Safe-approximation*



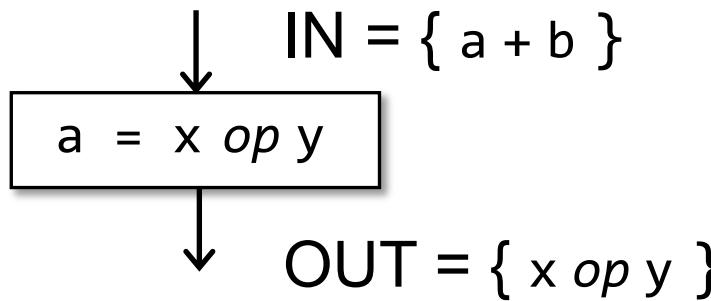
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# Understanding Available Expressions Analysis

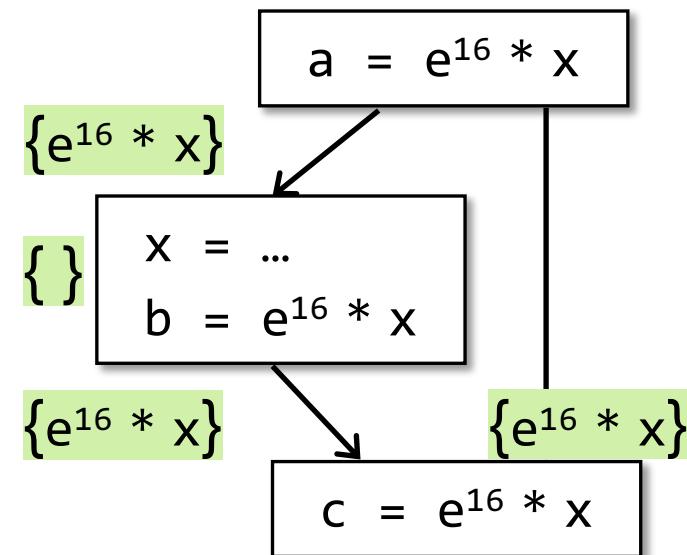
Safe-approximation



- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

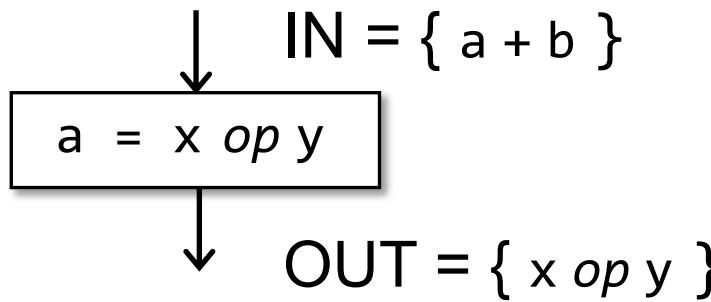
$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

$$IN[B] = \bigcap_{P \text{ a predecessor of } B} OUT[P]$$



# Understanding Available Expressions Analysis

Safe-approximation

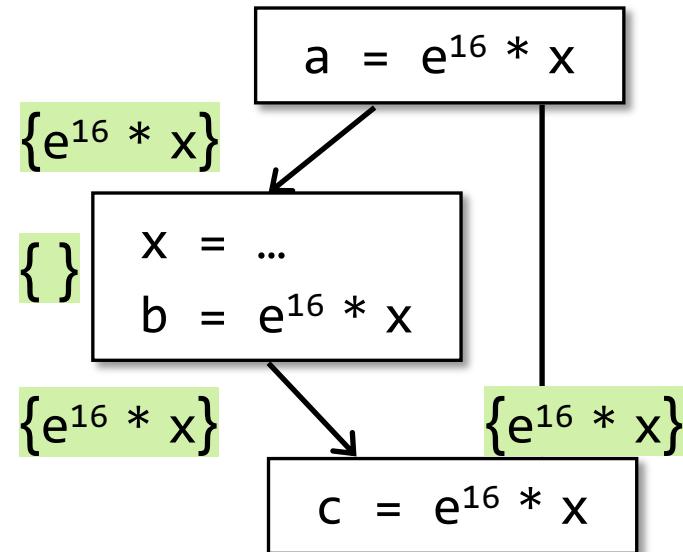


- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

$$IN[B] = \bigcap_{P \text{ a predecessor of } B} OUT[P]$$

All paths from entry to point p must pass through the evaluation of  $x \ op \ y$



# Understanding Available Expressions Analysis

$$\downarrow \quad \text{IN} = \{ a + b \}$$

*Safe-approximation*

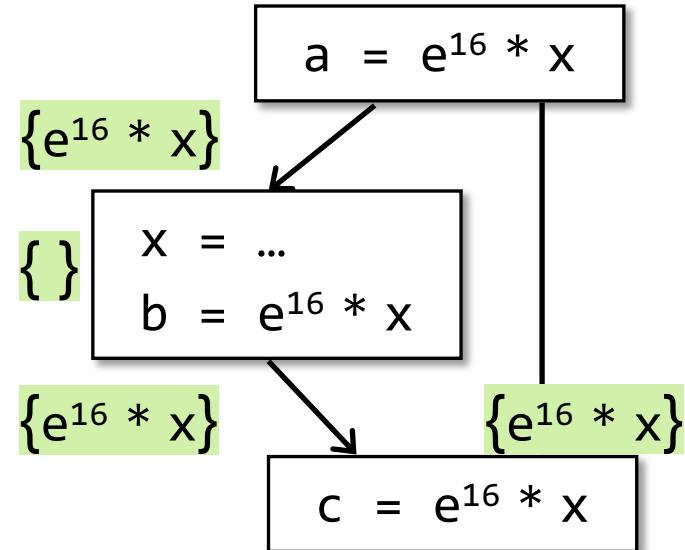
For safety of the analysis, it may report an expression as unavailable even if it is truly available (must analysis  $\rightarrow$  under-approximation)

- Add to OUT the expression  $x \ op \ y$  (gen)
- Delete from IN any expression involving variable a (kill)

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# Understanding Available Expressions Analysis

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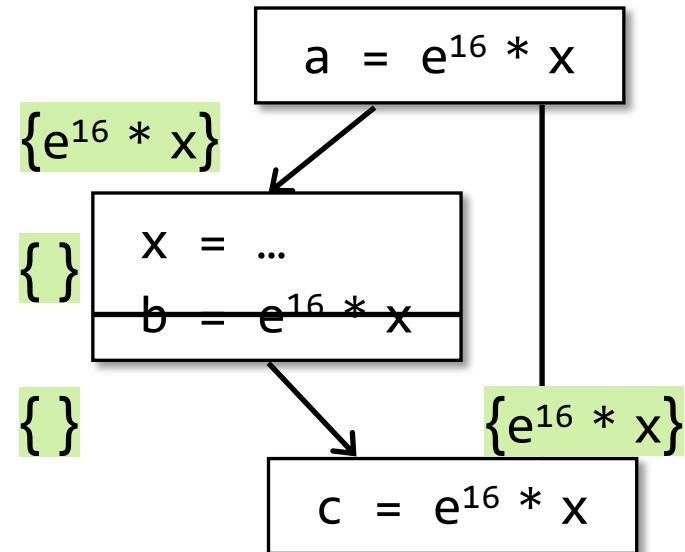
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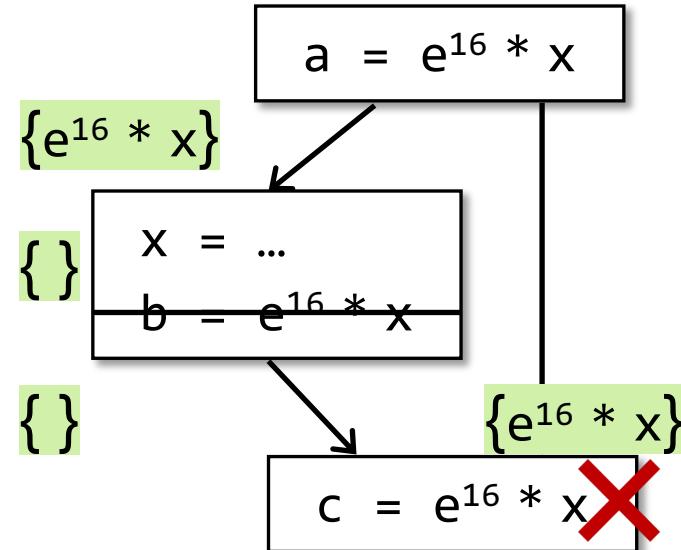
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# Understanding Available Expressions Analysis

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*Safe-approximation*

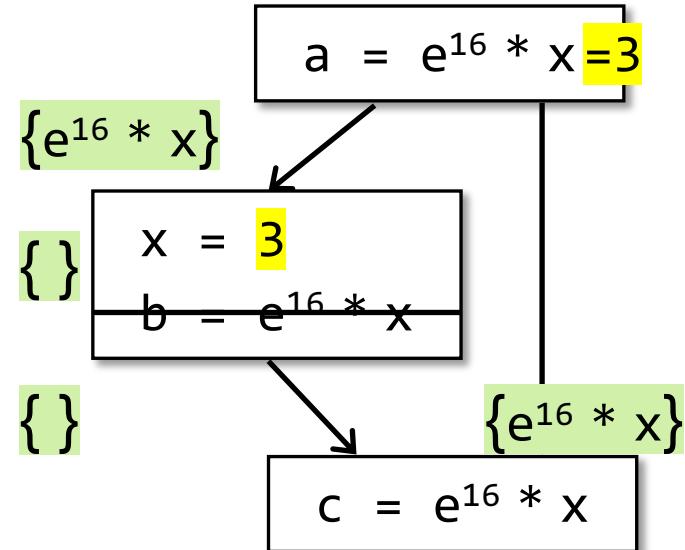
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All paths from entry to point p must pass through the evaluation of  $x \ op \ y$



# Understanding Available Expressions Analysis

$$\downarrow \quad \text{IN} = \{ a + b \}$$

**Safe**-approximation

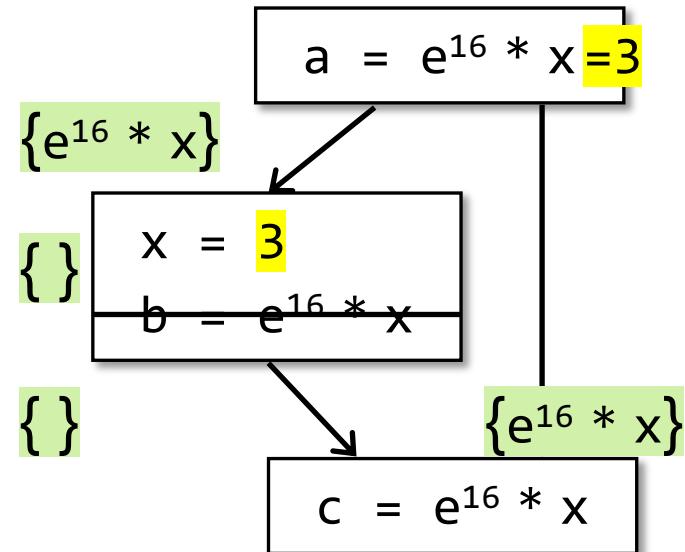
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$$\text{IN}[B] = \bigcap_{P \text{ a predecessor of } B} \text{OUT}[P]$$

All paths from entry to point p must pass through the evaluation of  $x \ op \ y$



# Algorithm of Available Expressions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

**OUTPUT:**  $IN[B]$  and  $OUT[B]$  for each basic block  $B$

**METHOD:**

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[B] = U;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[B] =  $\bigcap_{P \text{ a predecessor of } B} OUT[P];$   
            OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
        }  
    }
```

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        }  
    }
```

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        OUT[ $B$ ] =  $gen_B \cup (IN[B] - kill_B)$ ;  
    }
```

# Algorithm of Available Expressions Analysis

**INPUT:** CFG ( $kill_B$  and  $gen_B$  computed for each basic block  $B$ )

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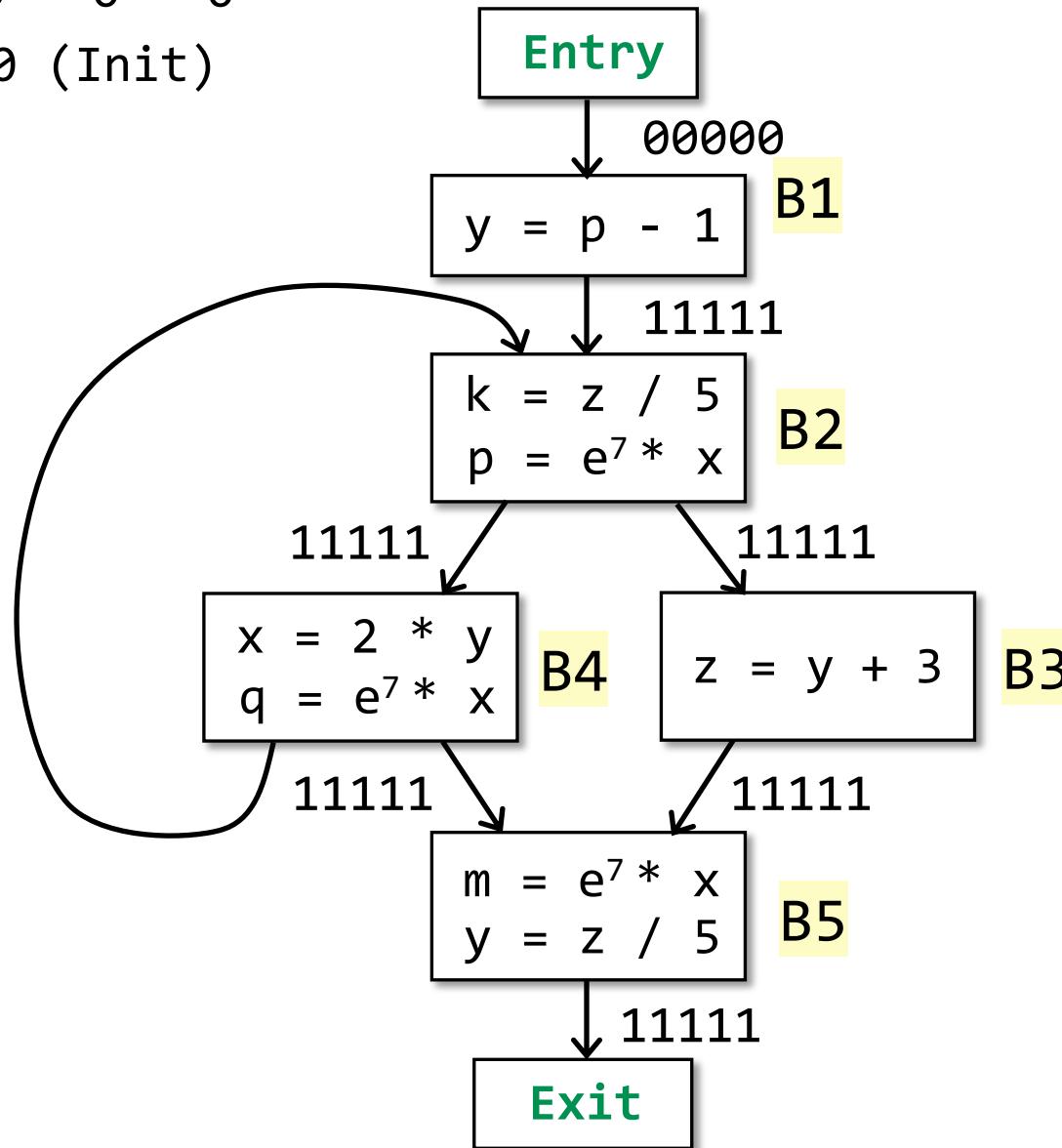
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            OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
        }
```



p-1 z/5 2\*y e<sup>7</sup>\*x y+3  
0 0 0 0 0

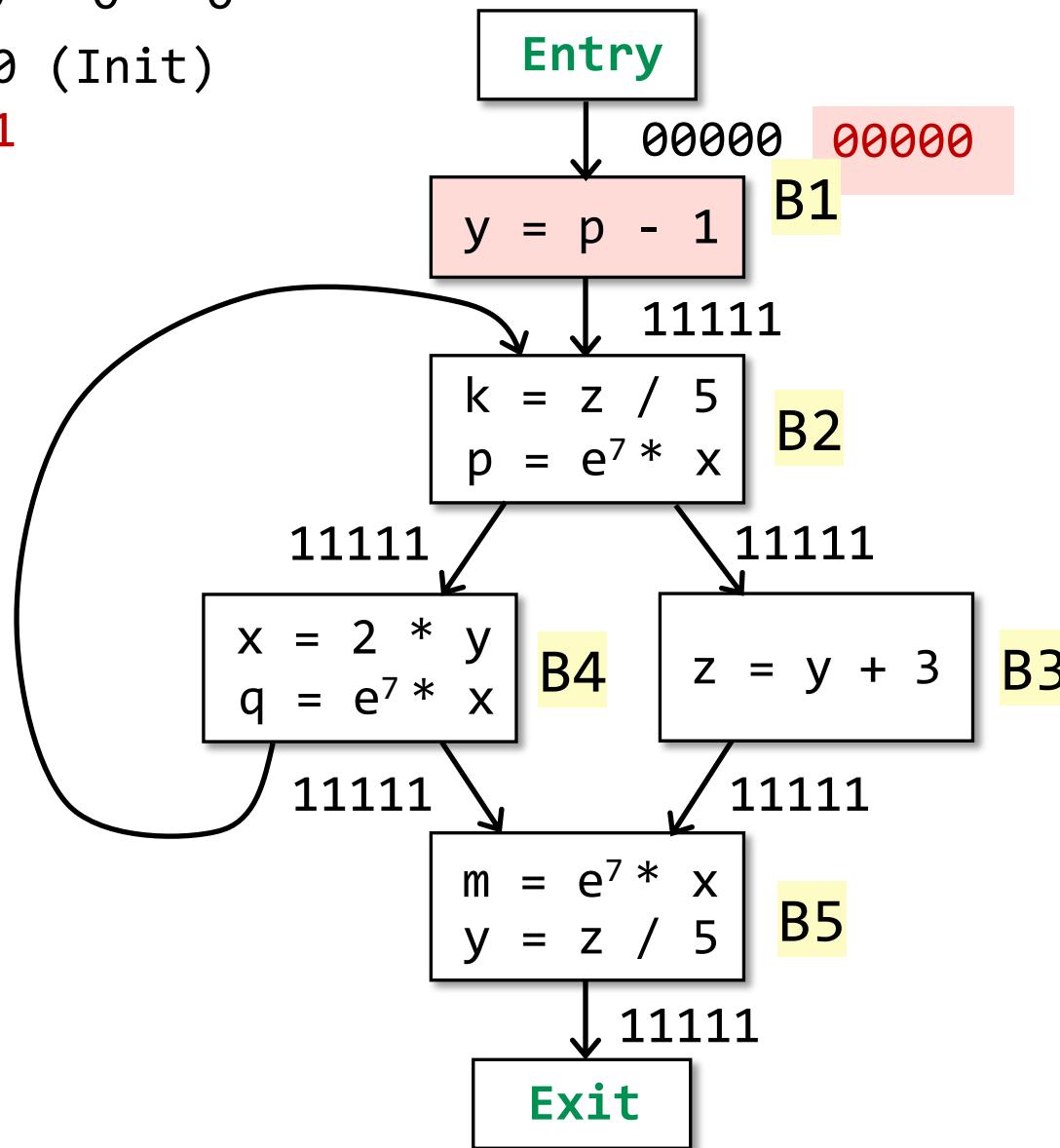
Iteration 0 (Init)



p-1 z/5 2\*y e<sup>7</sup>\*x y+3  
0 0 0 0 0

Iteration 0 (Init)

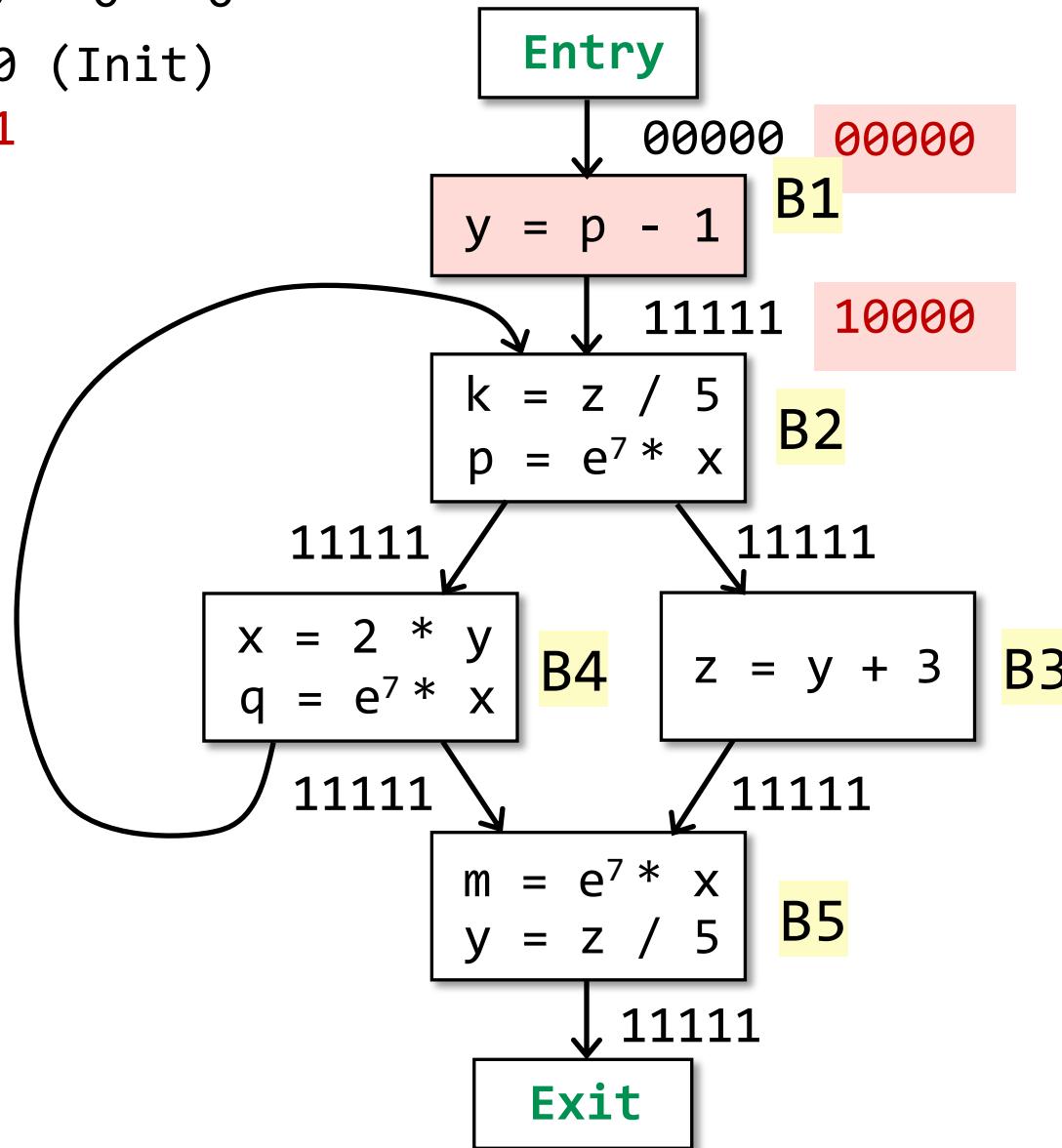
Iteration 1



p-1 z/5 2\*y e<sup>7</sup>\*x y+3  
0 0 0 0 0

Iteration 0 (Init)

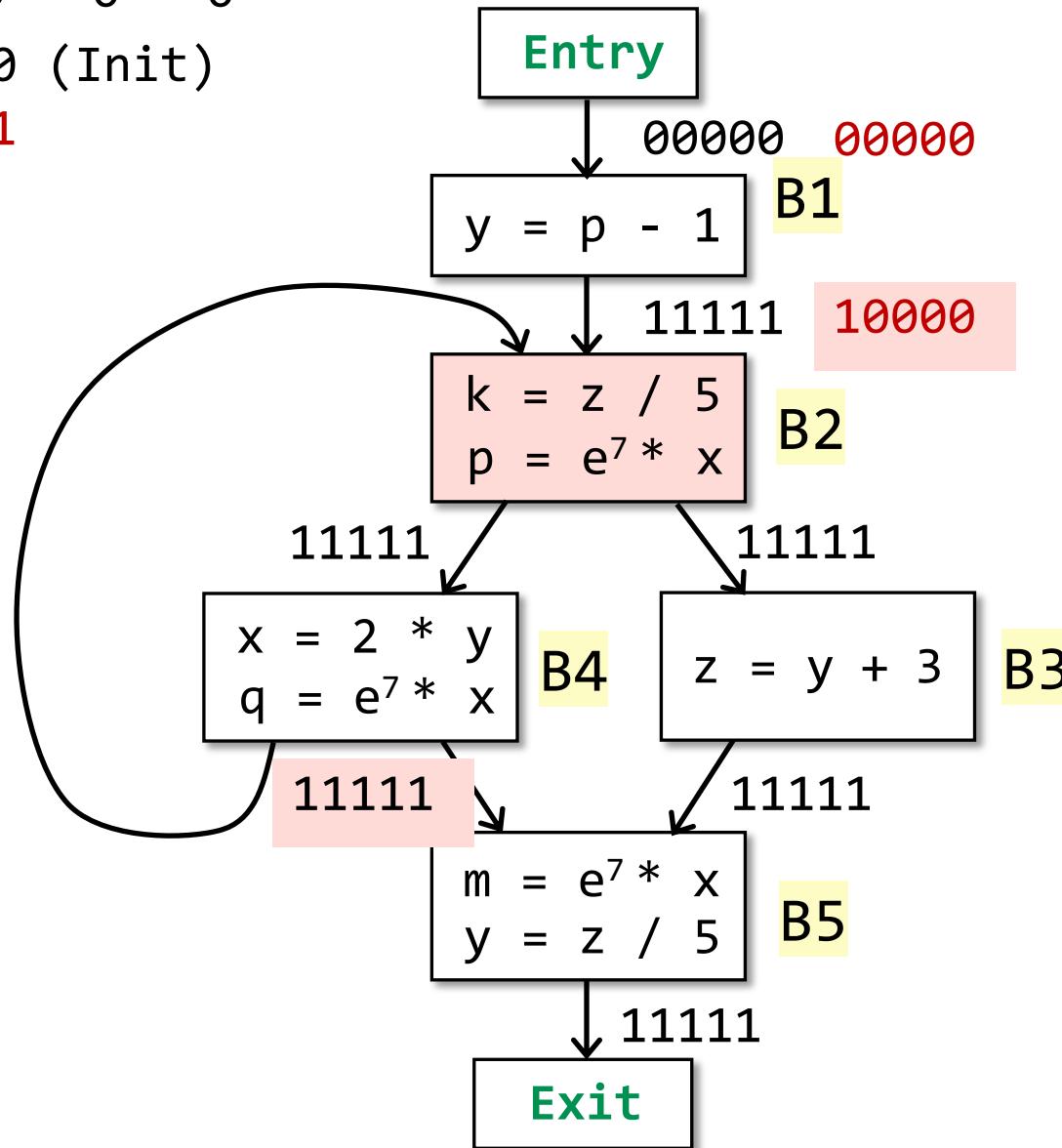
Iteration 1



p-1 z/5 2\*y e<sup>7</sup>\*x y+3  
0 0 0 0 0

Iteration 0 (Init)

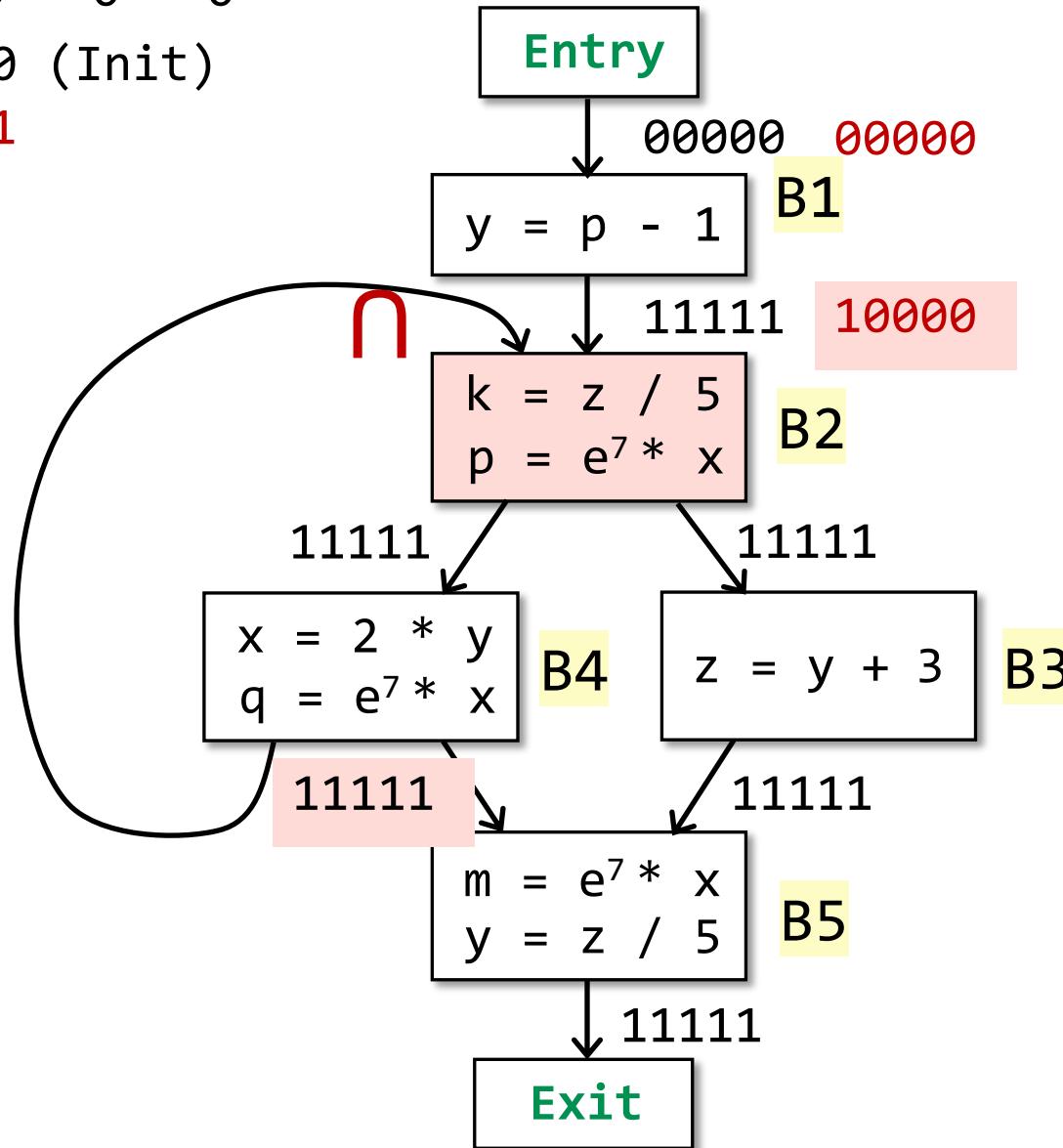
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
0 0 0 0 0

Iteration 0 (Init)

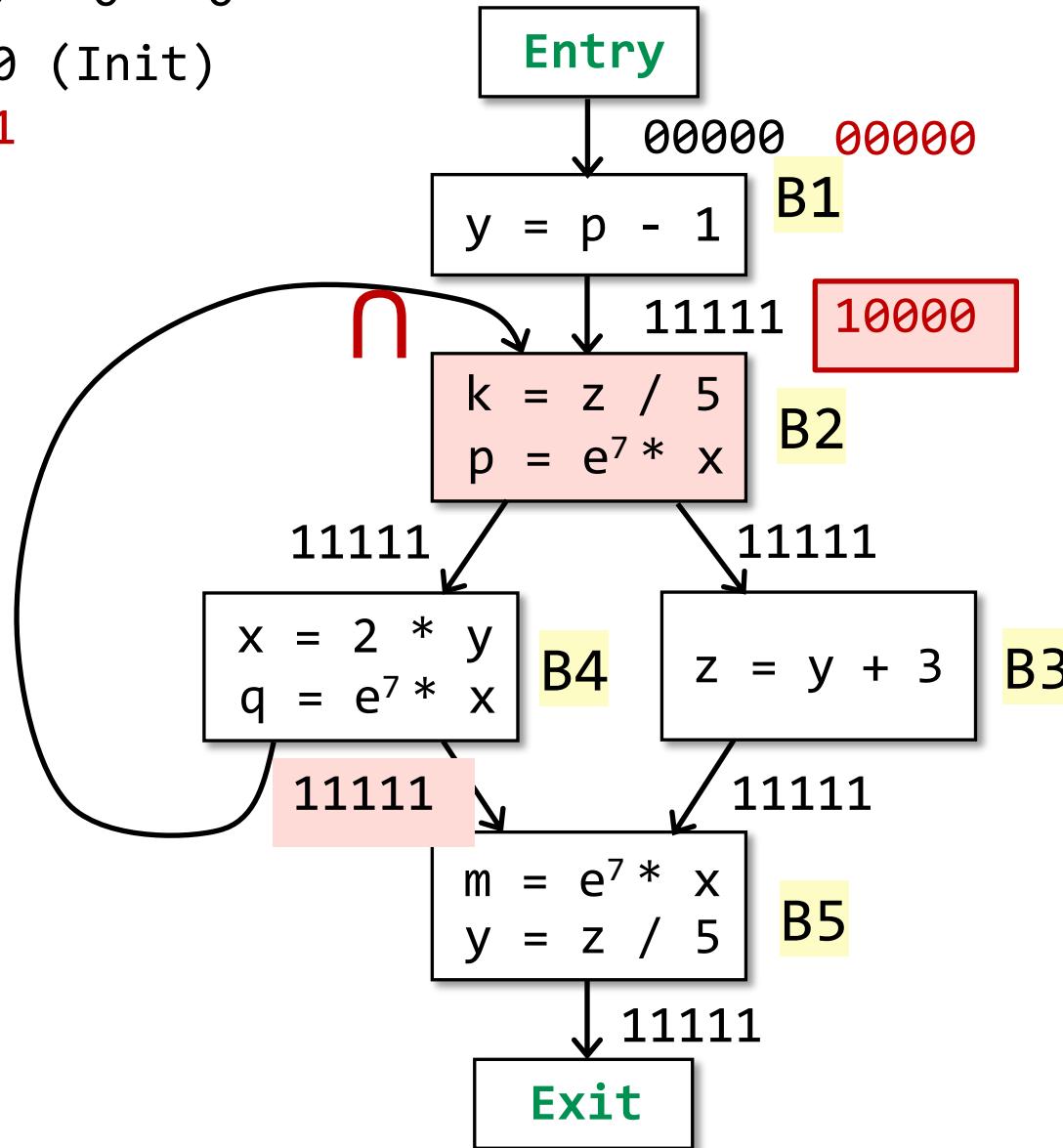
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

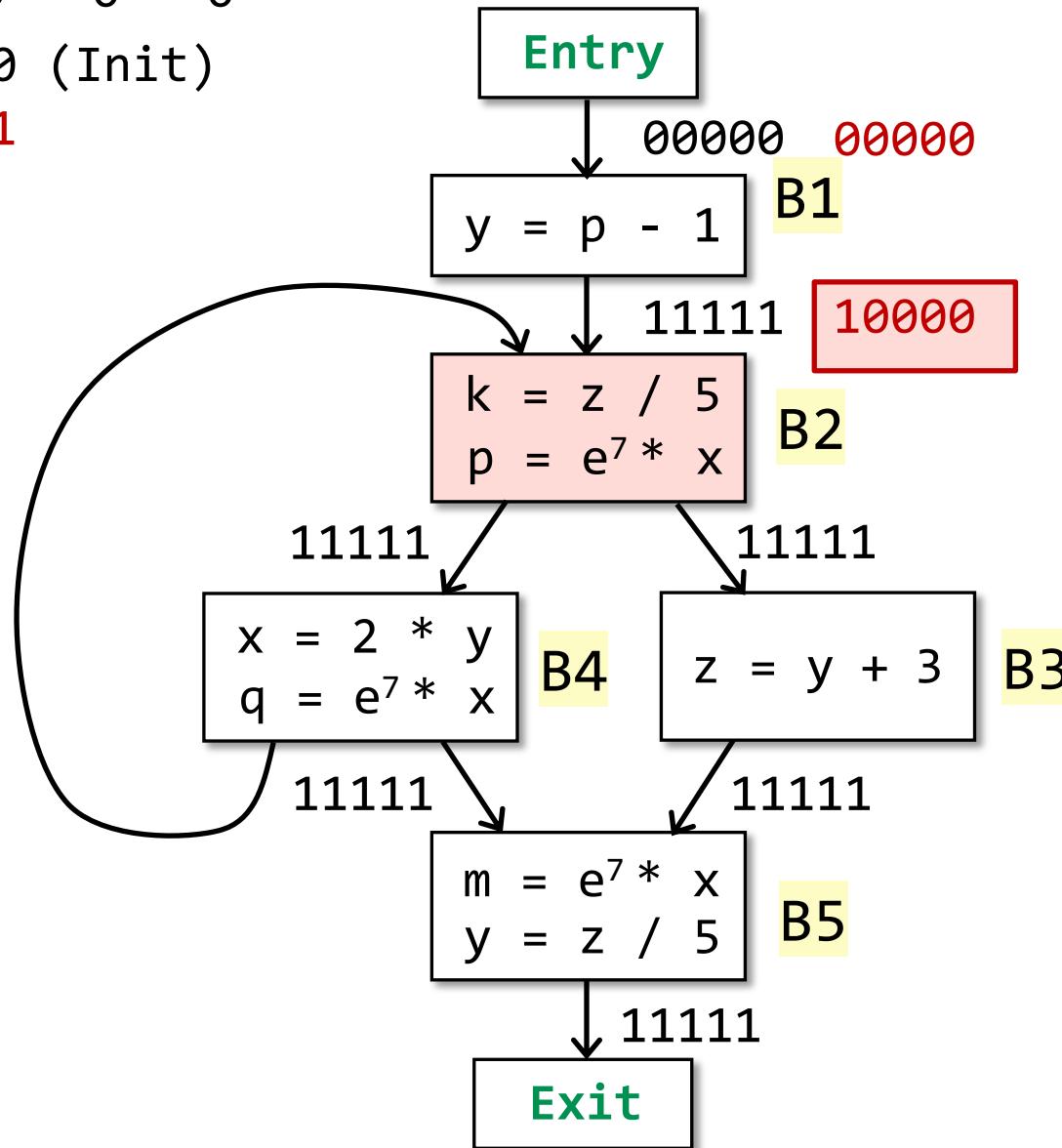
Iteration 1



p-1 z/5 2\*y e<sup>7</sup>\*x y+3  
0 0 0 0 0

Iteration 0 (Init)

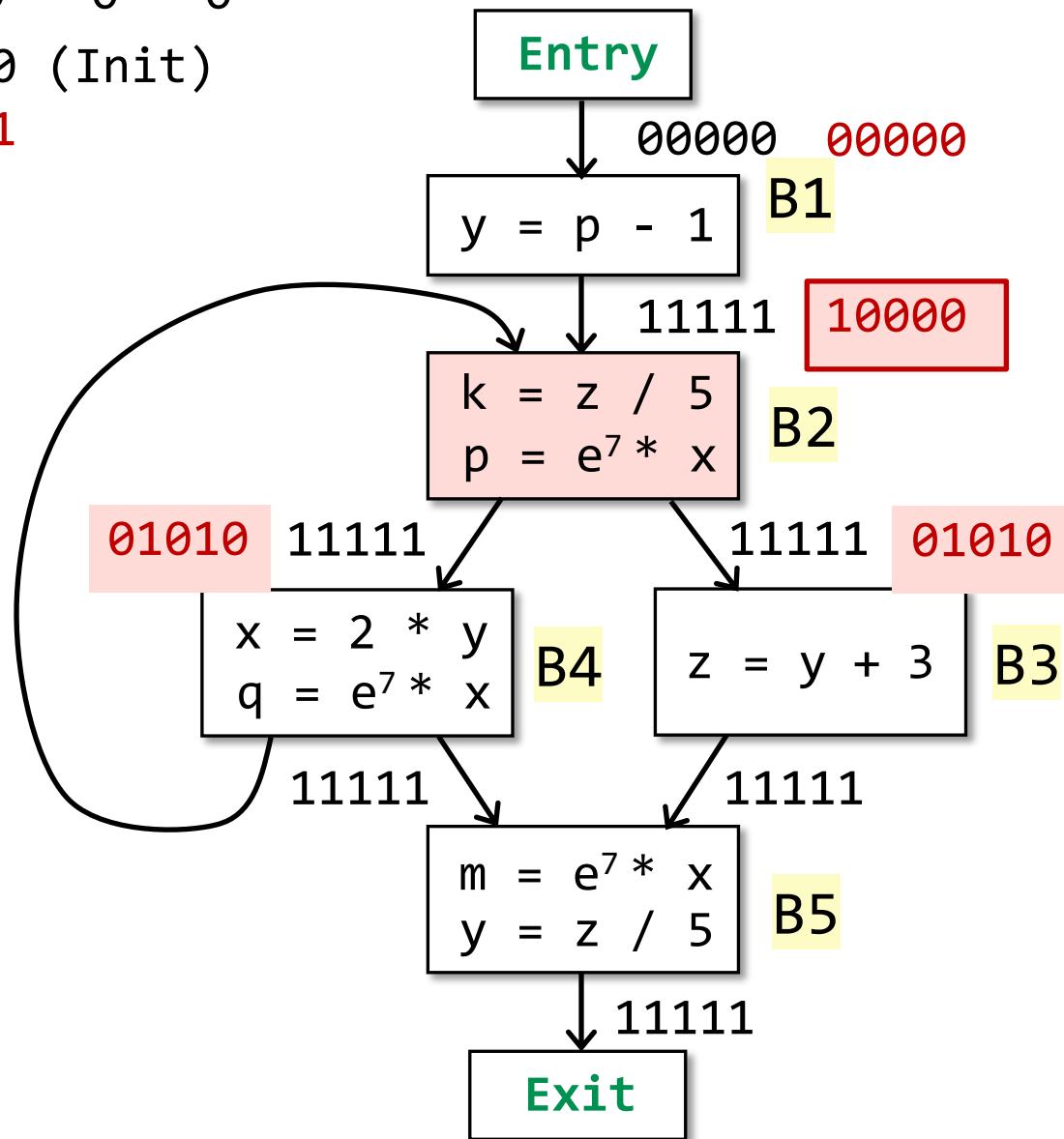
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
0 0 0 0 0

Iteration 0 (Init)

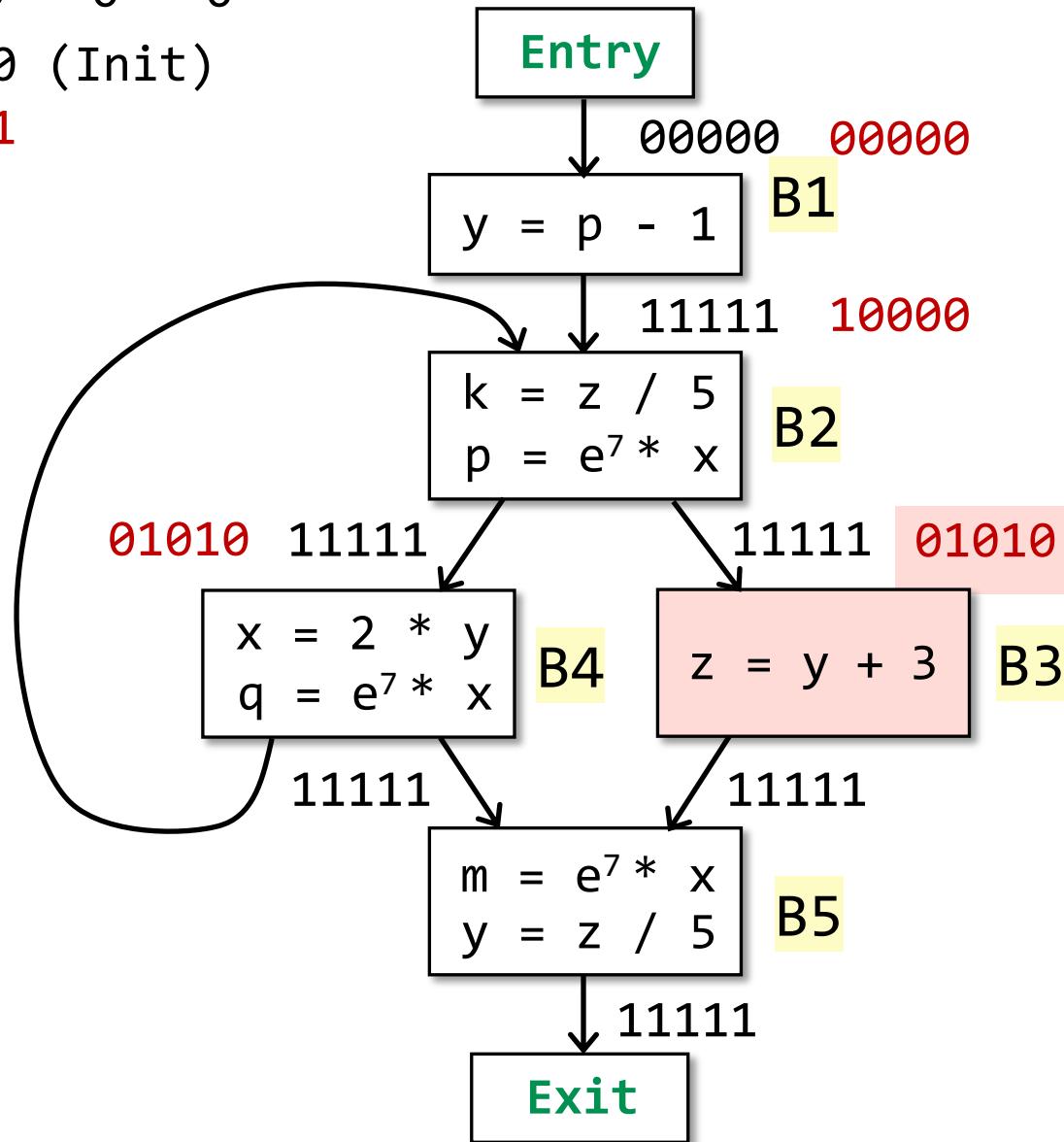
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
0 0 0 0 0

Iteration 0 (Init)

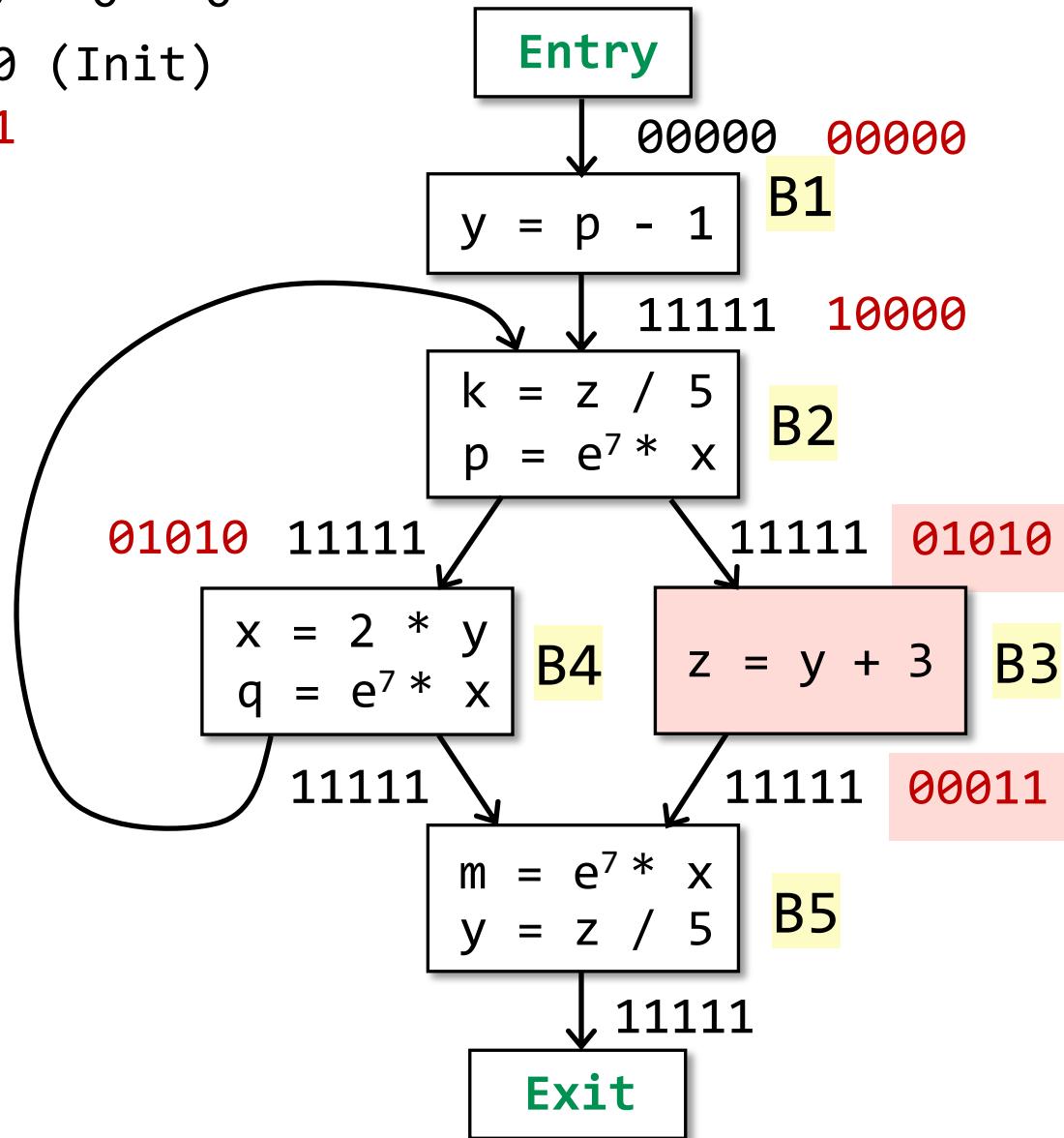
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

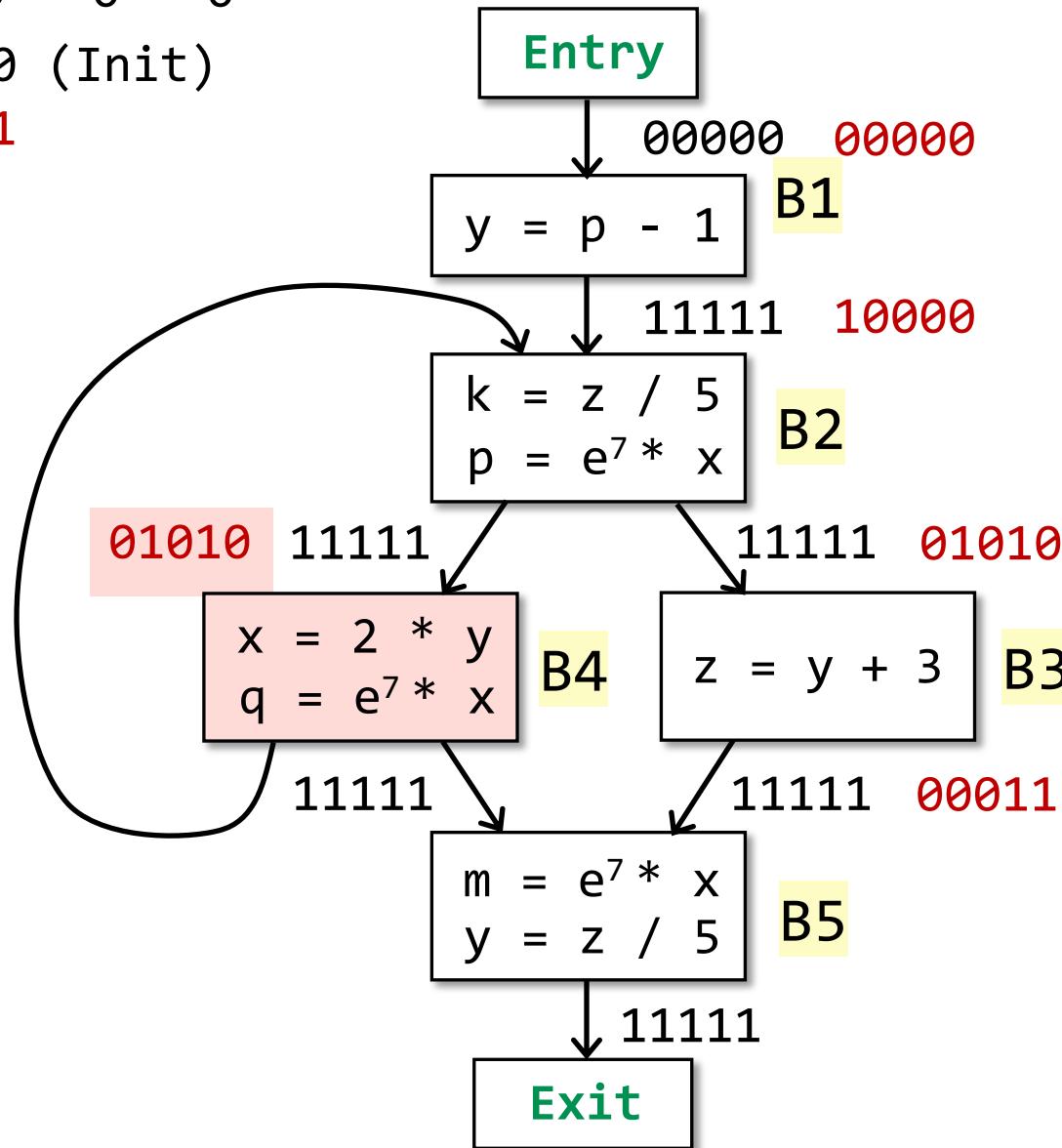
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

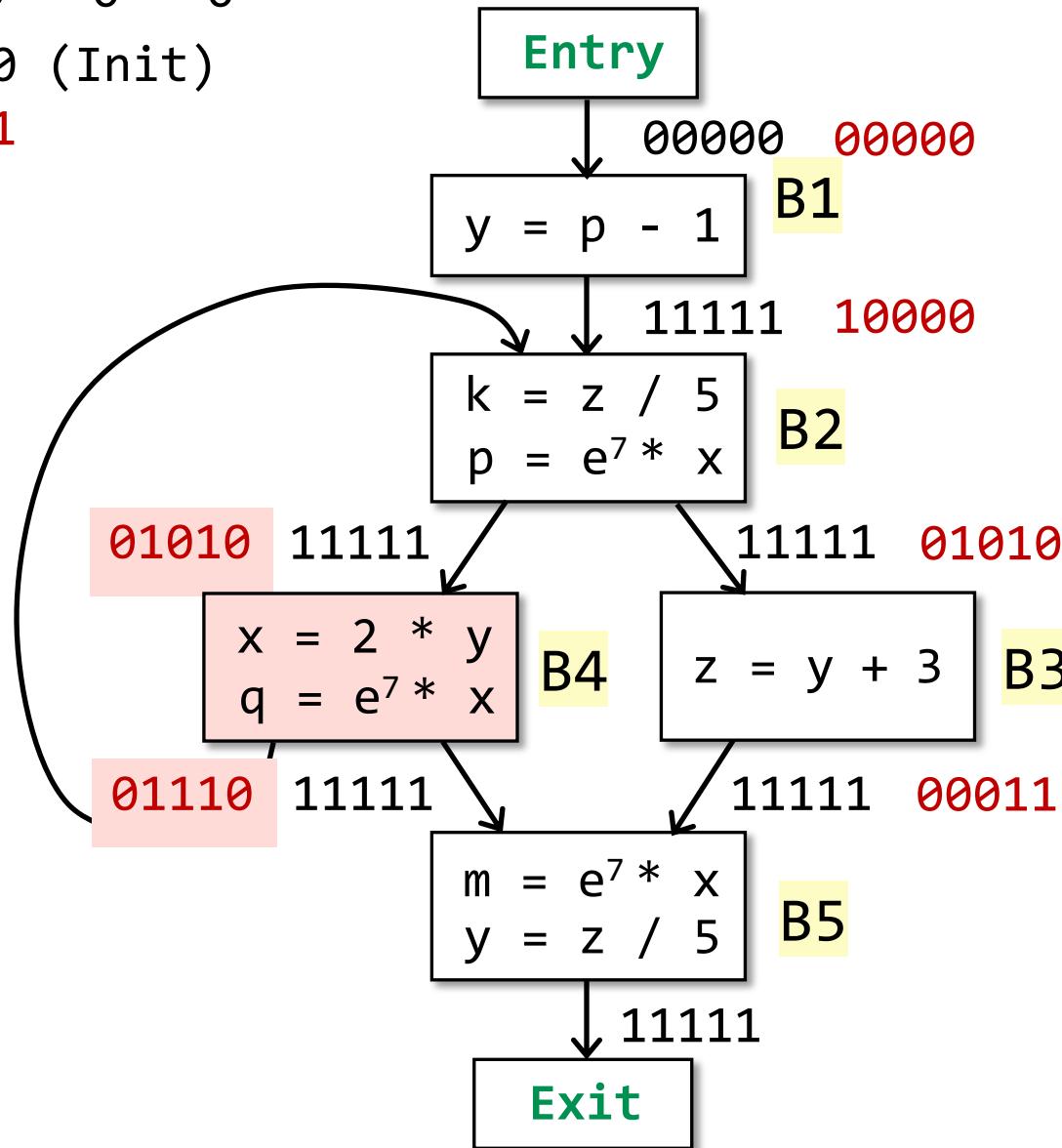
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

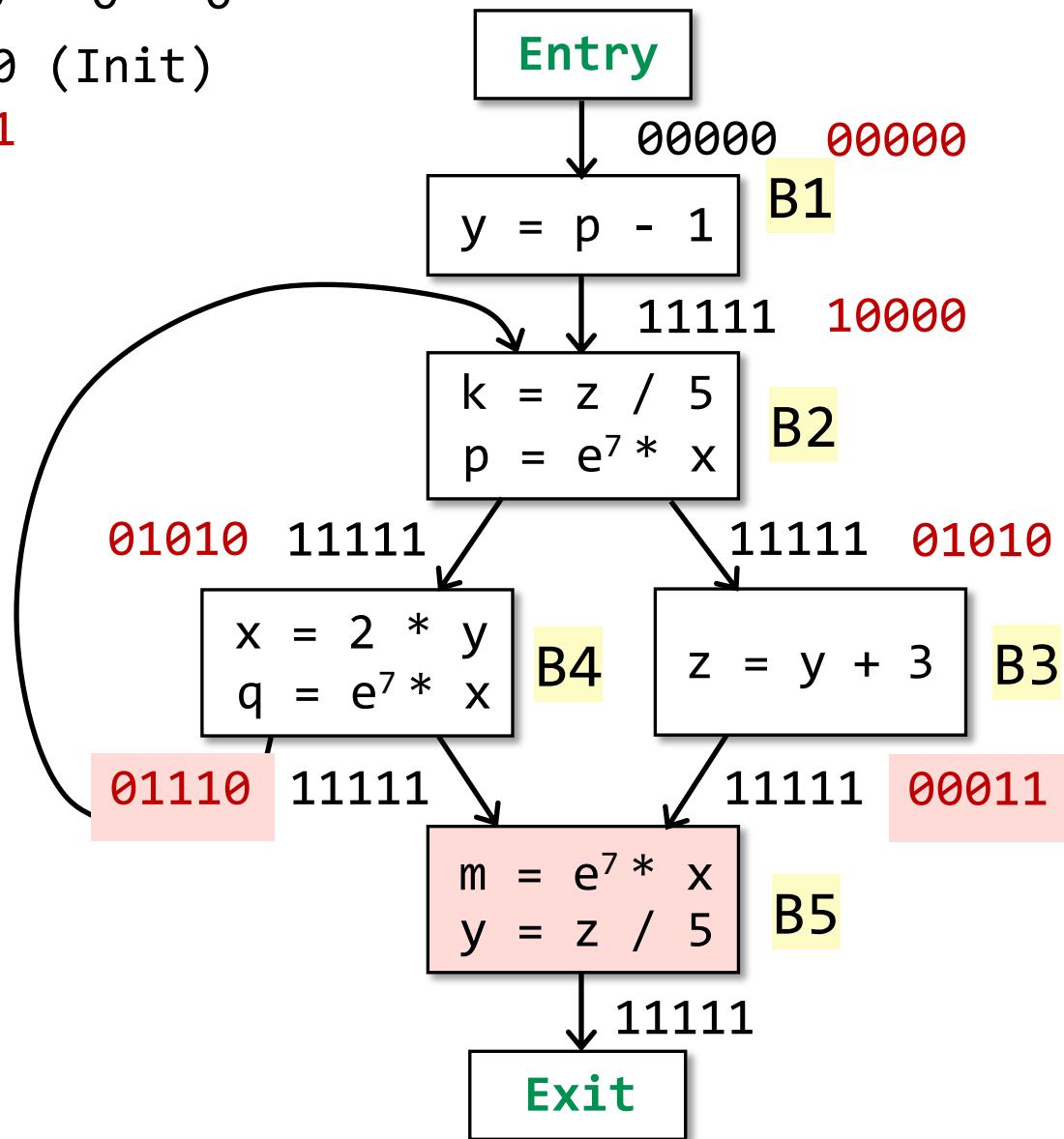
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

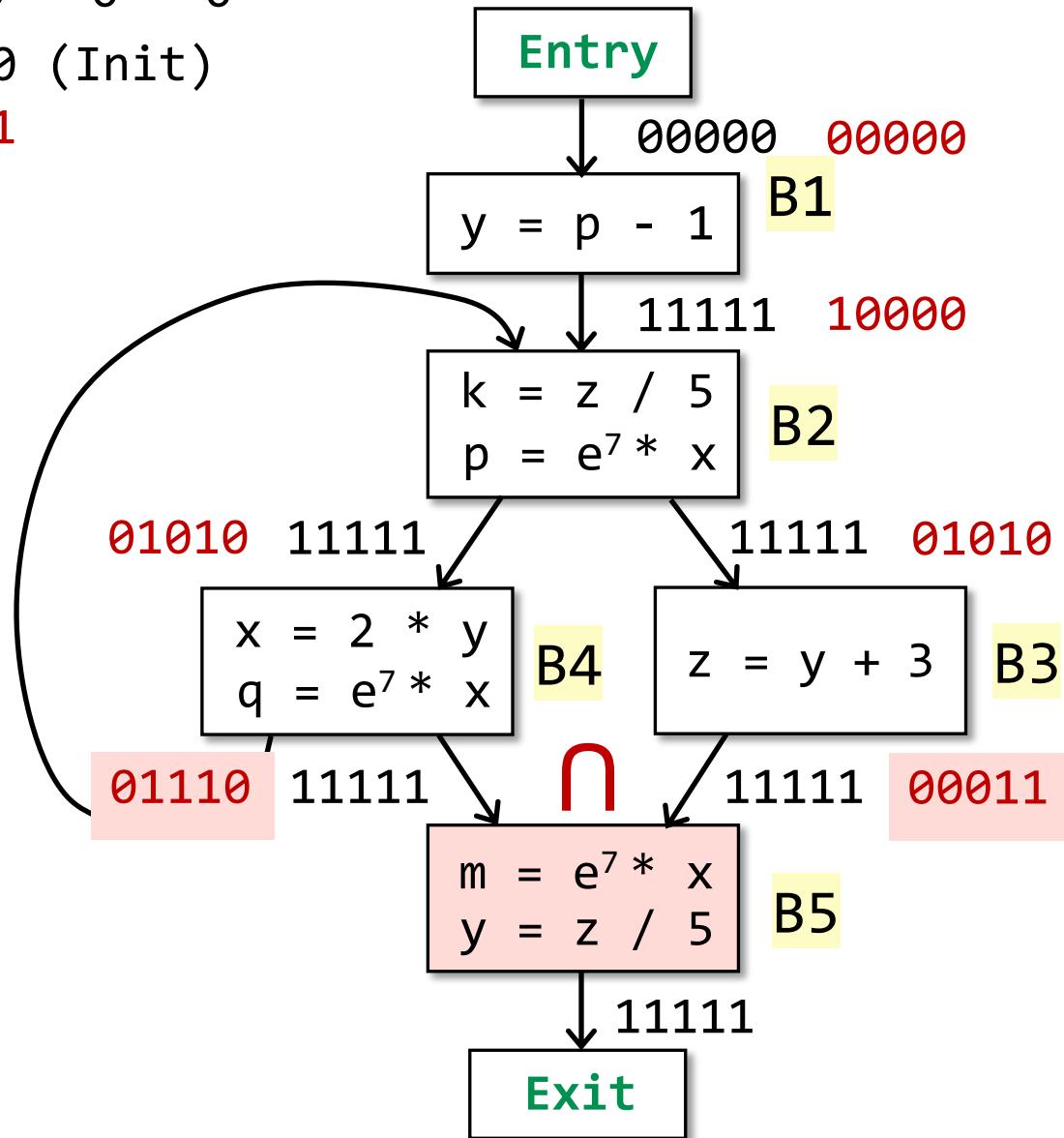
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

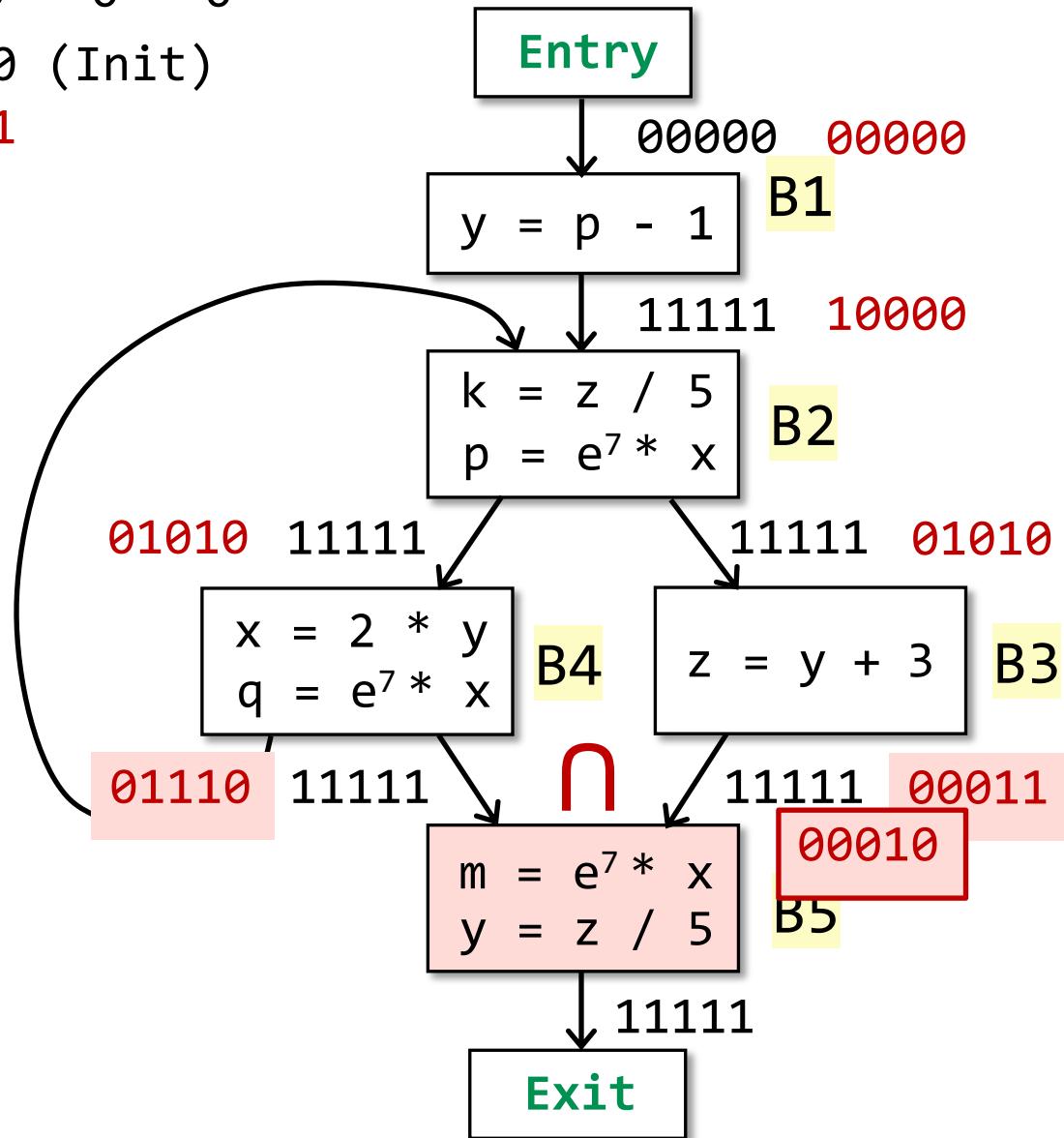
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

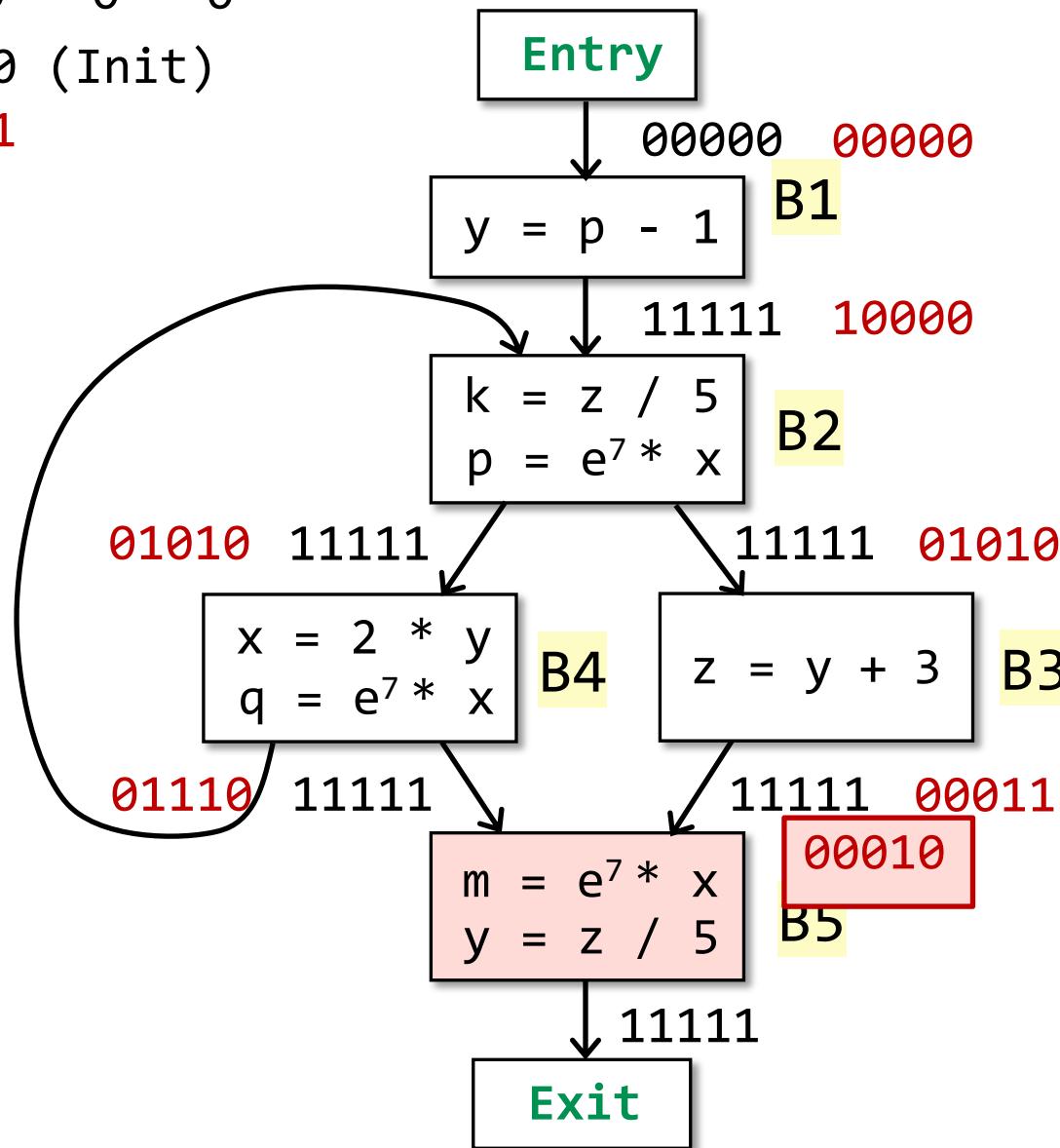
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

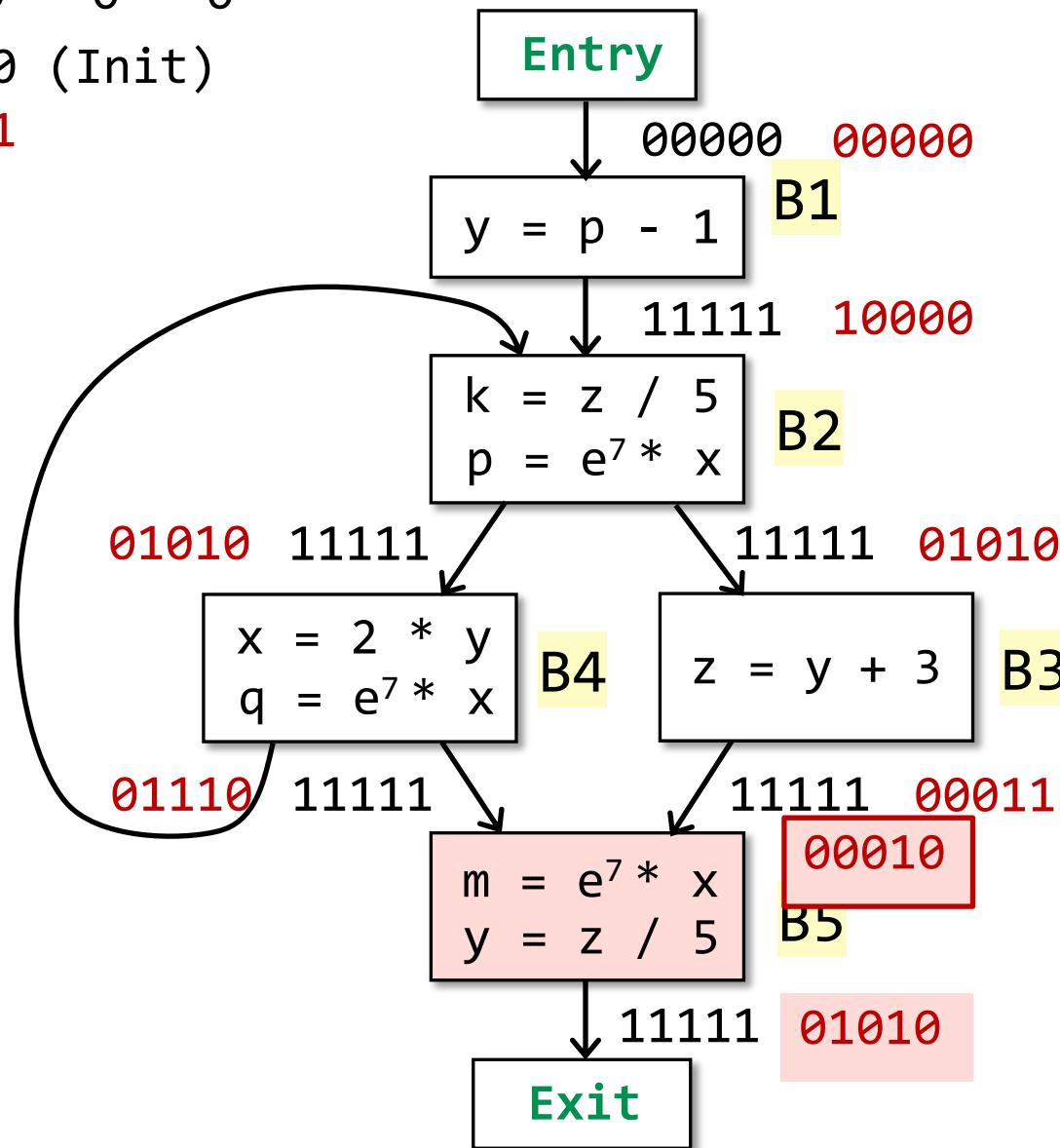
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

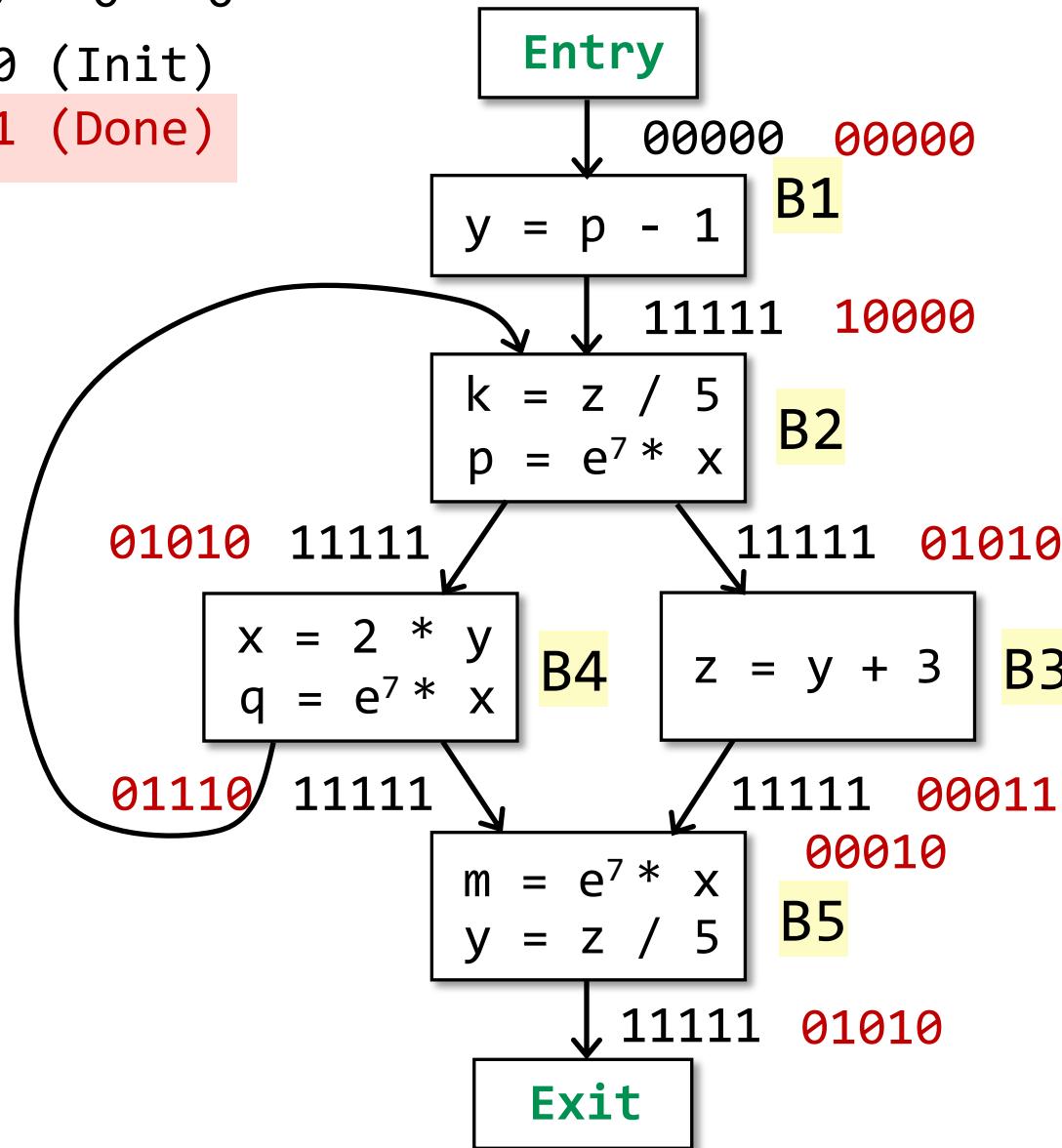
Iteration 1



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
0 0 0 0 0

Iteration 0 (Init)

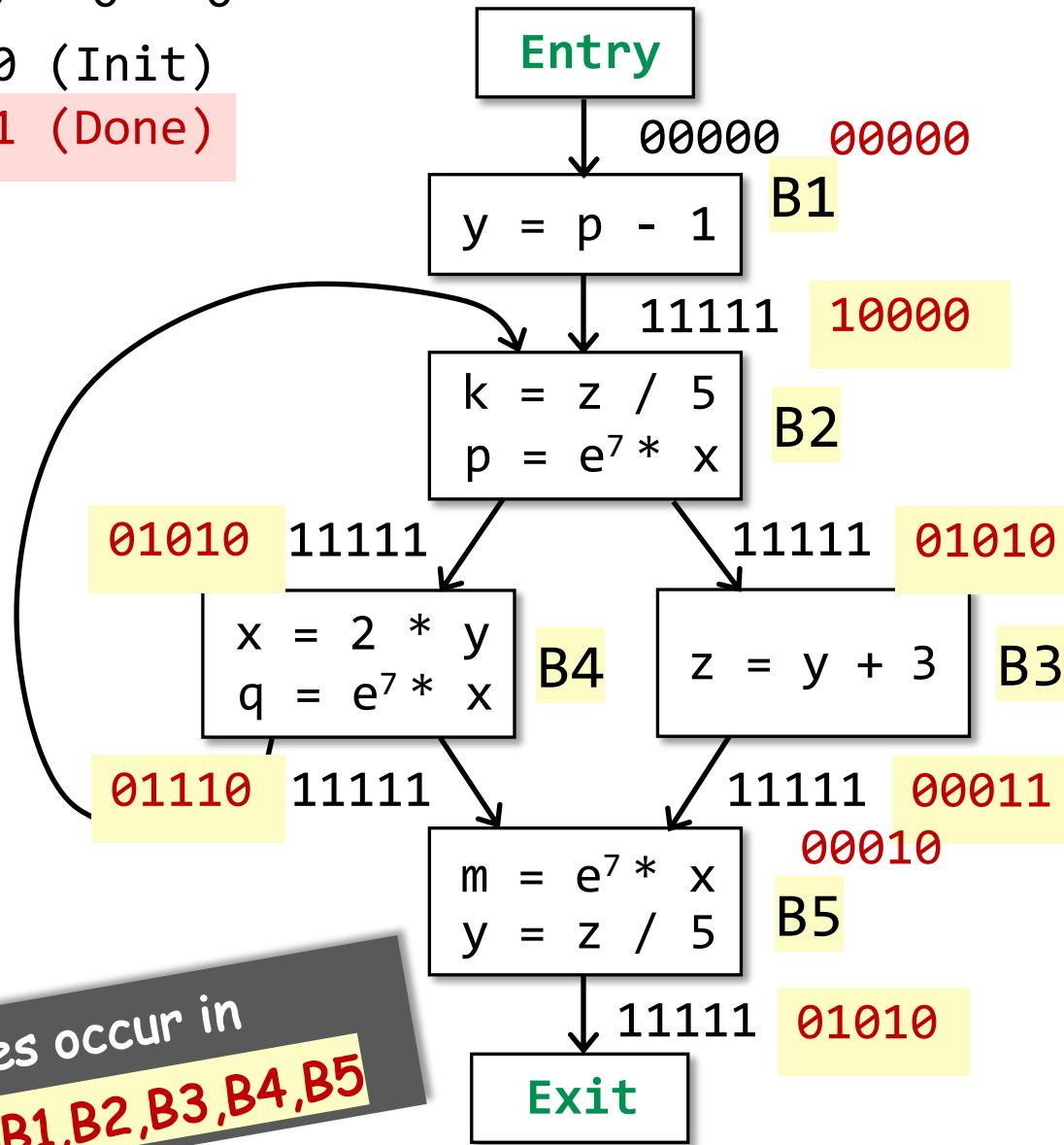
Iteration 1 (Done)



$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)



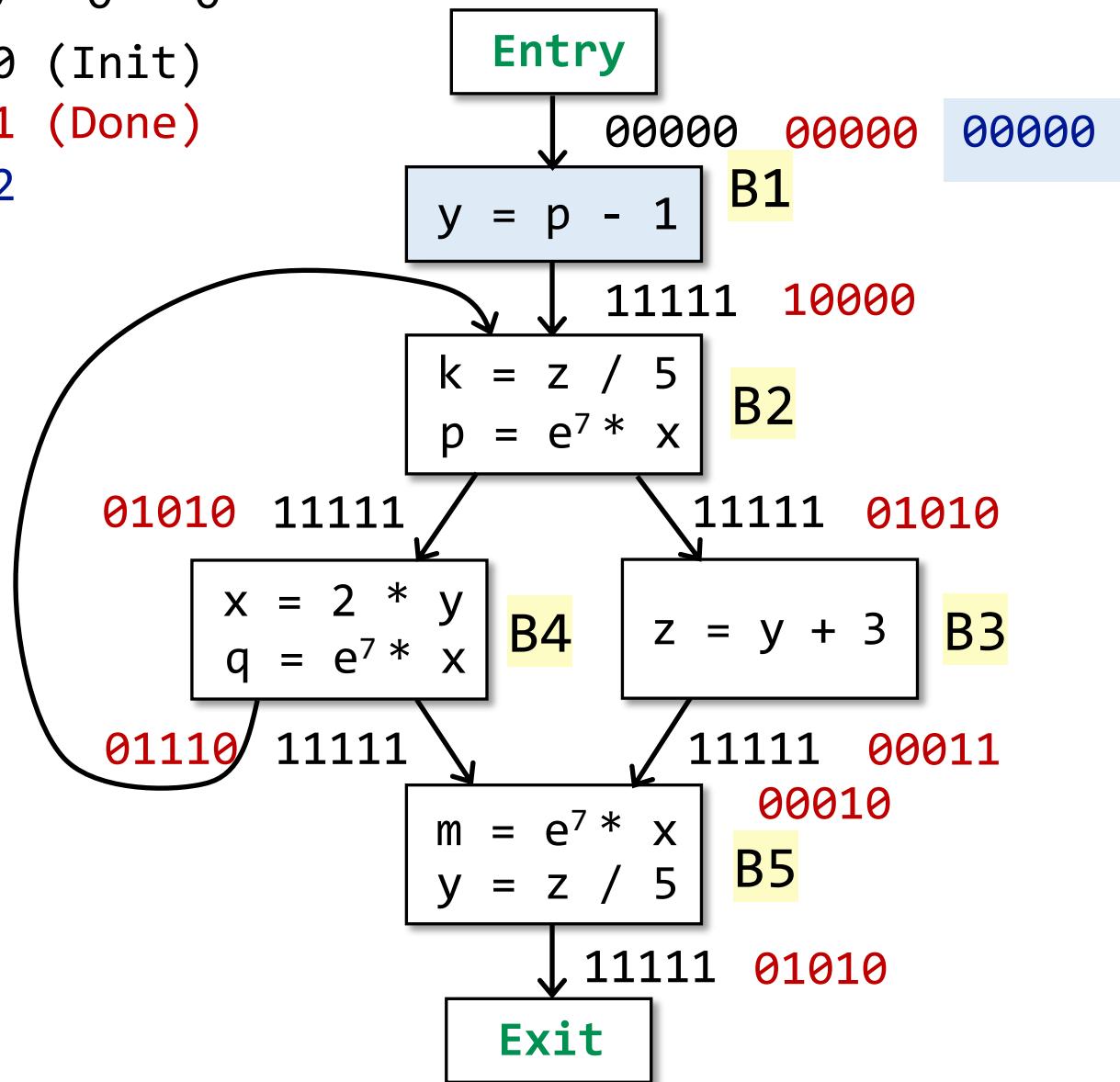
Changes occur in  
 OUT[] of B1,B2,B3,B4,B5

$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

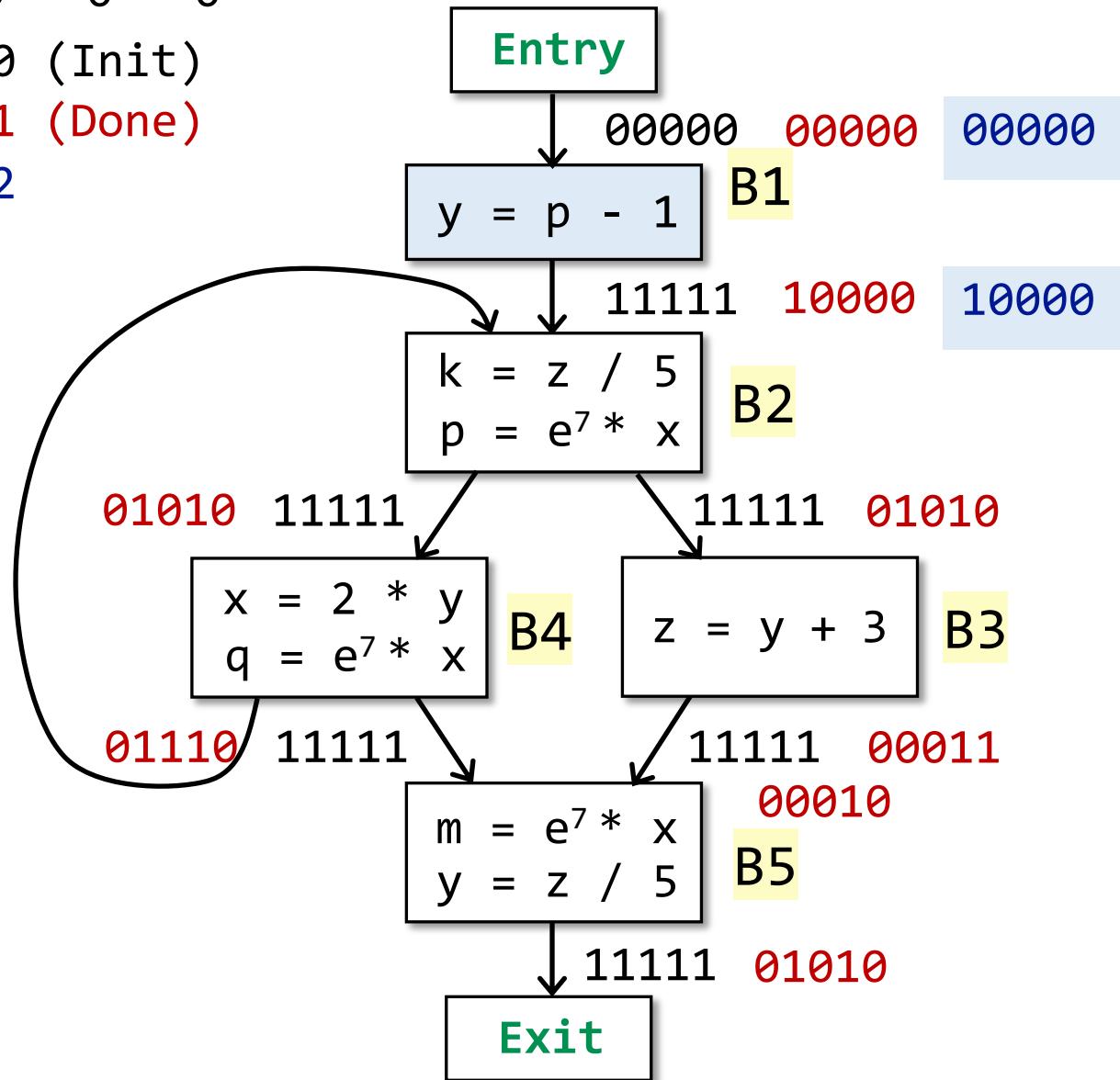


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

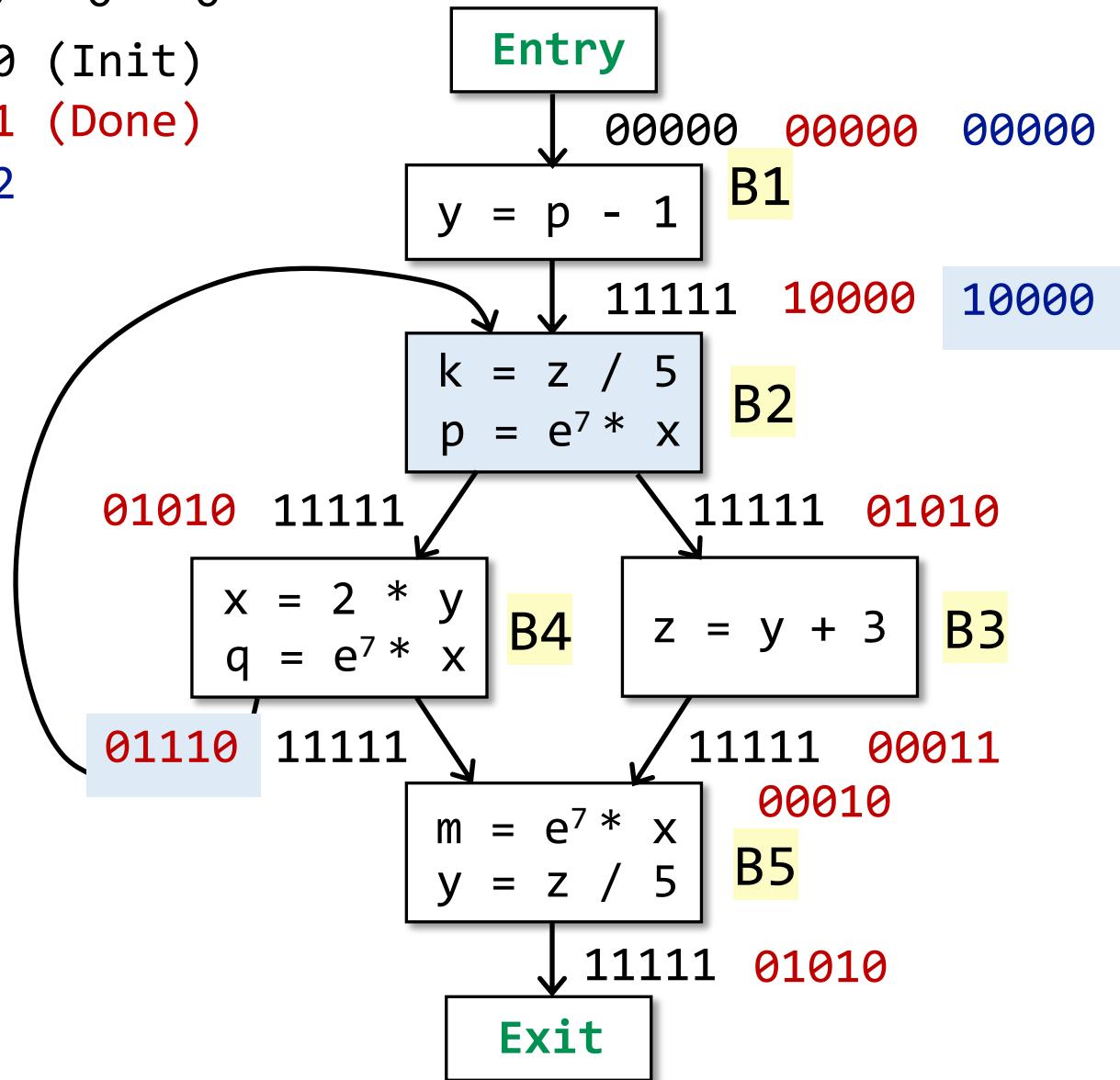


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

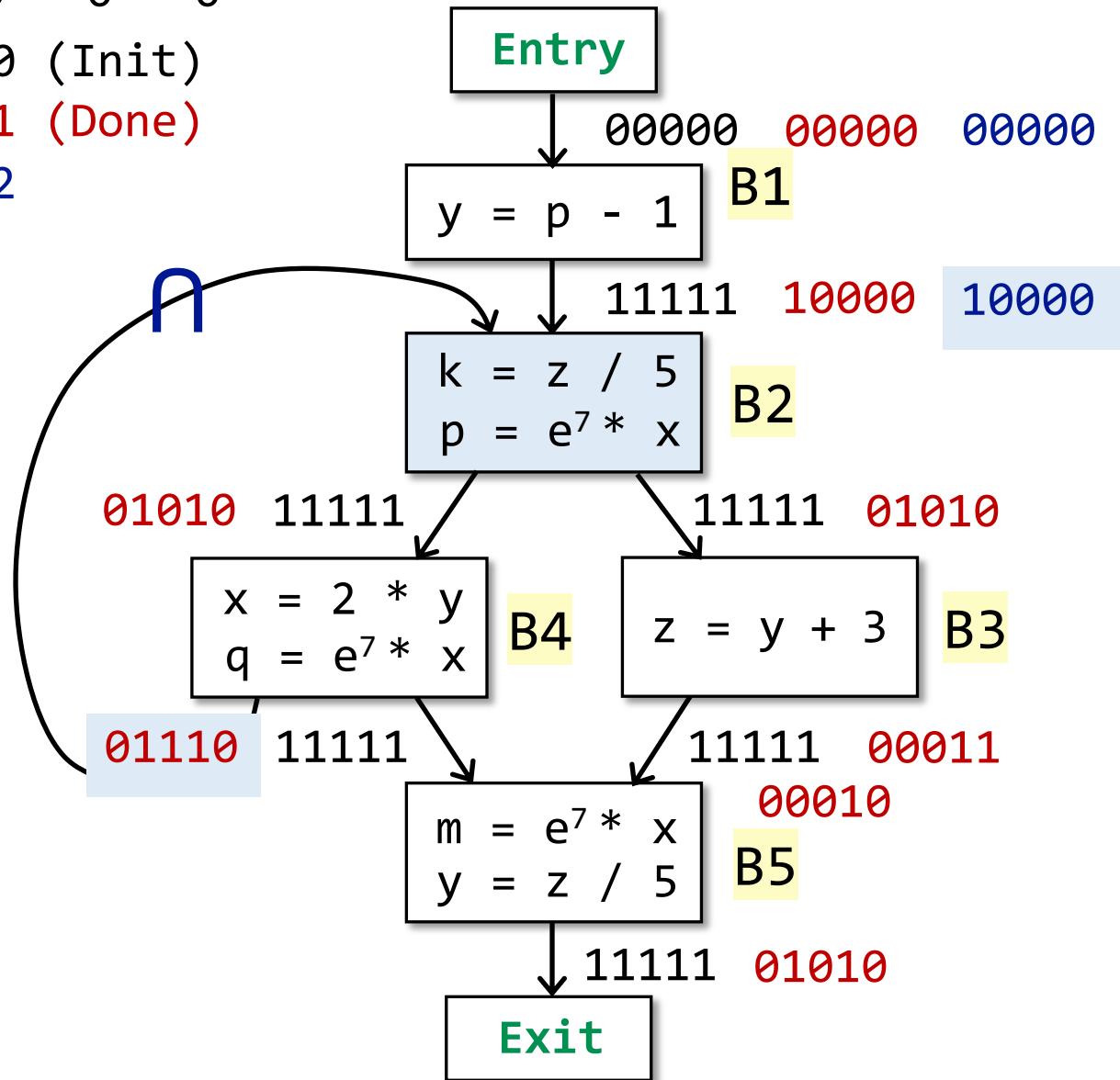


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

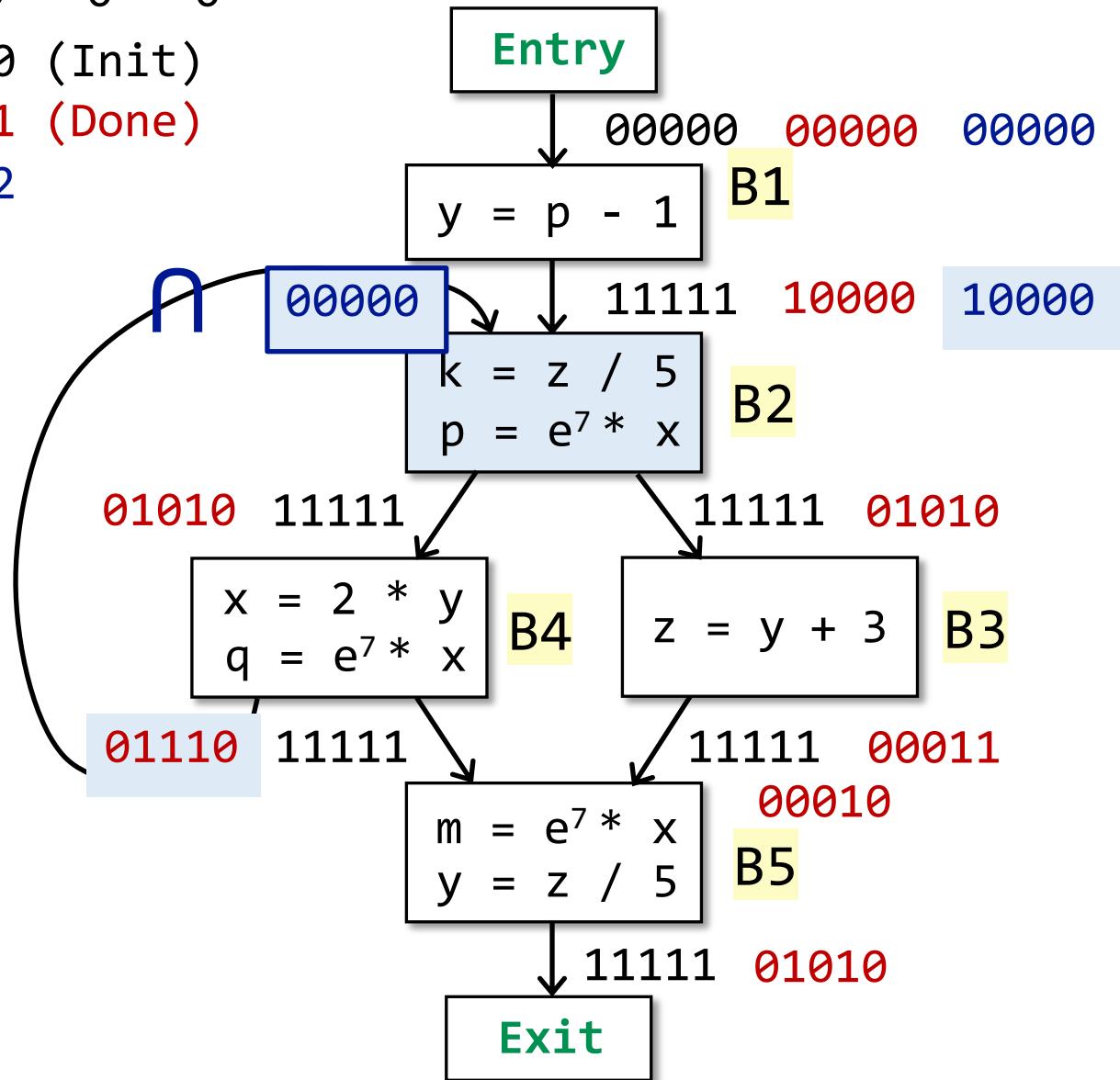


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

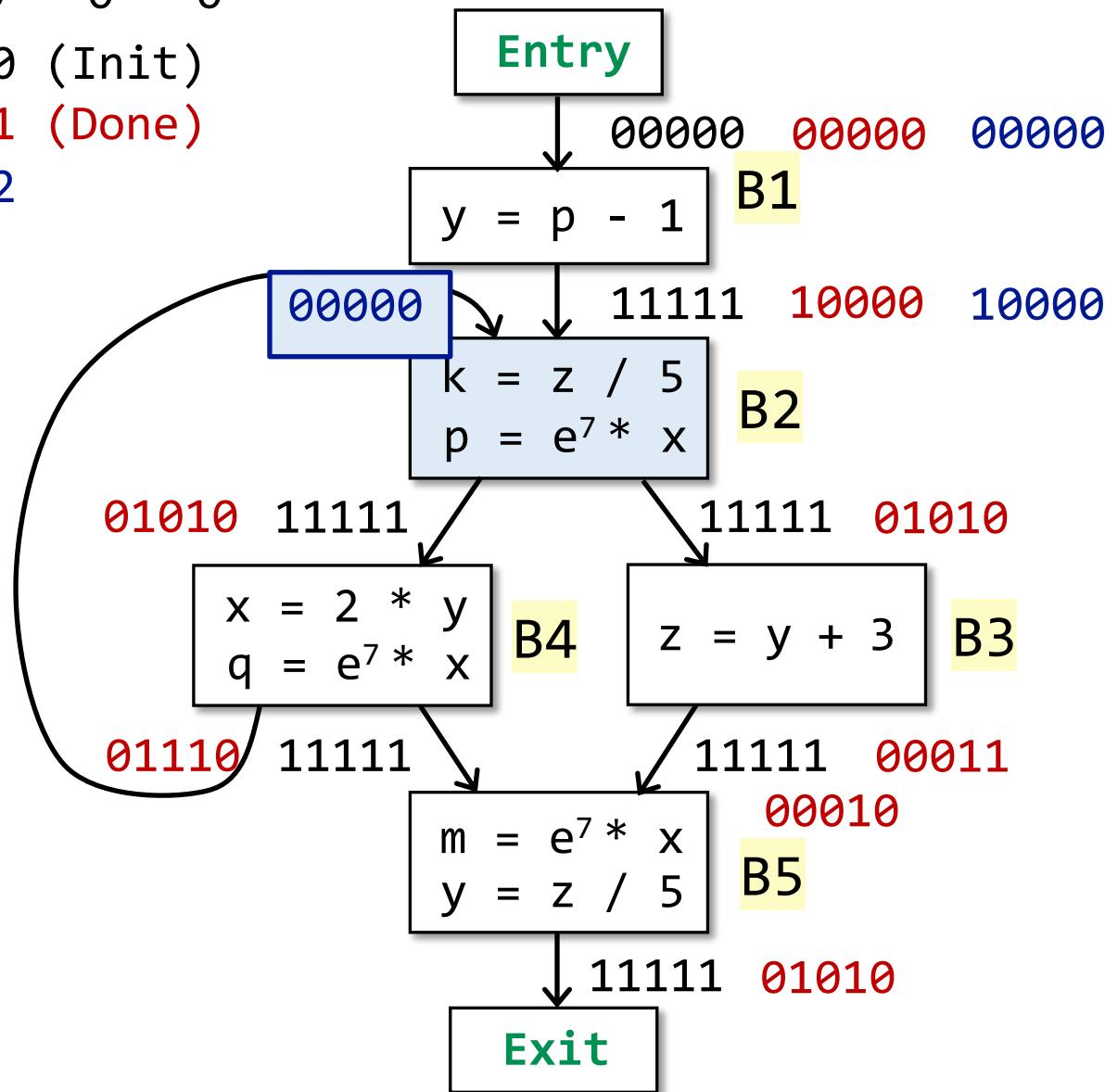


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

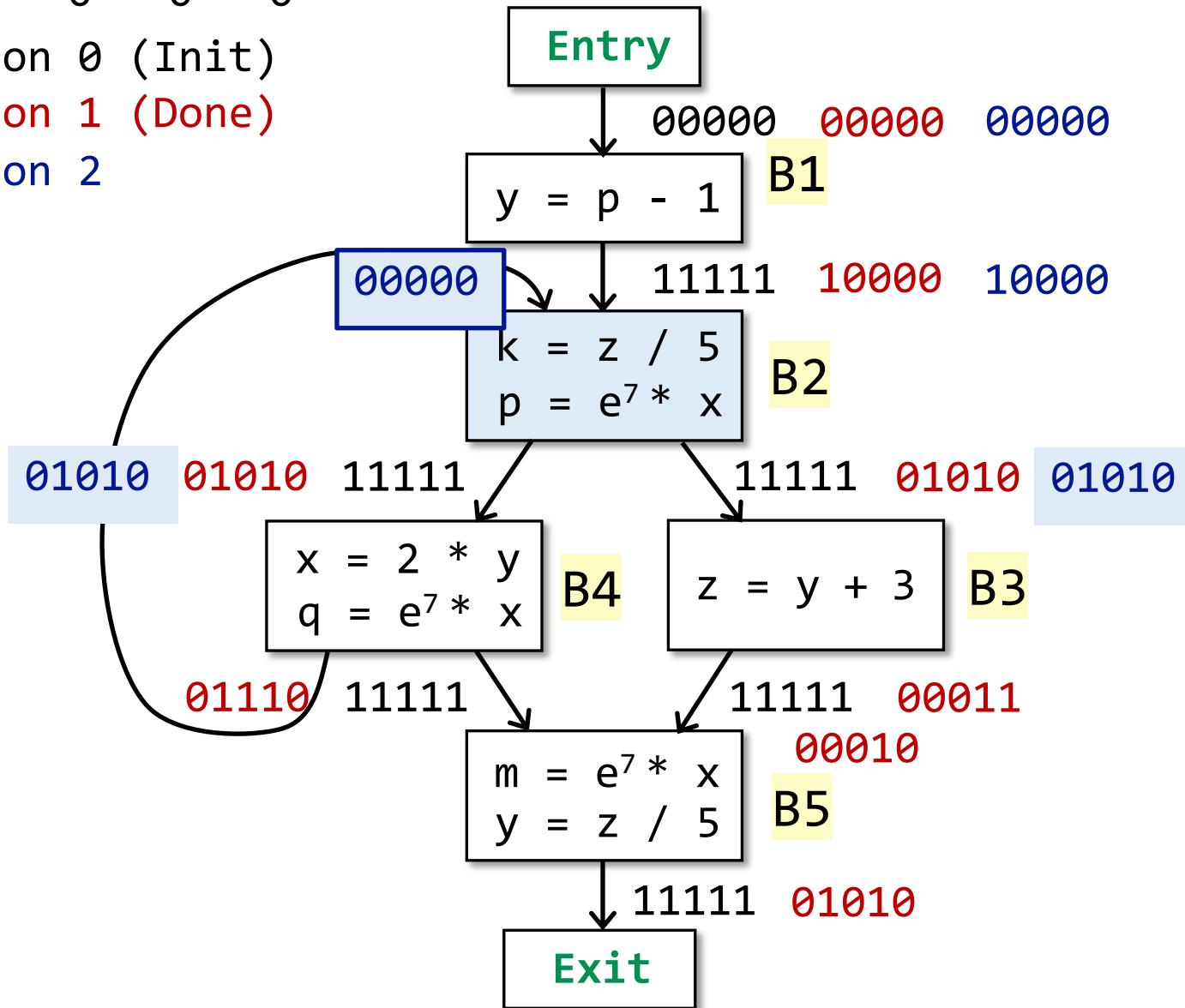


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

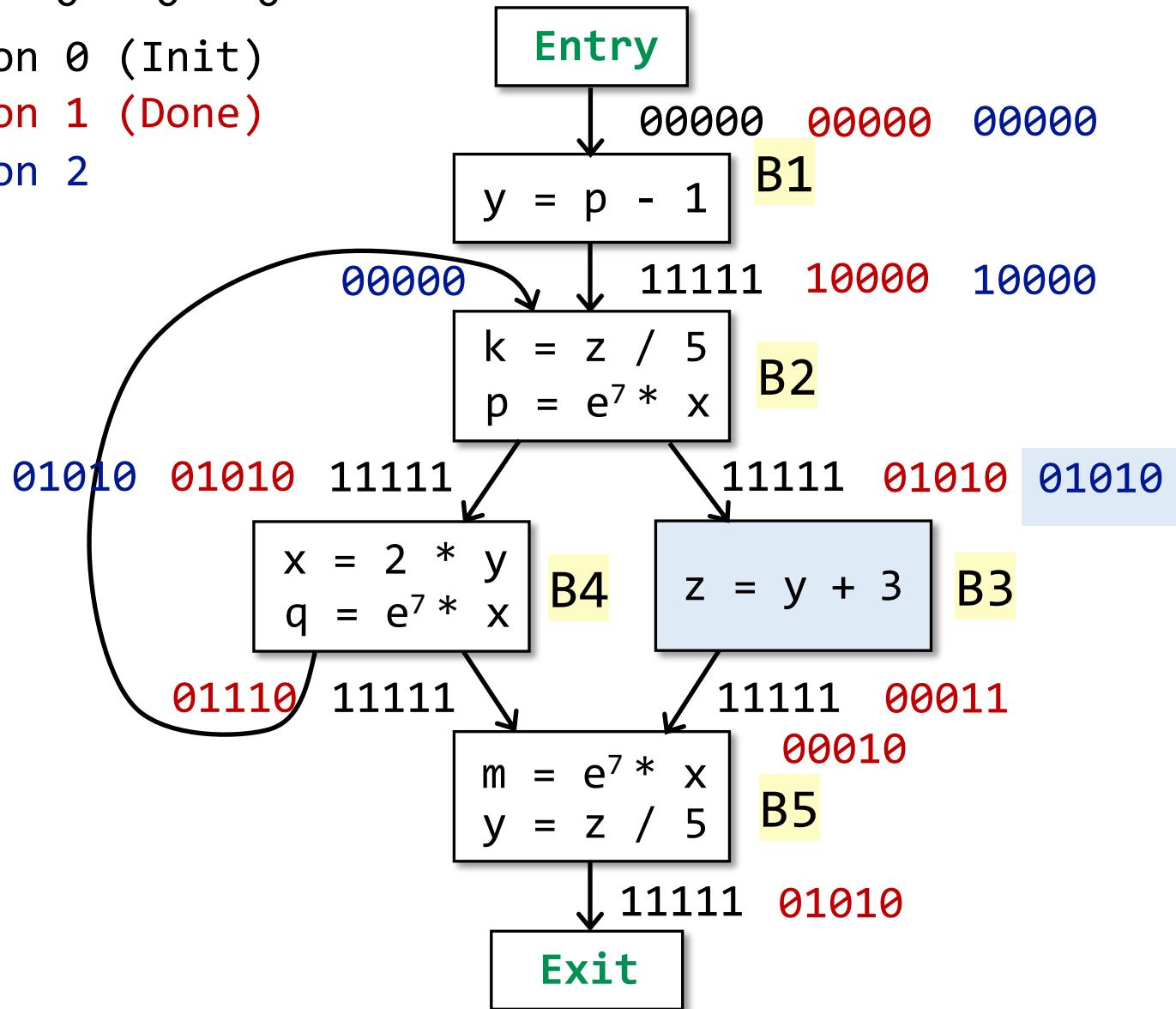


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

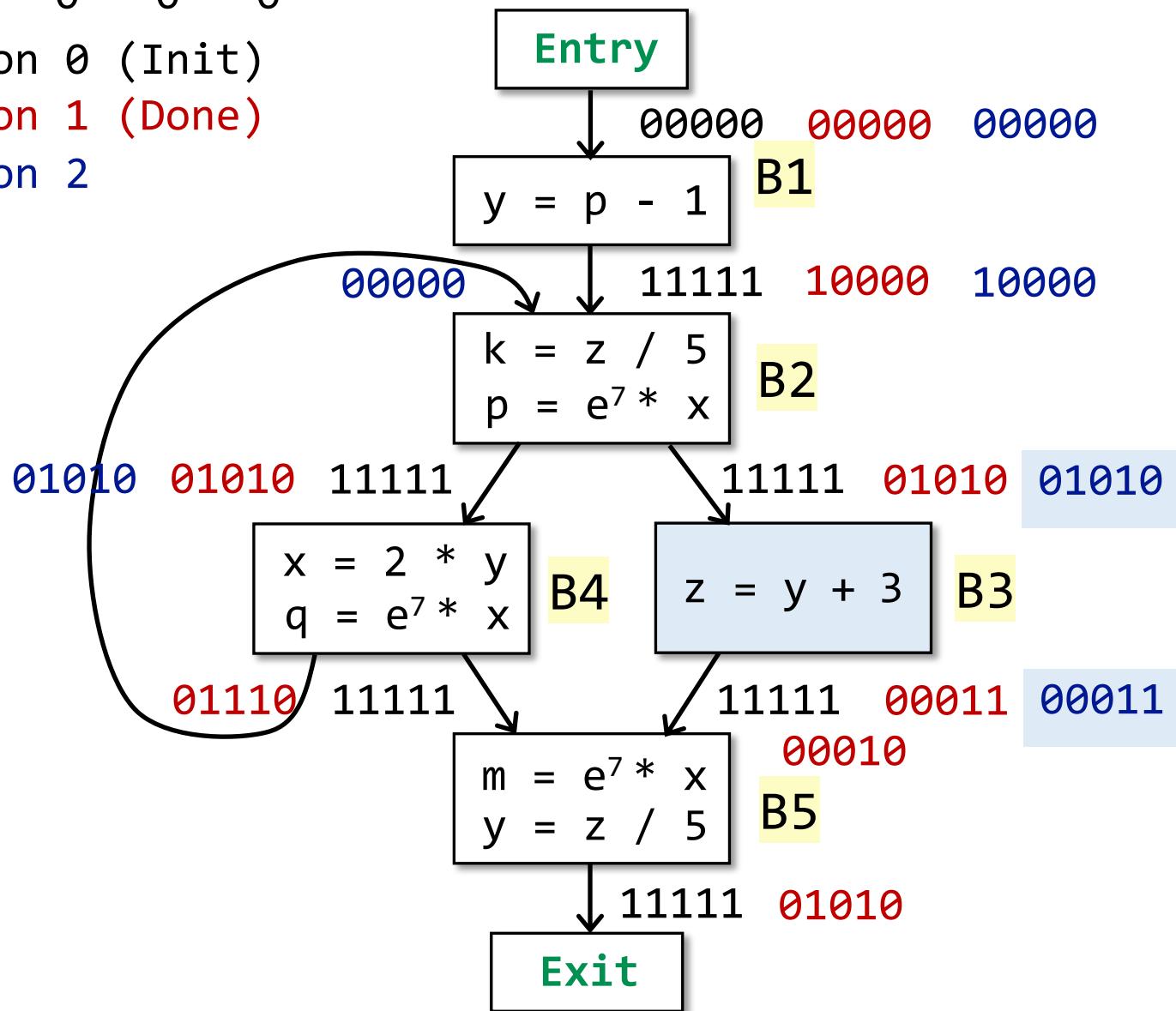


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

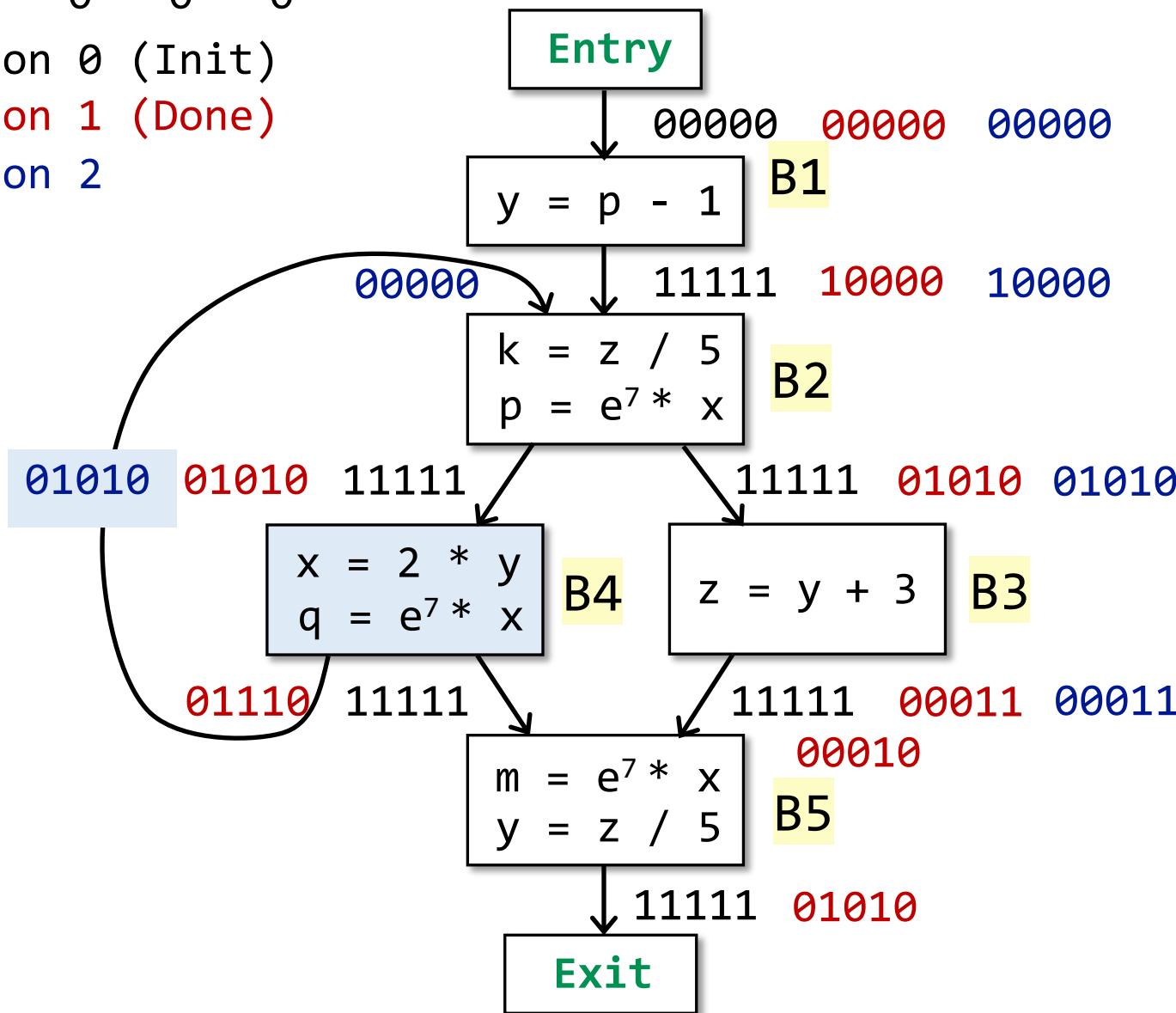


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

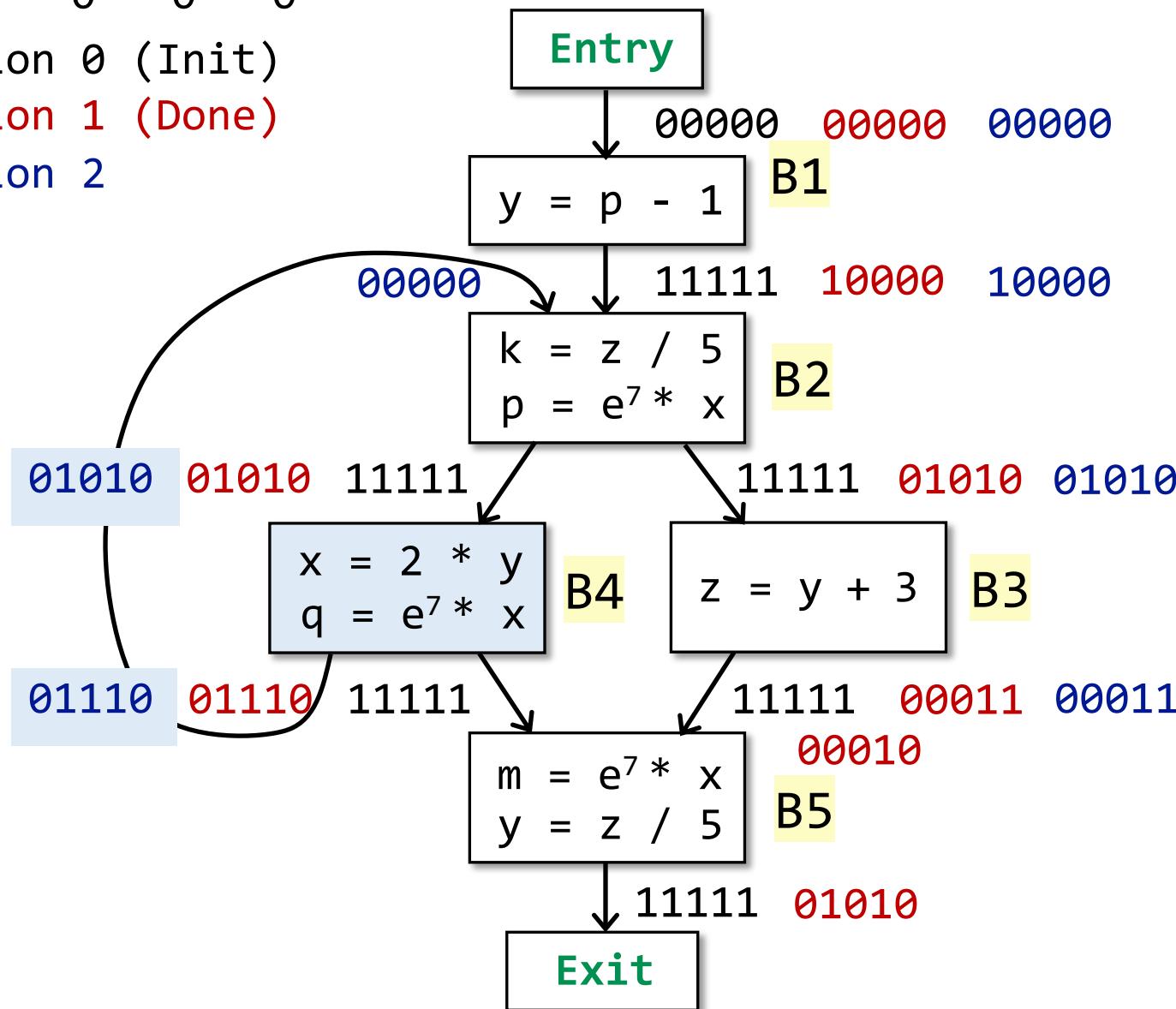


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

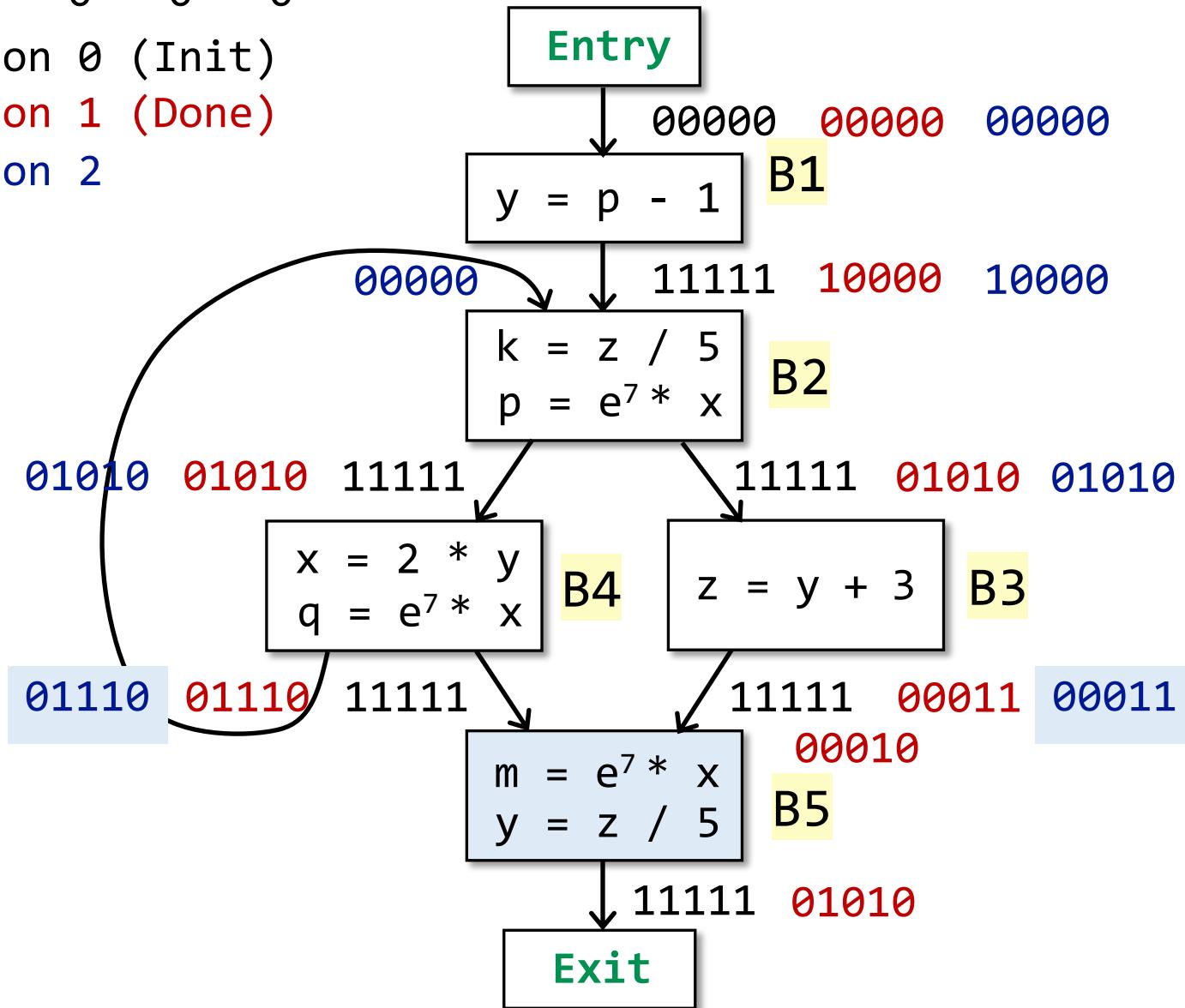


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

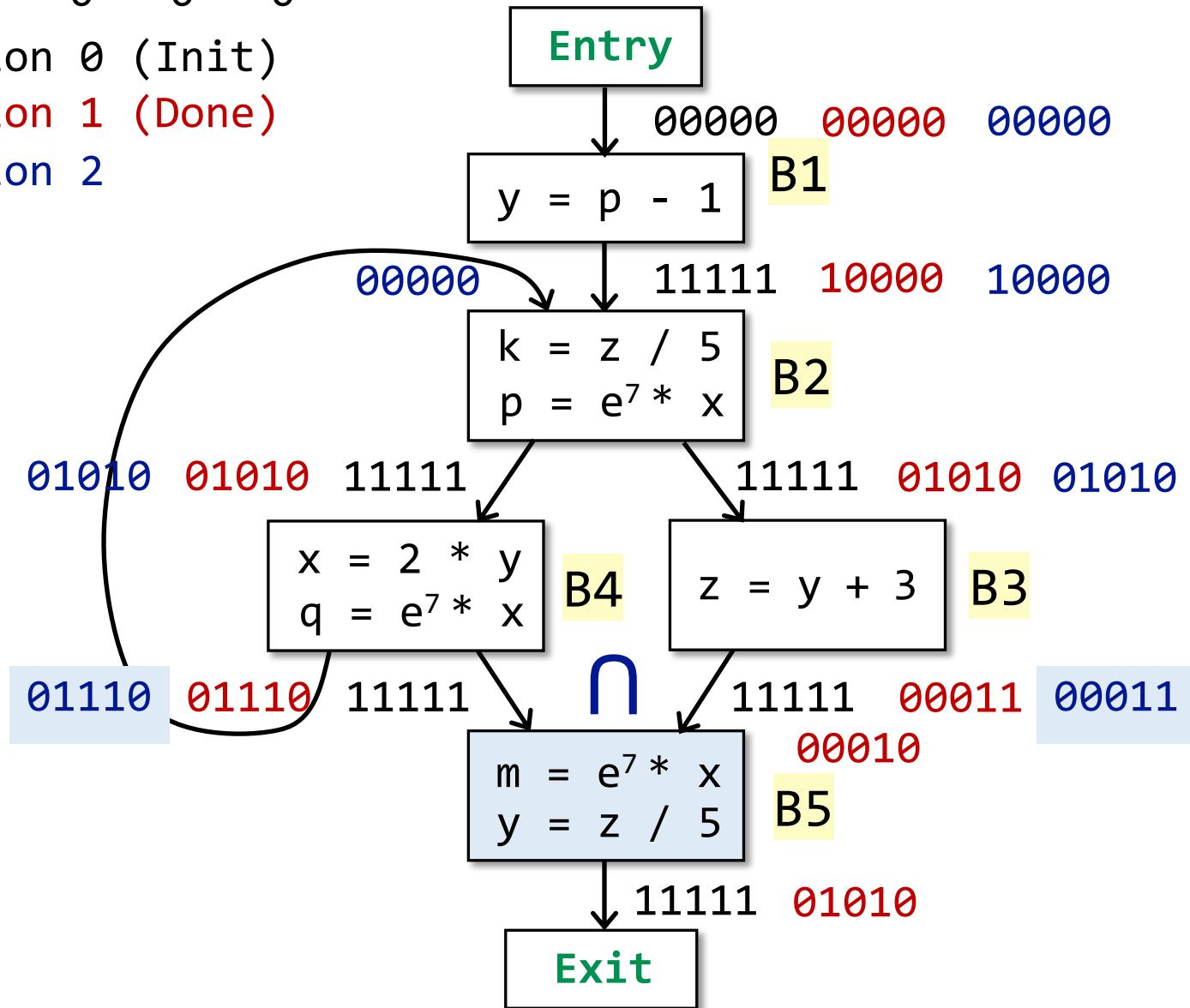


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

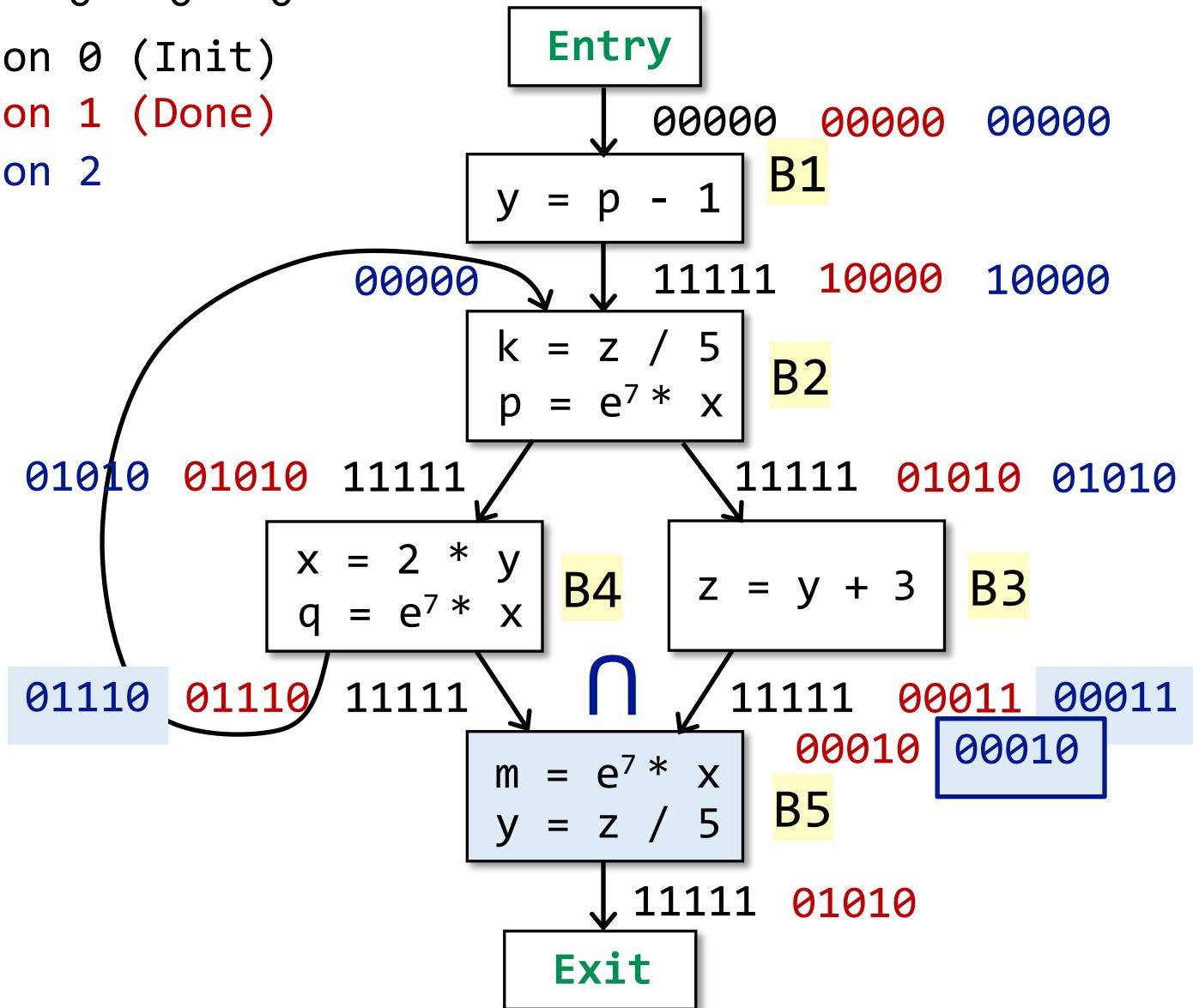


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

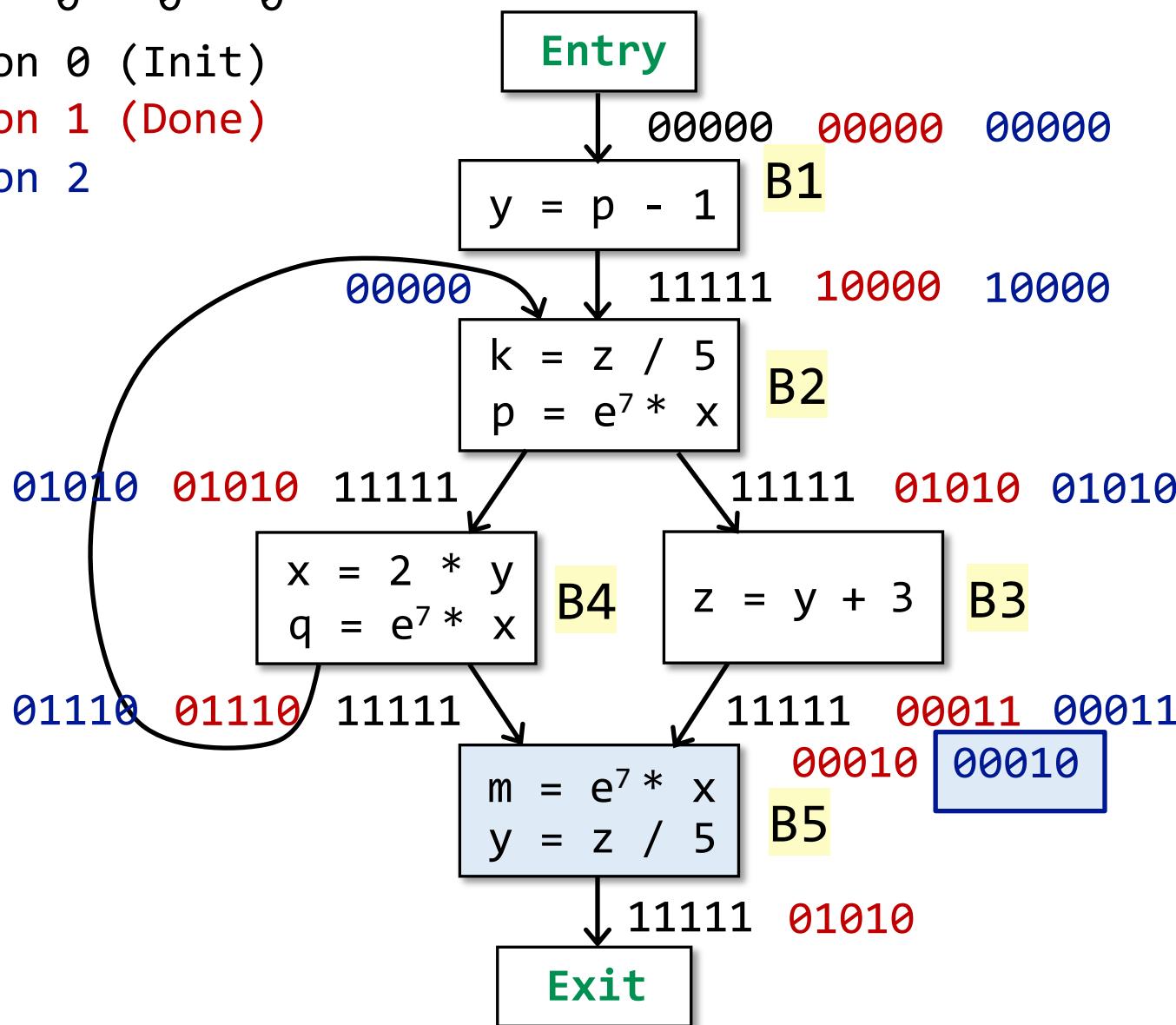


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

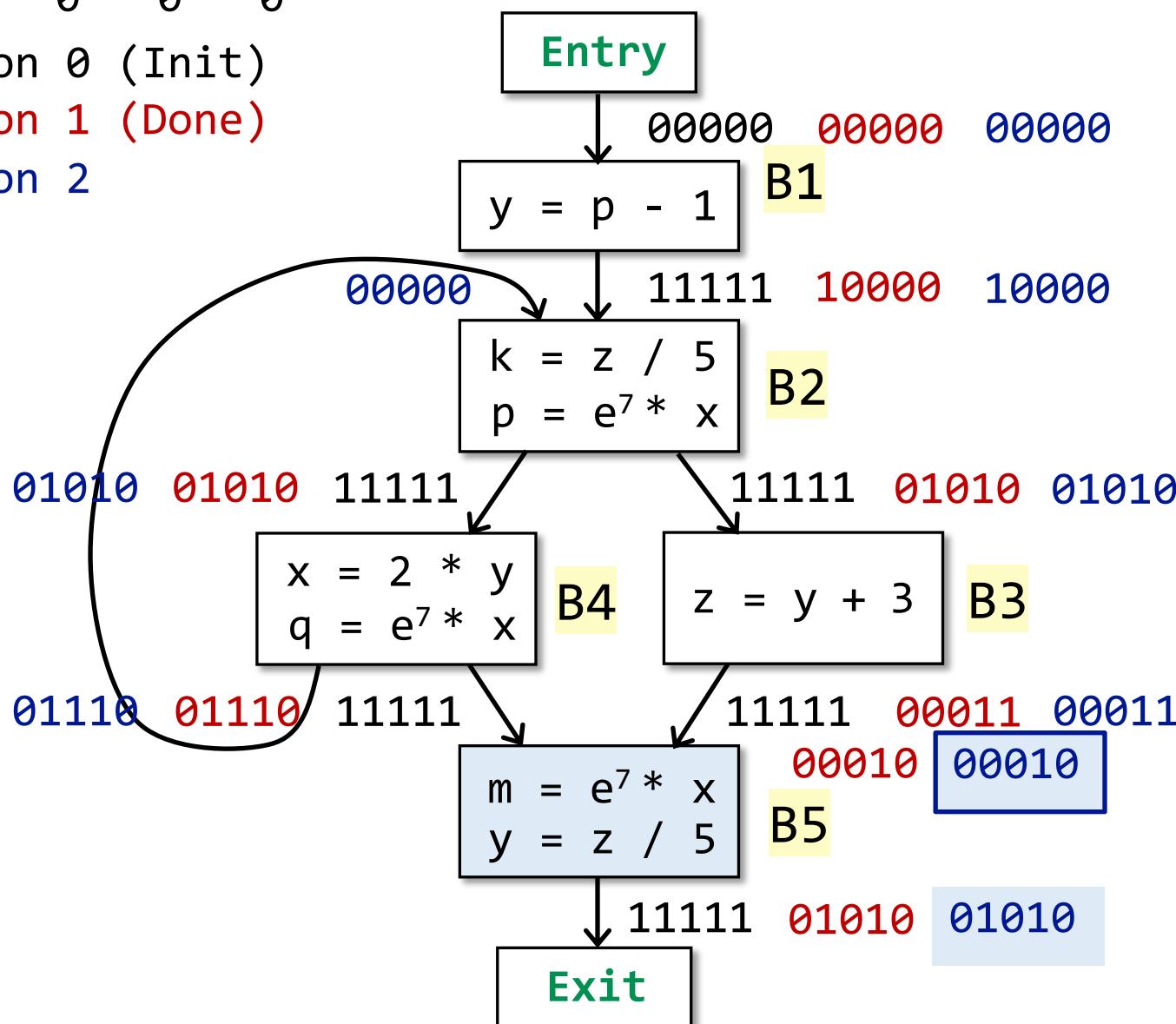


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

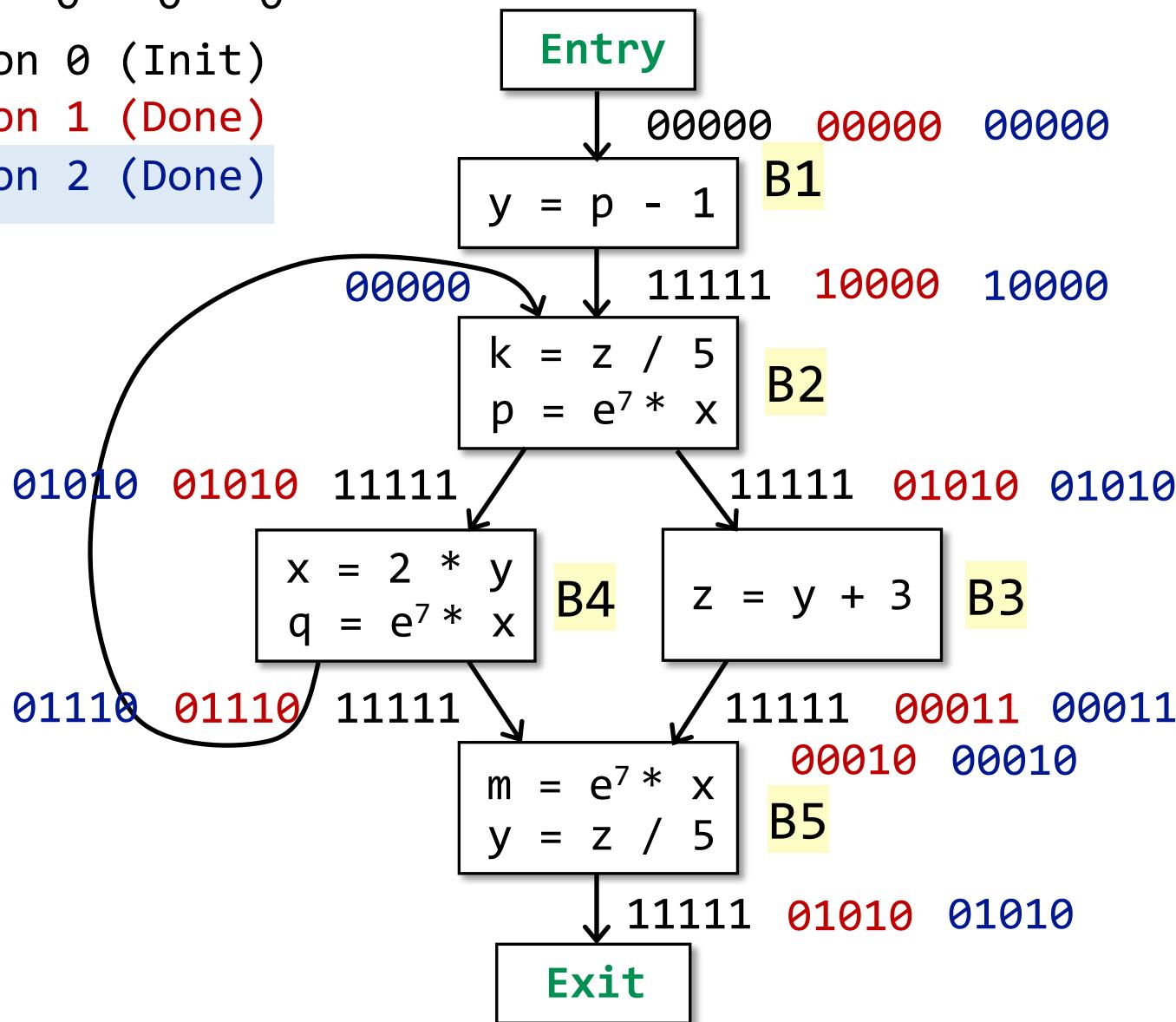


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

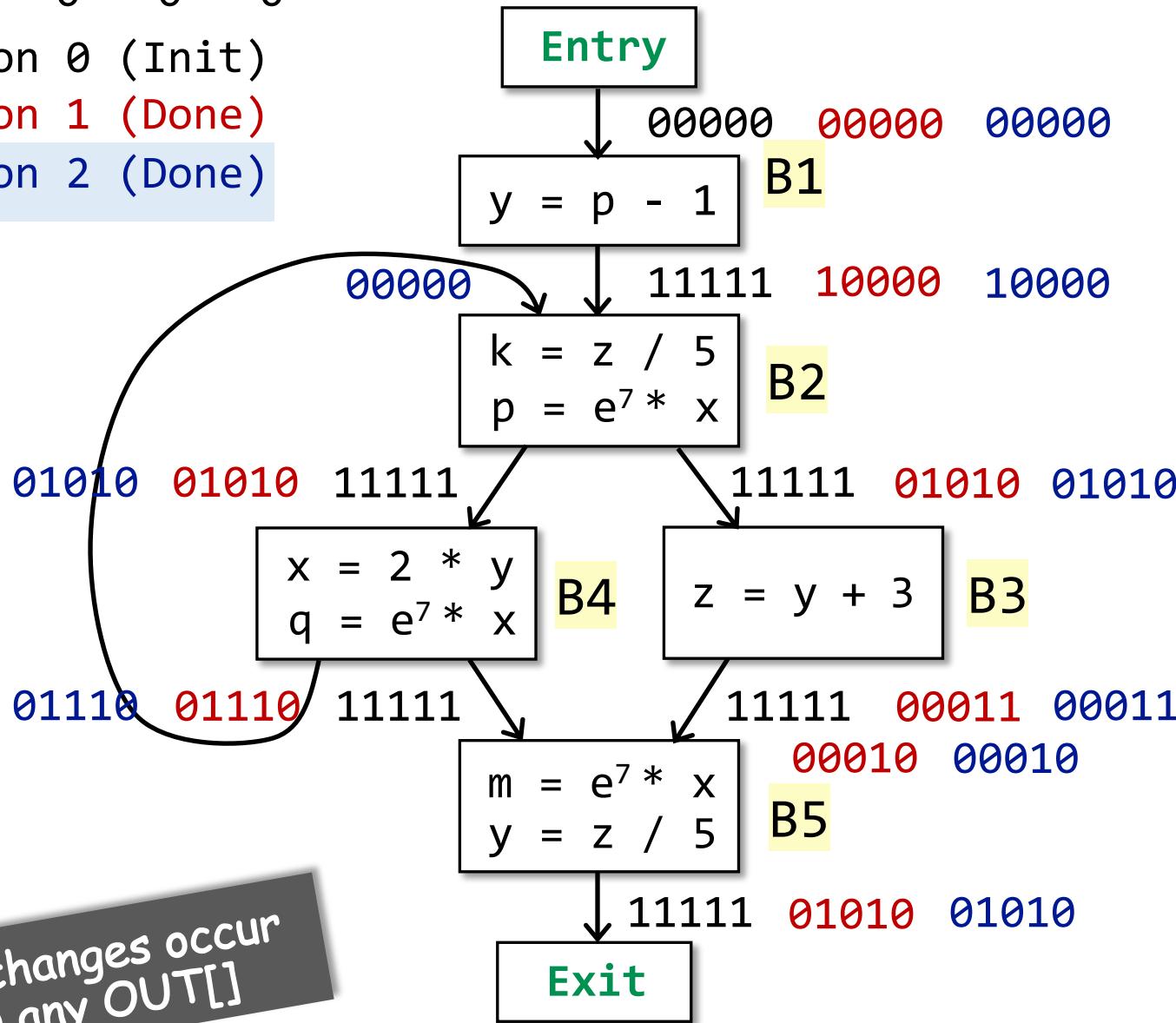


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

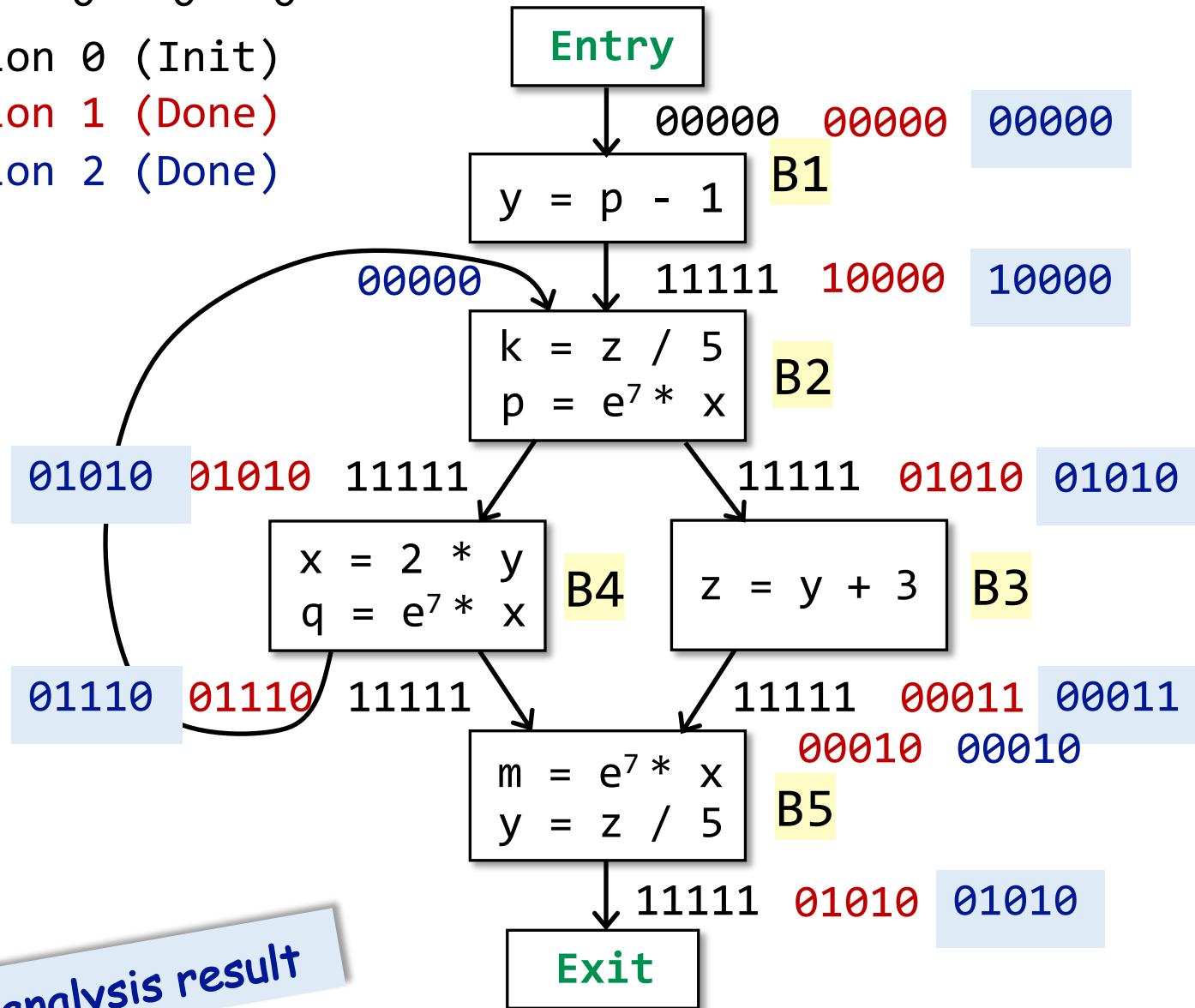


$p - 1$   $z / 5$   $2 * y$   $e^7 * x$   $y + 3$   
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)



# Analysis Comparison

|                   | Reaching Definitions | Live Variables | Available Expressions |
|-------------------|----------------------|----------------|-----------------------|
| Domain            |                      |                |                       |
| Direction         |                      |                |                       |
| May/Must          |                      |                |                       |
| Boundary          |                      |                |                       |
| Initialization    |                      |                |                       |
| Transfer function |                      |                |                       |
| Meet              |                      |                |                       |

# Analysis Comparison

|                   | Reaching Definitions | Live Variables   | Available Expressions |
|-------------------|----------------------|------------------|-----------------------|
| Domain            | Set of definitions   | Set of variables | Set of expressions    |
| Direction         |                      |                  |                       |
| May/Must          |                      |                  |                       |
| Boundary          |                      |                  |                       |
| Initialization    |                      |                  |                       |
| Transfer function |                      |                  |                       |
| Meet              |                      |                  |                       |

# Analysis Comparison

|                   | Reaching Definitions | Live Variables   | Available Expressions |
|-------------------|----------------------|------------------|-----------------------|
| Domain            | Set of definitions   | Set of variables | Set of expressions    |
| Direction         | Forwards             | Backwards        | Forwards              |
| May/Must          |                      |                  |                       |
| Boundary          |                      |                  |                       |
| Initialization    |                      |                  |                       |
| Transfer function |                      |                  |                       |
| Meet              |                      |                  |                       |

# Analysis Comparison

|                   | Reaching Definitions | Live Variables   | Available Expressions |
|-------------------|----------------------|------------------|-----------------------|
| Domain            | Set of definitions   | Set of variables | Set of expressions    |
| Direction         | Forwards             | Backwards        | Forwards              |
| May/Must          | May                  | May              | Must                  |
| Boundary          |                      |                  |                       |
| Initialization    |                      |                  |                       |
| Transfer function |                      |                  |                       |
| Meet              |                      |                  |                       |

# Analysis Comparison

|                   | Reaching Definitions                   | Live Variables                       | Available Expressions                  |
|-------------------|--|--------------------------------------|--|
| Domain            | Set of definitions                     | Set of variables                     | Set of expressions                     |
| Direction         | Forwards                               | Backwards                            | Forwards                               |
| May/Must          | May                                    | May                                  | Must                                   |
| Boundary          | $\text{OUT}[\text{entry}] = \emptyset$ | $\text{IN}[\text{exit}] = \emptyset$ | $\text{OUT}[\text{entry}] = \emptyset$ |
| Initialization    |  |                                      |  |
| Transfer function |  |                                      |  |
| Meet              |  |                                      |  |

# Analysis Comparison

|                   | Reaching Definitions                   | Live Variables                       | Available Expressions                  |
|-------------------|--|--------------------------------------|--|
| Domain            | Set of definitions                     | Set of variables                     | Set of expressions                     |
| Direction         | Forwards                               | Backwards                            | Forwards                               |
| May/Must          | May                                    | May                                  | Must                                   |
| Boundary          | $\text{OUT}[\text{entry}] = \emptyset$ | $\text{IN}[\text{exit}] = \emptyset$ | $\text{OUT}[\text{entry}] = \emptyset$ |
| Initialization    | $\text{OUT}[B] = \emptyset$            | $\text{IN}[B] = \emptyset$           | $\text{OUT}[B] = U$                    |
| Transfer function |  |                                      |  |
| Meet              |  |                                      |  |

# Analysis Comparison

|                   | Reaching Definitions                   | Live Variables   | Available Expressions                  |
|-------------------|--|--|--|
| Domain            | Set of definitions                     | Set of variables                                       | Set of expressions                     |
| Direction         | Forwards                               | Backwards  | Forwards                               |
| May/Must          | May                                    | May  | Must                                   |
| Boundary          | $\text{OUT}[\text{entry}] = \emptyset$ | $\text{IN}[\text{exit}] = \emptyset$                   | $\text{OUT}[\text{entry}] = \emptyset$ |
| Initialization    | $\text{OUT}[B] = \emptyset$            | $\text{IN}[B] = \emptyset$                             | $\text{OUT}[B] = U$                    |
| Transfer function |  | $\text{OUT} = \text{gen } U (\text{IN} - \text{kill})$ |  |
| Meet              |  |  |  |

# Analysis Comparison

|                   | Reaching Definitions                   | Live Variables   | Available Expressions                  |
|-------------------|--|--|--|
| Domain            | Set of definitions                     | Set of variables                                       | Set of expressions                     |
| Direction         | Forwards                               | Backwards  | Forwards                               |
| May/Must          | May                                    | May  | Must                                   |
| Boundary          | $\text{OUT}[\text{entry}] = \emptyset$ | $\text{IN}[\text{exit}] = \emptyset$                   | $\text{OUT}[\text{entry}] = \emptyset$ |
| Initialization    | $\text{OUT}[B] = \emptyset$            | $\text{IN}[B] = \emptyset$                             | $\text{OUT}[B] = U$                    |
| Transfer function |  | $\text{OUT} = \text{gen } U (\text{IN} - \text{kill})$ |  |
| Meet              | $U$                                    | $U$  | $\cap$                                 |

# Analysis Comparison

|                   | Reaching Definitions                   | Live Variables   | Available Expressions                  |
|-------------------|--|--|--|
| Domain            | Set of definitions                     | Set of variables                                       | Set of expressions                     |
| Direction         | Forwards                               | Backwards  | Forwards                               |
| May/Must          | May                                    | May  | Must                                   |
| Boundary          | $\text{OUT}[\text{entry}] = \emptyset$ | $\text{IN}[\text{exit}] = \emptyset$                   | $\text{OUT}[\text{entry}] = \emptyset$ |
| Initialization    | $\text{OUT}[B] = \emptyset$            | $\text{IN}[B] = \emptyset$                             | $\text{OUT}[B] = U$                    |
| Transfer function |  | $\text{OUT} = \text{gen } U (\text{IN} - \text{kill})$ |  |
| Meet              | $U$                                    | $U$  | $\cap$                                 |

# Analysis Comparison

According to the meaning of the analysis  
We'll draw a theoretical framework to systematically explain them in next lectures

|                   | Reaching Definitions                   | Live Variables   | Available Expressions                  |
|-------------------|--|--|--|
| Domain            | Set of definitions                     | Set of variables                                       | Set of expressions                     |
| Direction         | Forwards                               | Backwards  | Forwards                               |
| May/Must          | May                                    | May  | Must                                   |
| Boundary          | $\text{OUT}[\text{entry}] = \emptyset$ | $\text{IN}[\text{exit}] = \emptyset$                   | $\text{OUT}[\text{entry}] = \emptyset$ |
| Initialization    | $\text{OUT}[B] = \emptyset$            | $\text{IN}[B] = \emptyset$                             | $\text{OUT}[B] = U$                    |
| Transfer function |  | $\text{OUT} = \text{gen } U (\text{IN} - \text{kill})$ |  |
| Meet              | U                                      | U  | n                                      |

# summary

1. Overview of Data Flow Analysis
2. Preliminaries of Data Flow Analysis
3. Reaching Definitions Analysis
4. Live Variables Analysis
5. Available Expressions Analysis

# The X You Need To Understand in This Lecture

- Understand the three data flow analyses:
  - reaching definitions
  - live variables
  - available expressions
- Can tell the differences and similarities of the three data flow analyses
- Understand the iterative algorithm and can tell why it is able to terminate

注意注意！  
划重点了！

